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Homological Algebra

(Master course)

Bernard Le Stum (November 5, 2025)

A decorative footer image featuring a bokeh effect with numerous out-of-focus yellow and orange circles of varying sizes against a dark orange background.

– If you're capable of doing the math yourself with confidence and efficiency, you're often better off doing it yourself.

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Introduction

Teaser

Invariant theory associates a simple mathematical object to a more complex one. For instance, one can associate to a set E its number of elements $\#E$.

Invariant theory is frequently employed to establish impossibilities. For instance, the Cantor-Bernstein theorem demonstrates that when the cardinality of a set E is less than that of another set F ($\#E < \#F$), there exists no injective map from F to E .

Another basic concept in invariant theory is the dimension $\dim V$ of a finite-dimensional real vector space V . It can be shown that when the dimensions of V and another vector space W are different ($\dim V \neq \dim W$), there exists no homeomorphism between V and W (Brouwer's open mapping theorem).

However, to prove this theorem, a more sophisticated invariant is required: to a topological manifold, one can associate a family of abelian groups $H_i(X)$ indexed by natural numbers, called homology groups, which is invariant under homeomorphisms.

From a finite-dimensional vector space V endowed with an Euclidean norm, one can construct a sphere $S := (V \setminus \{0\})/\| - \|$. It can then be shown that $H_i(S) = 0$ unless $i = 0$ or $i = \dim(V) - 1$ in which case it is free of rank one. This result leads to Brouwer's theorem.

Homology and cohomology theories provide invariants of abelian nature (for instance, abelian groups) for objects of geometrical nature (for example, manifolds). The purpose of homological algebra is to elucidate the theory behind these invariants. It appears that the most suitable setting for this study is category theory.

It is common in mathematics to present an object by generators and relations.

For example, the field \mathbb{C} of complex numbers may be seen as the \mathbb{R} -algebra generated by an element x satisfying $x^2 + 1 = 0$. We may think of $x = i$ but can as well choose $x = -i$ if we wish. In any case, we can write $\mathbb{C} = \mathbb{R}[X]/(X^2 + 1)$: it means that there exists a surjection from the polynomial ring $\mathbb{R}[X]$ whose kernel is the ideal generated by $X^2 + 1$. More generally, any commutative algebra over a commutative ring R may be written on the form $R[\{X_i\}_{i \in I}]/(f_j)_{j \in J}$.

Let us now consider the dihedral group D_n . This is the group generated by (a rotation of order n) r and (a reflection) s subject to the relations $r^n = 1$, $s^2 = 1$ and $rsrs = 1$. One usually writes for $n \in \mathbb{N}$,

$$D_n = \langle r, s \mid r^n = 1, s^2 = 1, rsrs = 1 \rangle = \langle R, S \rangle / \langle R^n, S^2, RSRS \rangle.$$

It means that there exists a surjection from the free group generated by two elements R, S whose kernel is the subgroup generated by $R^n, S^2, RSRS$. This group in turn is the target of a surjection from the free group on three elements X, Y, Z sending X to R^n , Y to S^2 and Z to $RSRS$. More generally, any group has a free *presentation*

$$F' \xrightarrow{\psi} F \xrightarrow{\phi} G \rightarrow 0 :$$

it means that both F and F' are free groups, $\text{im } \psi = \ker \phi$ and ϕ is surjective ($\text{im } \phi = \ker 0$).

We now consider the *commutative* group μ_n of n th roots of unity for some fixed integer n . If we choose a primitive n th root of unity ζ , then there exists a surjection $\mathbb{Z} \rightarrow \mu_n$ sending k to ζ^k . Of course, \mathbb{Z} is a free abelian group. But the kernel $n\mathbb{Z}$ of our surjective map also is a free abelian group. This is a general fact: any commutative group has a free (*left*) *resolution*

$$0 \rightarrow P_1 \xrightarrow{i} P_0 \xrightarrow{p} M \rightarrow 0$$

(of length 2). It means that P_0 and P_1 are free abelian groups, i is injective ($\text{im } 0 = \ker i$), $\text{im } i = \ker p$ and p is surjective ($\text{im } p = \ker 0$). One also calls $P_1 \rightarrow P_0$ a free (*left*) *replacement* for M .

Before going any further, let us consider the effect of a “transformation” on a free resolution. For example, for a fixed integer n , we may turn an abelian group M into its quotient M/n and consequently, any morphism of abelian groups $f : M \rightarrow N$ into the corresponding map $\bar{f} : M/n \rightarrow N/n$. If we apply this transformation to the resolution $0 \rightarrow \mathbb{Z} \xrightarrow{n} \mathbb{Z} \rightarrow \mathbb{Z}/n\mathbb{Z} \rightarrow 0$, then we get $0 \rightarrow \mathbb{Z}/n\mathbb{Z} \xrightarrow{0} \mathbb{Z}/n\mathbb{Z} \xrightarrow{\sim} \mathbb{Z}/n\mathbb{Z} \rightarrow 0$ which is *not* a resolution anymore because the first map is not injective. The point is that, if we are given $M \xrightarrow{f} N \xrightarrow{g} P$, it may happen that $\text{im } f = \ker g$ but $\text{im } \bar{f} \neq \ker \bar{g}$. This is called a *default of exactness*. However, it should be noticed that our transformation is compatible with composition: we have $\overline{g \circ f} = \bar{g} \circ \bar{f}$. As a consequence, the fact that $\text{im } f \subset \ker g$ will imply $\text{im } \bar{f} \subset \ker \bar{g}$.

Let us now consider the case of a (left) R -module M where R is a ring. In order to generalize the previous case (which is essentially the case $R = \mathbb{Z}$), we need to allow 1) direct summands of free modules, which are called *projective* modules and 2) longer, possibly infinite, *projective* (left) *resolutions*

$$\cdots \rightarrow P_2 \xrightarrow{d_2} P_1 \xrightarrow{d_1} P_0 \xrightarrow{d_0} M \rightarrow 0,$$

meaning that $\text{im } d_{n+1} = \ker d_n$ (so that, in particular, $\text{im } d_0 = \ker 0 = M$ which means that d_0 is surjective). And we shall call $P_\bullet := \cdots \rightarrow P_2 \xrightarrow{d_2} P_1 \xrightarrow{d_1} P_0$ a projective (left) replacement for M .

Again, a transformation will not preserve resolutions in general but it will preserve chain complexes (that we will shortly define) as long as the transformation is compatible with composition. A *chain complex* is a sequence of maps

$$K_\bullet := \cdots \rightarrow K_2 \xrightarrow{d_2} K_1 \xrightarrow{d_1} K_0 \xrightarrow{d_0} K_{-1} \rightarrow \cdots$$

satisfying $\text{im } d_{n+1} \subset \ker d_n$, or equivalently $d_n \circ d_{n+1} = 0$. In order to measure the default of exactness, we may then introduce the *homology* groups

$$H_n(K_\bullet) := \ker d_n / \text{im } d_{n+1}.$$

Let us briefly explain how these techniques can be used in topology. We shall not enter the details at this point, but one may consider on any topological space X , the notion of an abelian sheaf M : it is a compatible¹ local² family of abelian groups $\Gamma(U, M)$ for all open subsets U of X . As a baby example, one may consider the sheaf $\mathcal{C}_\mathbb{R}$ of continuous maps with values in \mathbb{R} so that $\Gamma(U, \mathcal{C}_\mathbb{R}) = \mathcal{C}(U, \mathbb{R})$ or the *constant sheaf* \mathbb{Z}_X given by

$$\Gamma(U, \mathbb{Z}_X) = \mathcal{C}(U, \mathbb{Z}) \simeq \mathbb{Z}^{\pi_0(U)}.$$

It happens that there are *not* enough projective abelian sheaves in general. We need to rely on the dual construction and introduce the notions of an *injective* abelian sheaf as well as an injective (right) resolution

$$0 \rightarrow M \rightarrow I^0 \rightarrow I^1 \rightarrow I^2 \rightarrow \cdots$$

(and injective (right) replacement). One also replaces chain complex with *cochain complex*

$$K^\bullet := \cdots \rightarrow K^{-1} \xrightarrow{d_{-1}} K^0 \xrightarrow{d_0} K^1 \xrightarrow{d_1} K^2 \rightarrow \cdots$$

meaning $d_n \circ d_{n-1} = 0$ and homology with *cohomology*

$$H^n(K^\bullet) = \ker d_n / \text{im } d_{n-1}.$$

One can show that if I^\bullet is an injective replacement for an abelian sheaf M , then

$$H^n(X, M) := H^n(\Gamma(X, I^\bullet))$$

depends only on M . This is the n th cohomology group of M . When X is locally contractible, we recover the singular cohomology of X :

$$H^n(X, \mathbb{Z}_X) \simeq H_{\text{sing}}^n(X).$$

¹There exists restrictions maps $s \mapsto s|_V$ for s on U and $V \subset U$ with $(s|_V)_W = s|_W$ and $s|_U = s$.

²Given s_i on U_i such that s_i and s_j coincide on $U_i \cap U_j$, then there exists a unique s on $\bigcup U_i$ extending all s_i simultaneously.

Sheaf cohomology however is a lot more flexible since it allows coefficients in any abelian sheaf (and not merely abelian group) and this may even be extended to cochain complexes of abelian sheaves. There also exists a relative version: given any sheaf M on X and any continuous map $f : X \rightarrow Y$, there exists cohomology sheaves $R^n f_* M$ on Y which may be used to compute the cohomology of M . There is no such thing as singular cohomology.

It is also possible to apply these methods to differential geometry. On a differential manifold X , there exists the de Rham complex:

$$\Omega_X^\bullet := \mathcal{O}_X \xrightarrow{d} \Omega_X^1 \xrightarrow{d} \cdots \xrightarrow{d} \Omega_X^n.$$

Poincaré lemma states that this is actually a (right) replacement for the constant sheaf \mathbb{R}_X . The cohomology of the complex of global sections

$$\Gamma(X, \Omega_X^\bullet) := \Gamma(X, \mathcal{O}_X) \xrightarrow{d} \Gamma(X, \Omega_X^1) \xrightarrow{d} \cdots \xrightarrow{d} \Gamma(X, \Omega_X^n)$$

is called de Rham cohomology. It happens actually that

$$H_{\text{dR}}^n(X) := H^n(\Gamma(X, \Omega_X^\bullet)) \simeq H^n(X, \Omega_X^\bullet) \simeq H^n(X, \mathbb{R}_X) \simeq H_{\text{sing}}^n(X, \mathbb{R}),$$

showing that de Rham and singular cohomology coincide. In particular, a differential invariant appears to be a topological invariant. This generalizes to a Riemann-Hilbert correspondence between differential systems and local systems but this is another story.

Throughout the course, there will be numerous exercises. Some of these exercises will illustrate previous properties or definitions, while others will present general results that are usually straightforward to prove. It's important to ensure that you can complete all these exercises but also write down full solutions occasionally.

We will use a smaller font to indicate optional material that can be skipped.

Many thanks to the students who helped make this course more readable.

Background

We briefly review some fundamental concepts from general mathematics as well as specific developments that will be used throughout the course. Our main goal is to refresh the reader's memory with familiar definitions and standard properties, as well as introduce some notation. We assume that the reader is proficient in elementary logical operations, including quantifiers, and has mastered the induction process. Additionally, we expect the reader to be familiar with the concepts of sets, elements, inclusion, intersection, union, and (possibly infinite) product³ of sets. Also, we will not revisit the definitions and basic properties of \mathbb{N} , \mathbb{Z} , \mathbb{Q} , \mathbb{R} , or \mathbb{C} nor delve into issues related to logic or set theory. Finally, we shall use von Neumann notation $n := \{0, 1, \dots, n - 1\}$.

0.1 Relation

A (binary) relation $R : X \rightarrow Y$ is the data of two sets X, Y called respectively *domain* and *codomain*, and another set $R \subset X \times Y$ called the *graph* of the relation. We will write xRy instead of $(x, y) \in R$ and call y the *image* of x (resp. x the *preimage* of y).

There exists the notion of a (direct) image by R of a subset A of X which is defined as⁴

$$R(A) := \{y \in Y, \exists x \in X, xRy\}.$$

We also set $\text{im}(R) := R(X)$. We shall simply write $R(x) := R(\{x\})$ and more generally remove all embedded parenthesis when there is no risk of confusion.

³Let us however recall that the product $\prod_{i \in I} X_i$ is the set of all families $(x_i)_{i \in I}$ of terms $x_i \in X_i$ for all $i \in I$ and we shall write $X^I := \prod_I X$.

⁴To avoid confusion later, we should write $R_*(A)$ in which case, we would write $R^*(B) = R_*^{-1}(B)$ below.

There exists an *inverse* relation $R^{-1} : Y \rightarrow X$ for R obtained by switching factors so that

$$\forall x \in X, y \in Y, \quad yR^{-1}x \Leftrightarrow xRy.$$

If $B \subset Y$, then $R^{-1}(B)$ is called the *inverse image* of B by R .

One can make the *product* of relations $\prod_{i \in I} R_i : \prod_{i \in I} X_i \rightarrow \prod_{i \in I} Y_i$ so that

$$(x_i)_{i \in I} \left(\prod_{i \in I} R_i \right) (y_i)_{i \in I} \Leftrightarrow \forall i \in I, x_i R_i y_i.$$

We may consider the *conjunction* $R := \bigcap_{i \in I} R_i : X \rightarrow Y$ of relations so that

$$\forall x \in X, y \in Y, \quad xRy \Leftrightarrow \forall i \in I, xR_i y.$$

A relation $R' : X' \rightarrow Y'$ is a *restriction* (resp. *the restriction*) of R to $X' \subset X$ and $Y' \subset Y$ if

$$\forall x \in X', y \in Y', \quad xR'y \Rightarrow xRy \quad (\text{resp. } \Leftrightarrow xRy).$$

We may also say that R' is *induced* by R or that R is an *extension* of R' .

One can also *compose* a relation $R : X \rightarrow Y$ with a relation $S : Y \rightarrow Z$ by setting

$$x(S \circ R)z \Leftrightarrow \exists y \in Y, xRy \text{ and } ySz.$$

A map $f : X \rightarrow Y$ is a relation satisfying

$$\forall x \in X, \exists y \in Y, \quad f(\{x\}) = \{y\}.$$

We will then write $f(x) = y$ or $f : x \mapsto y$.

For a set X , we shall consider the set $\mathcal{P}(X)$ of subsets of X so that $A \subset X \Leftrightarrow A \in \mathcal{P}(X)$. If $R : X \rightarrow Y$ is any relation, then the direct and inverse images introduced above are maps

$$R : \mathcal{P}(X) \rightarrow \mathcal{P}(Y), \quad A \mapsto R(A) \quad \text{and} \quad R^{-1} : \mathcal{P}(Y) \rightarrow \mathcal{P}(X), \quad B \mapsto R^{-1}(B)$$

and any of them uniquely determines the relation.

A map $f : X \rightarrow Y$ is said to be *injective* (resp. *surjective*, resp. *bijective*) if, for all $y \in Y$, $f^{-1}(y)$ has at most (resp. at least, resp. exactly) one element. The map f is bijective if and only if f^{-1} is a map.

We shall denote by $\mathcal{F}(X, Y)$ (resp. $\mathcal{S}(X, Y)$) the set of all maps (resp. bijective maps) from X to Y and $\mathcal{F}(X)$ (resp. $\mathcal{S}(X)$) when $X = Y$. There always exists a bijection

$$\mathcal{F}(X, Y) \simeq Y^X, \quad f \mapsto (f(x))_{x \in X}.$$

It means that a map is essentially the same thing as a family. By the way, to any set X , we can associate a family $(x)_{x \in X}$ corresponding to the map Id_X and to any family $(x_i)_{i \in I}$, corresponding to the map $f : I \rightarrow X, i \mapsto x_i$, we can associate the set $\text{im}(f) = \{x_i\}_{i \in I}$.

Also, there always exists a bijection (called *currying*)

$$\mathcal{F}(X \times Y, Z) \simeq \mathcal{F}(X, \mathcal{F}(Y, Z)), \quad f \mapsto (x \mapsto (f_x : y \mapsto f(x, y))).$$

The maps f_x are called *partial* maps. A map $X \times Y \rightarrow Z$ is also called an *operation* of X on Y with values in Z . This is essentially the same thing as a map $X \mapsto \mathcal{F}(Y, Z)$.

0.2 Relation in a set

A *relation* in a set X is a relation $R : X \rightarrow X$. For example, the *identity* Id_X , also called *equality*, is the relation in X defined by $x\text{Id}_X y \Leftrightarrow x = y$.

A relation in a set X is said to be

1. *reflexive* if $\forall x \in X, xRx$,
2. *transitive* if $\forall x, y, z \in X, (xRy \text{ and } yRz) \Rightarrow xRz$,
3. *symmetric* if $\forall x, y \in X, xRy \Rightarrow yRx$,
4. *antisymmetric* if $\forall x, y \in X, (xRy \text{ and } yRx) \Rightarrow x = y$,
5. *total* (or *linear*) if $\forall x, y \in X, xRy$ or yRx ,

All five properties are inherited by (the restriction to) any subset and satisfied by the opposite relation. The first four properties are stable under conjunction. Therefore, there always exists a smallest relation in X extending R and having one or more of these properties (but not always the last one).

A *preorder* is a relation in X which is both reflexive and transitive. This is called an *equivalence relation* (resp. a *partial order*) if, moreover, it is symmetric (resp. antisymmetric). A set endowed with a partial order is called a *poset* but we shall concentrate on preordered sets. A preorder \leq on $X \neq \emptyset$ is *directed* if $\forall x, y \in X, \exists z \in X, x \leq z$ and $y \leq z$. A set endowed with a directed partial order is called a *directed set*. An element x of a preordered set (X, \leq) is said to be *maximal* if $\forall y \in X, x \leq y \Rightarrow x = y$ (and *minimal* if it is maximal for the opposite preorder \geq).

An element x of a preordered set (X, \leq) is an *upper bound* for a subset $A \subset X$ if $\forall y \in A, y \leq x$. An upper bound for X itself is called a *maximum* (or *greatest element*). There exists the opposite notions of *lower bound* and *minimum*. The *supremum* (or *join*) $\sup(A)$ (resp. *infimum* (or *meet*) $\inf(A)$) is the least upper bound (greatest lower bound). An ordered set is said to be *inductive* if any totally ordered subset has an upper bound. Zorn's lemma states that any inductive set has a maximal element.

If R is a relation on X and \tilde{R} denotes the equivalence relation generated by R , then the *quotient* of X by R is $X/R := \{\tilde{R}(x), x \in X\}$ and $\tilde{R}(x)$ is called an *equivalence class*. There exists an obvious surjective map

$$\pi : X \twoheadrightarrow X/R, \quad x \mapsto \tilde{R}(x)$$

called the *quotient map*. A map $f : X \rightarrow Y$ is said to be *compatible* with relations R and S in X and Y respectively if

$$\forall x, x' \in X, \quad xRx' \Rightarrow f(x)Sf(x').$$

When this is the case, f provides a map $\bar{f} : X/R \rightarrow Y/S$. And conversely when S is an equivalence relation. When a map is compatible with partial orders, it is said to be *order preserving*.

The *direct image* $f(R)$ of a relation R on X by a relation $f : X \rightarrow Y$ is defined by

$$\forall y, y' \in Y, \quad yf(R)y' \Leftrightarrow \exists x, x' \in X, xRx', xfy, x'fy'.$$

A direct image by f^{-1} is also called an *inverse image*. We shall only use this notion in the case f is a map in which case f will be compatible with R and $f(R)$ (and

with $f^{-1}(S)$ and S if S is a relation on Y). The inverse image under a map of a preorder (resp. an equivalence relation) is still a preorder (resp. an equivalence relation). In particular, the inverse image of equality on Y is an equivalence relation on X . Conversely, if \sim is an equivalence relation on X , then \sim is the inverse image of equality under the quotient map $\pi : X \rightarrow X/\sim$. Finally, if \leq is a preorder on a set X , then the image of \leq is a partial order on X/\sim where $x \sim y \Leftrightarrow x \leq y$ and $y \leq x$.

A *filter* on a partially ordered set P is a non empty proper subset $\mathcal{F} \subsetneq P$ such that

1. $\forall x, y \in \mathcal{F}, \exists z \in \mathcal{F}, z \leq x$ and $z \leq y$,
2. $\forall x \in \mathcal{F}, \forall y \in X, x \leq y \Rightarrow y \in \mathcal{F}$.

An *ultrafilter* is a maximal filter (for inclusion). If I is any set, then a filter on $\mathcal{P}(I)$ is also called a *filter* on I (this may be confusing when I itself is endowed with some preorder).

0.3 Group

A *magma* is a set G endowed with an internal operation

$$G \times G \rightarrow G, \quad (g, h) \mapsto gh$$

or, equivalently, a map

$$G \rightarrow \mathcal{F}(G), \quad g \mapsto (h \mapsto gh).$$

The notion of (non empty finite) *product* $g_1 \dots g_n$ is defined by induction. Note that $\mathcal{P}(G)$ then becomes also a magma for the rule $AB := \{gh, g \in A, h \in B\}$.

A *homomorphism of magmas* is a map $\varphi : G \rightarrow G'$ that preserves nonempty finite products. The homomorphism φ is called an *isomorphism* if it is bijective. If G is a magma, then the *opposite magma* is the magma G^{op} with the same elements as G but multiplication $(g, h) \mapsto hg$. A magma G is said to be *commutative* or *abelian* if $G^{\text{op}} = G$. A magma is called a *semigroup* if multiplication is *associative*:

$$\forall g, h, k \in G, \quad (gh)k = g(hk).$$

A *monoid* is a semigroup which is *unitary*:

$$\exists 1 \in G, \forall g \in G, \quad 1g = g = g1.$$

By convention, the *empty product* is 1. A *homomorphism of monoids* is required to preserve *all* finite products (including the empty product). A *submonoid* of a monoid G is a subset which is stable under (all) finite products. Any intersection of submonoids is a submonoid and we may always consider the submonoid generated by a given subset. If X is any set, then the set $\mathcal{F}(X)$ of all maps from X to itself is a monoid for composition. If X is any set and G is a monoid, then $\mathcal{F}(X, G)$ (or equivalently G^X) is a monoid for the termwise multiplication rule: $\forall x \in X, (\varphi\psi)(x) = \varphi(x)\psi(x)$. The subset $G \cdot X \subset \mathcal{F}(X, G)$ of maps with finite support (the support of φ is $\{x \in X, \varphi(x) \neq 1\}$) is a submonoid (and similarly $G^{(X)} \subset G^X$). More generally, given a family of monoids $\{G_i\}_{i \in I}$, the subset $\oplus_{i \in I} G_i \subset \prod_{i \in I} G_i$ of families $(g_i)_{i \in I}$ that are almost always 1, is a submonoid.

An element g of a monoid G is said to be *regular* (resp. *invertible*) if multiplication by g is injective (resp. bijective) on both sides. A monoid is *integral* (resp. a *group*) if all elements are regular (resp. invertible). A *subgroup* of a group G is a submonoid which is a group. In general, we shall denote by G^\times the group of invertible elements of a monoid G . Any group homomorphism $\varphi : G \rightarrow H$ induces a bijection between the set of subgroups of G containing the *kernel* $\ker \varphi := \varphi^{-1}(1)$ and the set of subgroups of H contained in the *image* $\text{im } \varphi := \varphi(G)$. If M and N are two abelian groups, then the set $\text{Hom}_{\text{Ab}}(M, N)$ of all homomorphisms from M to N is a subgroup of $\mathcal{F}(M, N)$. We shall write $\text{End}_{\text{Ab}}(M)$ in the case $M = N$. A subgroup of a free abelian group (isomorphic to $\mathbb{Z} \cdot I$ for some I) is also a free abelian group.

If G is a monoid, then a G -set is a set X endowed with a homomorphism of monoids

$$G \rightarrow \mathcal{F}(X), \quad g \mapsto (x \mapsto gx).$$

It means that X is endowed with an *action* of G (an operation with values in X itself) which is associative ($(gh)x = g(hx)$) and unitary ($1.x = x$). One may then define the notions of a G -morphism between G -sets as well as a *sub- G -set* in the usual way (exercise). If G is a group and X is a G -set, one may consider the *quotient* X/G which is the set of *orbits* $Gx := \{gx, g \in G\}$ as well as the quotient map

$$X \rightarrow X/G, \quad x \mapsto Gx.$$

As a particular case, if H is a subgroup of a group G , then H acts by translation on G and we may consider the quotients G/H as well as $H \backslash G := G^{\text{op}}/H^{\text{op}}$. The subgroup H is said to be *normal* in G if $G/H = H \backslash G$.

An abelian group is said to be *divisible* if $\forall a \in Q, \forall n \in \mathbb{N} \setminus 0, \exists x \in Q, a = nx$. A \mathbb{Q} -vector space is divisible. A quotient of a divisible group is divisible.

0.4 Ring and module

A (*unitary associative*) *ring* is an (additive) abelian group A endowed with a monoid structure that factors through a homomorphism of abelian groups

$$A \rightarrow \text{End}_{\text{Ab}}(A) \subset \mathcal{F}(A), \quad a \mapsto (b \mapsto ab).$$

Note that this extra condition simply means that *distributivity* holds on both sides:

$$\forall a, b, c \in A, \quad a(b + c) = ab + ac \quad \text{and} \quad (a + b)c = ac + bc.$$

A *ring homomorphism* is a map which is a homomorphism for both operations (*isomorphism* if it is bijective). A *subring* is a subset which is at the same time a subgroup of $(A, +)$ and a submonoid of (A, \times) . If M is an abelian group, then $\text{End}_{\text{Ab}}(M)$ is a ring for composition.

If A is a ring, a (*left*) A -module is an abelian group M endowed with a ring homomorphism $A \rightarrow \text{End}_{\text{Ab}}(M), a \mapsto (x \mapsto ax)$. One then defines the notions of an A -linear map between A -modules as well as a *sub- A -module* and *quotient A -module* in the usual way (exercise). The kernel of a morphism of A -modules $\varphi : M \rightarrow N$ is a

sub- A -module and its *cokernel* is defined as $\text{coker } \varphi := M/\text{im } \varphi$. One also defines the notion of submodule $\langle S \rangle$ generated by a subset S and call an A -module M *finite (type)* if it is generated by a finite subset. A (*left*) *ideal* \mathfrak{a} of A is a submodule of the A -module A . Be careful that, if \mathfrak{a} is an ideal of A and M an A -module, then $\mathfrak{a}M$ will denote the submodule *generated by* the naive product and not merely the product itself. A *right* A -module is a (left) A^{op} -module. The ring A is said to be *commutative* if the multiplicative monoid of A is commutative, or equivalently if $A^{\text{op}} = A$. When \mathfrak{a} is an ideal in a commutative ring A , there exists a unique structure of (commutative) ring on A/\mathfrak{a} making the quotient map $A \rightarrow A/\mathfrak{a}$ commutative.

When k is a *field* (meaning that $k \setminus \{0\}$ is a group), a k -module M is also called *vector space*. More generally, if k is a commutative ring, then the set $L(M, N)$ of all k -linear maps from M to N is automatically a k -module and $M^\vee := L(M, k)$ is called the *dual* of M . If M, N, P are k -modules, then a map $M \times N \rightarrow P$ is said to be *k-bilinear* if all partial maps are k -linear. The set $B(M, N; P)$ of all k -bilinear maps is naturally a k -module and we have an isomorphism

$$B(M, N; P) \simeq L(M, L(N, P)), \quad \varphi \mapsto (y \mapsto \varphi(x, y))$$

As an example, the action $k \times M \rightarrow M$ is bilinear and corresponds to a linear map $k \rightarrow L(M)$. A (*unitary associative*) *k-algebra* is a k -module A endowed with a k -bilinear map $A \times A \rightarrow A$ that turns A into a (unitary associative) ring. For example, $L(M) := L(M, M)$ is a k -algebra (for composition) and we shall write $\text{GL}(M) := L(M)^\times$. One sets $M_n(k) := L(k^n)$ and $\text{GL}_n(k) := \text{GL}(k^n)$.

If A is a ring, then the *tensor product* of a right A -module M with a left A -module N is the quotient $M \otimes_A N$ of the free abelian group $\mathbb{Z} \cdot (M \times N)$ by the following relation:

1. $\forall m \in M, n, n' \in N, \quad (m, n + n') \sim (m, n) + (m, n'),$
2. $\forall m, m' \in M, n \in N, \quad (m + m', n) \sim (m, n) + (m', n),$
3. $\forall a \in A, m \in M, n \in N, \quad (ma, n) \sim (m, an).$

We shall denote by $m \otimes n$ the class of (m, n) in $M \otimes_A N$. It should be noticed that

$$(M/M') \otimes_A N \simeq (M \otimes_A N)/(M' \otimes_A N)$$

if $M' \subset M$ is a submodule and that

$$(\oplus_{i \in I} M_i) \otimes_A N \simeq \oplus_{i \in I} (M_i \otimes_A N)$$

if $(M_i)_{i \in I}$ is a family of right A -modules. Of course we have the same results on the other side.

If k is a commutative ring and M, N are two k -modules, then $M \otimes_k N$ is endowed with the structure of a k -module through $a(m \otimes n) = am \otimes n$. Composition with the map

$$M \times N \mapsto M \otimes_k N, \quad (m, n) \mapsto m \otimes n$$

provides an isomorphism

$$L(M \otimes_k N, P) \simeq B(M, N; P)$$

for all k -module P . Note that we always have

$$M \otimes_k N \simeq N \otimes_k M, \quad (M \otimes_k N) \otimes_k P \simeq M \otimes_k (N \otimes_k P) \quad \text{and} \quad k \otimes_k M \simeq M.$$

Also,

$$(k \cdot E) \otimes_k (k \cdot F) \simeq k \cdot (E \times F)$$

when E, F are two sets.

If $f : A \rightarrow B$ is a ring homomorphism, then B becomes a right A -module via $ba = bf(a)$ and, if M is a left A -module, then $B \otimes_A M$ becomes a B -module via $b(b' \otimes s) = bb' \otimes s$. If A and B are two k -algebras, then $A \otimes_k B$ is endowed with a structure of a k -algebra through $(a \otimes b)(a' \otimes b') = aa' \otimes bb'$.

If k is a commutative algebra, then the n th *tensor power* of a k -module M is

$$T^n(M) := \underbrace{M \otimes_k \cdots \otimes_k M}_{n \text{ times}}.$$

If M is a k -module, then there exists a bilinear map $T^i(M) \times T^j(M) \rightarrow T^{i+j}(M)$ and the composite map

$$\left(\bigoplus_{i=0}^{\infty} T^i(M) \right) \times \left(\bigoplus_{j=0}^{\infty} T^j(M) \right) = \bigoplus_{i,j=0}^{\infty} T^i(M) \times T^j(M) \rightarrow \bigoplus_{n=0}^{\infty} T^n(M)$$

endows the *tensor algebra* $T(M) := \bigoplus_{n=0}^{\infty} T^n(M)$ with the structure of a ring. The *symmetric* (resp. *exterior*) algebra of M is the quotient $S(M)$ (resp. $\Lambda(M)$) $:= T(M)/I$ where I is the ideal generated by $x \otimes y - y \otimes x$ for $x, y \in M$ (resp. $x \otimes x$ for $x \in M$). The module of n th *symmetric* (resp. *exterior*) powers of M is the image $S^n(M)$ (resp. $\Lambda^n(M)$) of $T^n(M)$ in $S(M)$ (resp. $\Lambda(M)$). We have $S(M) = \bigoplus_{n=0}^{\infty} S^n(M)$ (resp. $\Lambda(M) = \bigoplus_{n=0}^{\infty} \Lambda^n(M)$). Any k -linear map $u : M \rightarrow N$ produces in an obvious way a k -linear map $T^n(u) : T^n(M) \rightarrow T^n(N)$ (resp. $S^n(u) : S^n(M) \rightarrow S^n(N)$, resp. $\Lambda^n(u) : \Lambda^n(M) \rightarrow \Lambda^n(N)$).

If M is free of rank n , then $\Lambda^n(M)$ is free of rank 1 and the “canonical” map $k \rightarrow L(\Lambda^n(M))$ is therefore an isomorphism. The *determinant* is the composite map

$$\det : L(M) \xrightarrow{\Lambda^n} L(\Lambda^n(M)) \simeq k.$$

On the other hand, there exists a general bilinear map

$$L(M, N) \times P \rightarrow L(M, N \otimes_k P), \quad (f, p) \mapsto (m \mapsto f(m) \otimes p)$$

providing a k -linear map $L(M, N) \otimes_k P \rightarrow L(M, N \otimes_k P)$. This is an isomorphism when P is free of finite rank. In particular, if M is a free module of finite rank, then there exists an isomorphism $M^\vee \otimes_k M \simeq L(M)$. Now, there exists a bilinear map

$$M^\vee \times M \rightarrow k, \quad (\varphi, m) \mapsto \langle \varphi, m \rangle := \varphi(m)$$

providing the *trace map*

$$\text{tr} : L(M) \simeq M^\vee \otimes_k M \rightarrow k$$

when M is free of finite rank.

If G is a monoid, then the composite map

$$k \cdot G \times k \cdot G \rightarrow k \cdot G \otimes_k k \cdot G \simeq k \cdot (G \times G) \rightarrow k \cdot G$$

turns the free k -module $k \cdot G$ into a k -algebra and this is functorial. The *polynomial ring* over k on a set E is $k[E] := k \cdot G$ where $G = \mathbb{N} \cdot E$ (free abelian monoid). Actually, $k[E] \simeq S(M)$ with $M = k \cdot E$.

A *lattice* is a set X endowed with two internal operations \vee and \wedge which are associative, commutative and *absorbant* with respect to each other:

$$\forall x, y \in X, \quad x \vee (x \wedge y) = x = x \wedge (x \vee y).$$

It is said to be *distributive* if

$$\forall x, y, z \in X, \quad x \vee (y \wedge z) = (x \vee y) \wedge (x \vee z) \quad \text{and} \quad x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z).$$

It is said to be *bounded* if both \vee and \wedge are unitary (with units denoted 0 and 1 respectively). Then, a *complement* for an $x \in X$ is an $\neg x \in X$ such that $x \vee \neg x = 1$ and $x \wedge \neg x = 0$. A *morphism of (bounded) lattices* is a homomorphism for both operations. The relation $x \leq y \Leftrightarrow x = x \wedge y$ is a partial order on X . Conversely, a partial order such that any pair has a supremum and an infimum defines a lattice by $x \vee y = \sup(x, y)$ and $x \wedge y = \inf(x, y)$. A morphism of lattices is an order preserving map that also preserves maxima and minima for pairs.

A *boolean ring* is a ring A such that all $a \in A$ are *idempotent*: $a^2 = a$. If we set $a \wedge b = ab$ and $a \vee b = a + b + ab$ and $\neg a = 1 + a$, then A becomes a bounded distributive lattice with complements. Conversely, if A is a bounded distributive lattice with complements, then it is a boolean ring.

0.5 Topology

A *topology* on a set X is a subset $\text{Open}(X) \subset \mathcal{P}(X)$ of the set of all subsets of X which is stable under any union (and in particular the empty union) and any finite intersection (and in particular empty intersection). A set X endowed with a topology is called a *topological space* and a subset $U \in \text{Open}(X)$ is said to be *open*. If there exists an open subset U such that $A \subset U \subset V \subset X$, then V is called a *neighborhood* of A . An $U \subset X$ is open if and only if it is a neighborhood of all its points.

Topologies are stable under intersection inside $\mathcal{P}(X)$. In particular, we may always consider the topology generated by any set of subsets of X . A topology is *finer* (resp. *coarser*) than another if it contains (resp. is contained in) the other: it has more (resp. less) open subsets. There exists a topology on X which is finer (resp. coarser) than any other: the *discrete* (resp. *coarse*) topology.

The complement of an open subset is said to be *closed*. There always exists a smallest closed (resp. biggest open) subset \overline{Y} (resp. $\overset{\circ}{Y}$) containing (resp. contained in) a given subset Y which is called its *closure* (resp. its *interior*). A subset Y is *dense* in X if $\overline{Y} = X$.

If $A \subset X$, then a *limit* for a map (remember that this is the same thing as a family) $f : A \rightarrow Y$ at $x \in \overline{A}$ is a $y \in Y$ such that, for all neighbourhood V of y in Y , there exists a neighborhood U of x in X such that $f(A \cap U) \subset V$. We then write $y = \lim_x f$ (but we should write \in because the limit may not be unique). A map $f : X \rightarrow Y$ is said to be *continuous* at $x \in X$ if $f(x)$ is a limit of f at x .

A map $f : X \rightarrow Y$ between two topological spaces is continuous (everywhere) if and only if

$$f^{-1}(\text{Open}(Y)) \subset \text{Open}(X).$$

We shall denote by $\mathcal{C}(X, Y)$ the set of all continuous maps $X \rightarrow Y$. On the other hand, a map f is said to be *open* (resp. *closed*) if it sends open (resp. closed) subsets to open (resp. closed) subsets. A *subspace* of a topological space X is a subset Y endowed with the *induced* topology: $\text{Open}(Y) := i^{-1}(\text{Open}(X))$ where $i : Y \hookrightarrow X$ denotes the inclusion map. If Y is a quotient of a topological space and $p : X \twoheadrightarrow Y$ denotes the quotient map, then the *quotient* topology on Y is defined by

$$V \in \text{Open}(Y) \Leftrightarrow p^{-1}(V) \in \text{Open}(X).$$

Unless otherwise specified, we will always implicitly endow a subset (resp. a quotient) with the induced (resp. quotient) topology. A topological space has *locally* a property P if any neighborhood V of a point x contains a neighborhood V' of x having property P . We shall encounter the notion of a *local homeomorphism* $f : X' \rightarrow X$. It means that there exists for all $x' \in X'$, an open neighborhood U' of x' (resp. U of $x := f(x')$) such that f induces a homeomorphism $U' \simeq U$.

A topological space X is said to be

- *Fréchet* if all points are closed,
- *Hausdorff* if any two distinct points have disjoint neighborhoods,
- *normal* if any two disjoint closed subsets have disjoint neighborhoods.

One also say T_1 for Fréchet, T_2 for Hausdorff and T_4 for normal Hausdorff (or equivalently normal Fréchet).

A *covering* of a topological space X is simply a set (or a family) E of subsets of X such that $X = \bigcup_{Y \in E} Y$. A *refinement* of E is a subset $E' \subset E$. A topological space X is said to be *compact* (resp. *paracompact*) if any open covering R of X has a finite (resp. locally finite) refinement $R' \subset R$ which is also a covering of X . A paracompact Hausdorff space is normal (and therefore T_4). Actually, if $X = \bigcup_{i \in I} U_i$ is a locally finite covering of a paracompact Hausdorff space, then there exists another covering $X = \bigcup_{j \in J} V_j$ such that $\forall j \in J, \exists i \in I, \overline{V_j} \subset U_i$.

A subset of a topological space X is said to be *clopen* if it is both open and closed. The space X is said to be *connected* if only \emptyset and X are clopen in X . The *connected components* in X are the maximal non-empty connected subspaces. The set of connected components of X is denoted by $\pi_0(X)$. The image of a connected space by a continuous map is always connected. The closure of a connected subset is connected. A connected component is closed (but not necessarily open). Connected components are disjoint. A clopen subset is a union of connected components.

A *semidistance* on a set X is a map $d : X \times X \rightarrow \mathbb{R}$ satisfying $d(x, x) = 0$, $d(x, y) = d(y, x)$ and $d(x, z) \leq d(x, y) + d(y, z)$ for $x, y, z \in E$. It is a *distance* if moreover $d(x, y) = 0 \Rightarrow x = y$. A set X endowed with a (semi) distance is called a (semi) *metric space*. A subset U of X is said to be *open* if

$$\forall x \in U, \exists \epsilon > 0, \forall y \in X, d(x, y) \leq \epsilon \Rightarrow y \in U.$$

This turns X into a topological space and d is a distance if and only if X is Hausdorff.

A *seminorm* on a real vector space E is a map

$$E \rightarrow \mathbb{R}, \quad x \mapsto \|x\|$$

satisfying

$$\|0\| = 0, \quad \|\lambda x\| = |\lambda| \|x\| \quad \text{and} \quad \|x + y\| \leq \|x\| + \|y\|$$

for $x, y \in E$ and $\lambda \in \mathbb{R}$. It is a *norm* if moreover $\|x\| = 0 \Rightarrow x = 0$. Formula $d(x, y) := \|y - x\|$ defines a (semi) distance on E . In the case of finite dimensional vector spaces, all norms are equivalent and all linear maps are continuous. Finally, if $f : A \rightarrow B$ is a map between subsets of finite dimensional vector spaces, we write $f(h) = o(h)$ (we should write \in instead of $=$) if there exists $\epsilon : A \rightarrow \mathbb{R}$ such that $\|f(h)\| = \|h\| \epsilon(h)$ for $h \in A$ and $\lim_0 \epsilon = 0$.

0.6 Some algebraic topology

A *homotopy* is a continuous map $s : X \times [0, 1] \rightarrow Y$. One then sets $s_t(x) := s(x, t)$ for $x \in X$ and $t \in [0, 1]$, says that s is a homotopy between s_0 and s_1 and writes $s : s_0 \sim s_1$. A continuous map $f : X \rightarrow Y$ is called a *homotopy equivalence* if there exists a continuous map $g : Y \rightarrow X$ such that $\text{Id}_X \sim g \circ f$ and $f \circ g \sim \text{Id}_Y$. A topological space said to be *contractible* if it is homotopy equivalent to a point.

For $n \in \mathbb{N}$, the *standard n -simplex*⁵ is the set

$$[n] := n + 1 := \{0, \dots, n\}$$

endowed with its natural order: $0 \leq 1 \leq \dots \leq n$. For $0 \leq i \leq n$, the (order preserving) map

$$\delta_n^i : [n - 1] \rightarrow [n] \quad (\text{resp. } \sigma_n^i : [n + 1] \rightarrow [n])$$

that forgets (resp. repeats) the i -th term is called the i th *face* (resp. *degeneracy*) map. Any order preserving map $u : [n] \rightarrow [m]$ is the composition of various degeneracy (surjective) and face (injective) maps. Moreover, we have the following relations

$$\begin{aligned} \delta_j^{n+1} \circ \delta_i^n &= \delta_i^{n+1} \circ \delta_{j-1}^n & \text{for } 0 \leq i < j \leq n + 1, \\ \sigma_j^n \circ \sigma_i^{n+1} &= \sigma_i^n \circ \sigma_{j+1}^{n+1} & \text{for } 0 \leq i \leq j \leq n, \\ \sigma_j^n \circ \delta_i^{n+1} &= \begin{cases} \delta_i^n \circ \sigma_{j-1}^{n-1} & \text{for } 0 \leq i < j \leq n \\ \text{Id}_{[n]} & \text{for } 0 \leq j \leq i \leq j + 1 \leq n + 1 \\ \delta_{i-1}^n \circ \sigma_j^{n-1} & \text{for } 0 \leq j < i - 1 \leq n. \end{cases} \end{aligned}$$

We now turn to the topological version. For $n \in \mathbb{N}$, the *standard topological n -simplex* is the subset:

$$|\Delta^n| := \left\{ (x_0, \dots, x_n) \in \mathbb{R}_{\geq 0}^{n+1}, \sum_{i=0}^n x_i = 1 \right\}.$$

⁵Sometimes also written Δ^n .

For $0 \leq i \leq n$, the map *ith face* (resp. *degeneracy*) map is (same notation as above)

$$\delta_n^i : |\Delta^{n-1}| \rightarrow |\Delta^n|, (x_0, \dots, x_{n-1}) \mapsto (x_0, \dots, x_{i-1}, 0, x_{i+1}, \dots, x_{n-1}) \quad (\text{resp.}$$

$$\sigma_n^i : |\Delta^{n+1}| \rightarrow |\Delta^n|, (x_0, \dots, x_{n+1}) \mapsto (x_0, \dots, x_{i-1}, x_i + x_{i+1}, x_{i+2}, \dots, x_{n+1})).$$

All above formulas still hold. More generally, any order preserving map $u : [m] \rightarrow [n]$ yields an analogous map (written in the same way) $u : |\Delta^m| \rightarrow |\Delta^n|$.

A *semisimplicial set*

$$S : \cdots S_2 \rightrightarrows S_1 \rightrightarrows S_0$$

is a family of maps $d_n^i : S_n \rightarrow S_{n-1}$ for $n > 0$ satisfying

$$d_n^i \circ d_{n+1}^j \circ = d_n^{j-1} \circ d_{n+1}^i \quad \text{for } 0 \leq i < j \leq n+1.$$

The *augmentation* map is the unique map $\epsilon : S_0 \rightarrow \{1\}$. A *simplicial set*

$$S : \cdots S_2 \rightrightarrows S_1 \rightrightarrows S_0$$

is the data of two families of maps $d_n^i : S_n \rightarrow S_{n-1}$ for $n > 0$ and $s_n^i : S_n \rightarrow S_{n+1}$ for $n \geq 0$ with $0 \leq i \leq n$ satisfying

$$d_n^i \circ d_{n+1}^j \circ = d_n^{j-1} \circ d_{n+1}^i \quad \text{for } 0 \leq i < j \leq n+1,$$

$$s_{n+1}^i \circ s_n^j = s_{n+1}^{j+1} \circ s_n^i \quad \text{for } 0 \leq i \leq j \leq n,$$

$$d_{n+1}^i \circ s_n^j = \begin{cases} s_{n-1}^{j-1} \circ d_n^i & \text{for } 0 \leq i < j \leq n \\ \text{Id}_{S_n} & \text{for } 0 \leq j \leq i \leq j+1 \leq n+1 \\ s_{n-1}^j \circ d_n^{i-1} & \text{for } 0 \leq j < i-1 \leq n. \end{cases}$$

If T is another (semi) simplicial set, then a *map of (semi) simplicial sets* $f : S \rightarrow T$ is a family of compatible maps $f_n : S_n \rightarrow T_n$.

When S is a (semi) simplicial set, we let $C_n(S) := \mathbb{Z} \cdot S_n$ and, when M is an abelian group, $C^n(S, M) := \mathcal{F}(S_n, M)$. The map d_n^i extends uniquely to a group homomorphism (still written) $d_n^i : C_n(S) \rightarrow C_{n-1}(S)$. On the other hand, the map d_{n+1}^i induces by composition a group homomorphism $d_n^i : C^n(S, M) \rightarrow C^{n+1}(S, M)$. In both cases, we shall be concerned with the alternating sum

$$d_n := \sum_{i=0}^{n-1} (-1)^i d_n^i.$$

Note that the augmentation map provides a homomorphism $e : C_0(S) \rightarrow \mathbb{Z}$ (sending S_0 to 1) and a homomorphism $e : M \rightarrow C^0(S)$ (sending m to the constant map m).

If X is a topological space, there exists a simplicial set $S_\bullet(X)$ with

$$S_n(X) := \mathcal{C}(|\Delta^n|, X), \quad d_i^n(\sigma) := \sigma \circ \delta_n^i, \quad \text{et} \quad s_i^n(\sigma) := \sigma \circ \sigma_n^i$$

and we shall write $C_n(X) := C_n(S_\bullet(X))$ as well as $C^n(X, M) := C^n(S_\bullet(X), M)$.

0.7 Some analytic geometry

Let $U \subset E$ and $V \subset F$ be open subsets of finite dimensional real vector spaces. A map $f : U \rightarrow V$ is said to be *differentiable* at $x \in U$ if there exists $f'(x) \in L(E, F)$ such that

$$f(x+h) - f(x) - f'(x)(h) = o(h).$$

When f is everywhere differentiable, this defines a map $f' : U \rightarrow L(E, F)$. One can then define $f^{(k)}$ and the notion of a k -differentiable map by induction on $k \in \mathbb{N}$. The map f is said to be \mathcal{C}^k -differentiable (resp. \mathcal{C}^∞ -differentiable or *smooth*, resp. \mathcal{C}^{an} -differentiable or *analytic*) if moreover $f^{(k)}$ is continuous (resp. if f is k -differentiable for all $k \in \mathbb{N}$, resp. and if moreover

$$f(x+h) = \sum_{k=0}^{\infty} \frac{1}{k!} f^{(k)}(x)(h, \dots, h)$$

in a neighborhood of any $x \in U$ – seeing $f^{(k)}(x)$ as a multilinear map). There exists an obvious complex analog but any complex differentiable map is then automatically \mathcal{C}^{an} and then said to be *holomorphic*. Everything below has an equivalent in this situation.

Let us fix some $k \in \{0, 1, \dots, \infty, \text{an}\}$.

A *chart* on a set X is a bijection $\varphi : U \xrightarrow{\sim} U'$ where $U \subset X$ and U' is an open subset of \mathbb{R}^n for some $n \in \mathbb{N}$. The components x_1, \dots, x_n of the composite map $U \xrightarrow{\sim} U' \hookrightarrow \mathbb{R}^n$ are called *local coordinates* on X . A \mathcal{C}^k -*atlas* on X is a family of charts $\varphi_i : U_i \xrightarrow{\sim} U'_i$ for $i \in I$ such that $X = \bigcup_{i \in I} U_i$ and for all $i, j \in I$, if we let $U_{ij} := U_i \cap U_j$, then $U'_{ij} := \varphi_i(U_{ij})$ is open and the map

$$\varphi_{j|U_{ji}} \circ \varphi_{i|U'_{ij}}^{-1} : U'_{ij} \simeq U'_{ji}$$

is \mathcal{C}^k -differentiable. Then, there exists a unique topology on X making each U_i open and φ_i a homeomorphism. Any \mathcal{C}^k -atlas is contained in a unique maximal \mathcal{C}^k -atlas and a \mathcal{C}^k -differentiable manifold is a set X endowed with a maximal \mathcal{C}^k -atlas. One also makes the extra topological assumption that X is Hausdorff and *countable at infinity* (a countable union of compact subsets). A map $f : X \rightarrow Y$ between \mathcal{C}^k -differentiable manifolds is said to be \mathcal{C}^k -differentiable if for all charts $\varphi : U \xrightarrow{\sim} U'$ and $\psi : V \xrightarrow{\sim} V'$ of X and Y respectively, if we set $W := U \cap f^{-1}(V)$, then $W' = \varphi(W)$ is open and the map

$$\psi \circ f|_W \circ \varphi_{|W'}^{-1} : W' \rightarrow V'$$

is \mathcal{C}^k -differentiable. It is called a \mathcal{C}^k -*diffeomorphism* if it is bijective and f^{-1} also is \mathcal{C}^k -differentiable. We shall denote by $\mathcal{C}^k(X, Y)$ the set of all \mathcal{C}^k -differentiable maps from X to Y and by $\mathcal{O}(X)$ the ring of \mathcal{C}^k -differentiable maps $f : X \rightarrow \mathbb{R}$. Note that, if $f \in \mathcal{O}(X)$ does not vanish, then $1/f \in \mathcal{O}(X)$.

Again, there exists a complex analog of the above theory leading to the notion of *complex manifold* (independent of $k \neq 0$ in which case we write \mathcal{C}^{hol}).

A *vector bundle* on a \mathcal{C}^k -differentiable manifold X is a \mathcal{C}^k -differentiable map $p : E \rightarrow X$ endowed with the structure of a vector space on each fiber $E_x := p^{-1}(x)$

such that there exists a covering of X by open subsets and for each such open subset U , a \mathcal{C}^k -diffeomorphism $E|_U := p^{-1}(U) \simeq U \times \mathbb{R}^n$ inducing an isomorphism of vector spaces $E_x \simeq \{x\} \times \mathbb{R}^n$ for all $x \in U$. As a set, E is simply the disjoint union of the vector spaces E_x for $x \in X$ (but they vary smoothly). A first example is given by the *trivial* vector bundle $X \times V$ if V is a finite dimensional vector space. It may be convenient to write $\mathcal{O}_X := X \times \mathbb{R}$.

A *morphism of vector bundles* is a \mathcal{C}^k -differentiable map $u : E \rightarrow F$ inducing a linear map $u_x : E_x \rightarrow F_x$ for all $x \in X$. We shall denote by $\text{Hom}(E, F)$ the set of all morphisms of vector bundles from E to F . One defines a subbundle as well as an isomorphism of vector bundles in the usual way. One can generalize the definitions of $E \oplus F$, $L(E, F)$ (and in particular $L(E)$ and E^\vee) and $E \otimes F$ from vector spaces to vector bundles so that they specialize to fibers (for example $(E \oplus F)_x = E_x \oplus F_x$). One can also extend the notions of tensor, symmetric and exterior powers and define $T^n(E)$, $S^n(E)$ and $\Lambda^n(E)$ when E is a vector bundle. A *section* of a vector bundle E on an open subset U of X is a \mathcal{C}^k -differentiable map $s : U \rightarrow E$ such that the composite $p \circ s$ is merely the inclusion $U \hookrightarrow X$. The set $E(U)$ (or $\Gamma(U, E)$) of all sections of E on U is an $\mathcal{O}(U)$ -module (note that $\mathcal{O}_X(U) \simeq \mathcal{O}(U)$ showing that our notations are compatible). Finally, if E, F are two vector bundles on X then there exists a bijection

$$\text{Hom}(E, F) \simeq \Gamma(X, L(E, F)), \quad u \mapsto (x \mapsto u_x \in L(E_x, F_x) = L(E, F)_x).$$

We assume from now on that $k \neq 0$ (and henceforth say *topological manifold* in the case $k = 0$).

A *curve* on a \mathcal{C}^k -manifold X is a \mathcal{C}^k -differentiable map $\gamma : I \rightarrow X$ defined on an open neighborhood I of $0 \in \mathbb{R}$. Two curves γ_1, γ_2 are said to be *tangent* if there exists a chart $\varphi : U \xrightarrow{\sim} U'$ of X such that⁶

$$\gamma_1(0) = \gamma_2(0) \in U \quad \text{and} \quad (\varphi \circ \gamma_1)'(0) = (\varphi \circ \gamma_2)'(0).$$

This is an equivalence relation on the set of all curves γ on X and the set TX of equivalence classes $[\gamma]$ is the *tangent bundle*: there exists an obvious map

$$p : TX \rightarrow X, \quad [\gamma] \mapsto \gamma(0)$$

and if $\varphi : U \xrightarrow{\sim} U'$ is a chart of X , a bijection

$$TX|_U \simeq U \times \mathbb{R}^n, \quad [\gamma] \mapsto (\gamma(0), (\varphi \circ \gamma)'(0)).$$

Any morphism $f : X \rightarrow Y$ provides a morphism of vector bundles

$$Tf : TX \rightarrow TY, \quad [\gamma] \mapsto [f \circ \gamma]$$

(on charts, it corresponds to $(x, h) \mapsto (x, f'(x)(h))$).

We assume from now on that $k = \infty$ or $k = \text{an}$.

A *differential form* of degree p on X is a (differentiable) section of $\Lambda^p(TX)^\vee$. They form an $\mathcal{O}(X)$ -module $\Omega^p(X)$. The maps $\Lambda^p(TX)^\vee \times \Lambda^q(TX)^\vee \rightarrow \Lambda^{p+q}(TX)^\vee$ coming from the ring structure of exterior powers provide maps

$$\Omega^p(X) \times \Omega^q(X) \rightarrow \Omega^{p+q}(X), \quad (\omega, \eta) \mapsto \omega \wedge \eta.$$

⁶After replacing γ_i by its restriction to $\gamma_i^{-1}(U)$.

There exists an \mathbb{R} -linear map

$$d : \mathcal{O}(X) = \mathcal{C}^k(X, \mathbb{R}) \xrightarrow{T} \text{Hom}(TX, T\mathbb{R}) \simeq \Gamma(X, (TX)^\vee) = \Omega^1(X).$$

It extends uniquely to maps $d : \Omega^p(X) \rightarrow \Omega^{p+1}(X)$ for all $p = 1, \dots, n-1$ such that

$$d(\omega \wedge \eta) = d(\omega) \wedge \eta + (-1)^p \omega \wedge d\eta \quad (1)$$

for $\omega \in \Omega^p(X)$ and $\eta \in \Omega^q(X)$. The sequence of maps

$$\mathcal{O}(X) \xrightarrow{d} \Omega^1(X) \xrightarrow{d} \dots \xrightarrow{d} \Omega^n(X)$$

is called the *de Rham complex* of X .

0.8 Some algebraic geometry

All rings are assumed to be commutative.

If \mathfrak{a} is an ideal in a ring A , then its *radical* is

$$\sqrt{\mathfrak{a}} = \{f \in A, \exists n \in \mathbb{N}, f^n \in \mathfrak{a}\}.$$

An $f \in \sqrt{(0)}$ is then said to be *nilpotent*. The ring A is said to be *reduced* if 0 is the only nilpotent element. It is called a *domain* (resp. a *field*) if $A \setminus \{0\}$ is a multiplicative sub-monoid (resp. group). If \mathfrak{a} is an ideal in a ring A and M is an A -module, then A/\mathfrak{a} becomes a ring and $M/\mathfrak{a}M$ an A/\mathfrak{a} module. The ideal \mathfrak{a} is said to be *radical* (resp. *prime*, resp. *maximal*) if A/\mathfrak{a} is reduced (resp. a domain, resp. a field). The ring A is called a *local ring* if there exists a unique maximal ideal \mathfrak{m} in which case $k := A/\mathfrak{m}$ is called the *residue field* of A .

If $S \subset A$ is a multiplicative submonoid and M is an A -module, then the *localization* of M at S is

$$S^{-1}M = (M \times S) / \sim \quad \text{with} \quad (a, s) \sim (b, t) \Leftrightarrow \exists u \in S, u(ta - sb) = 0.$$

The set $S^{-1}A$ is actually a ring and $S^{-1}M$ an $S^{-1}A$ -module. If S is generated by G , then there exists isomorphisms

$$S^{-1}A \simeq A[\{T_s\}_{s \in G}] / (sT_s - 1)_{s \in G} \quad \text{and} \quad S^{-1}M \simeq S^{-1}A \otimes_A M.$$

As a special case, if $f \in A$ (resp. $\mathfrak{p} \subset A$ is a prime ideal), then the *localization* of M at f (resp. \mathfrak{p}) is M_f (resp. $M_{\mathfrak{p}} := S^{-1}M$ where S is the submonoid generated by f (resp. $S = A \setminus \mathfrak{p}$)). If A is a domain, then $\text{Frac}(A) := A_{(0)}$ is called the *fraction field* of A . In general, when \mathfrak{p} is a prime ideal of A , $A_{\mathfrak{p}}$ is a local ring with maximal ideal $\mathfrak{p}A_{\mathfrak{p}}$.

The *(prime) spectrum* of A is the set $X := \text{Spec}(A)$ of prime ideals of A . Unless otherwise specified, everything below also holds when $X = \text{Spm}(A)$ is the set of maximal ideals (the *maximal spectrum*). If $\mathfrak{p} \in X$, we let $\kappa(\mathfrak{p}) := \text{Frac}(A/\mathfrak{p})$. If $f \in A$, we shall denote by $f(\mathfrak{p})$ the image of f in $\kappa(\mathfrak{p})$:

$$\begin{aligned} A &\longrightarrow A/\mathfrak{p} \twoheadrightarrow \kappa(\mathfrak{p}) \\ f &\longmapsto \bar{f} \longmapsto f(\mathfrak{p}) \end{aligned}$$

so that $f(\mathfrak{p}) = 0 \Leftrightarrow f \in \mathfrak{p}$. If $S \subset A$, we may consider the *zero locus*

$$V(S) := \{\mathfrak{p} \in X \mid \forall f \in S, f(\mathfrak{p}) = 0\}$$

of S , so that $\mathfrak{p} \in V(S) \Leftrightarrow S \subset \mathfrak{p}$. We always have

$$V(\emptyset) = X, \quad V(A) = \emptyset, \quad V(ST) = V(S) \cup V(T) \quad \text{and} \quad V\left(\bigcup_{i \in I} S_i\right) = \bigcap_{i \in I} V(S_i).$$

We define the *Zariski topology* on X by requiring a set to be closed if and only if it is the zero locus of some $S \subset A$.

On the other hand, if $E \subset X$, then we set

$$I(E) := \{f \in A \mid \forall \mathfrak{p} \in E, f(\mathfrak{p}) = 0\}.$$

The maps $S \mapsto V(S)$ et $E \mapsto I(E)$ induce inverse bijections between radical ideals of A and closed subsets of X . More precisely,

$$\forall S \subset A, \quad I(V(S)) = \sqrt{(S)} \quad \text{and} \quad \forall E \subset X, \quad V(I(E)) = \overline{E}.$$

A basis for the Zariski topology of X is given by the *special domains*

$$D(f) := \{\mathfrak{p} \in X \mid f(\mathfrak{p}) \neq 0\}$$

for $f \in A$ and we have $D(f) \cap D(g) = D(fg)$. One easily checks that X is compact (but not Hausdorff in general).

If $\varphi : A \rightarrow B$ is a ring homomorphism and $Y := \text{Spec}(B)$, then there exists⁷ a continuous map $u := \varphi^{-1} : Y \rightarrow X$. If \mathfrak{a} is an ideal of A (resp. $f \in A$), then the “canonical” map $A \rightarrow A/\mathfrak{a}$ (resp. $A \rightarrow A[1/f]$) provides a homeomorphism

$$\text{Spec}(A/\mathfrak{a}) \simeq V(\mathfrak{a}) \quad (\text{resp. } \text{Spec}(A[1/f]) \simeq D(f)).$$

If $R \rightarrow A$ is a morphism of commutative rings and $I = \ker(A \otimes_R A \rightarrow A)$, then $\Omega_{A/R} := I/I^2$ is the *module of relative differential forms* of A/R . The R -linear map $A \rightarrow A \otimes A, f \mapsto 1 \otimes f - f \otimes 1$ induces an R -linear map $d : A \rightarrow \Omega_{A/R}$. If we set $\Omega_{A/R}^p := \Lambda^p \Omega_{A/R}$, then d extends through formula (1) in order to give the *de Rham complex* $\Omega_{A/R}^\bullet$ of A/R .

⁷This is not true in general for maximal ideals.

1. Categories and functors

For those who might worry about set-theoretic issues (see [Shu08] for example), we shall stay in a fixed *universe* (some large set, see definition 1.1.1 of [KS06] for example). We shall then call *set* only those sets that belong (\in) to our universe¹ and rename *collection* (or call it *large*) a set that is only contained (\subset) in our universe.

1.1 Category

1.1.1 Definition/Examples

Definition 1.1.1 A *category*^a consists in the following data:

1. a collection \mathcal{C} of *objects*,
2. for all $X, Y \in \mathcal{C}$, a set $\text{Hom}(X, Y)$ of *morphisms*,
3. for all $X \in \mathcal{C}$, an *identity* morphism $\text{Id}_X \in \text{End}(X) := \text{Hom}(X, X)$,
4. for all $X, Y, Z \in \mathcal{C}$, a composition rule

$$\text{Hom}(X, Y) \times \text{Hom}(Y, Z) \rightarrow \text{Hom}(X, Z), \quad (f, g) \mapsto g \circ f$$

such that

- (a) $\text{Id}_Y \circ f = f = f \circ \text{Id}_X$,
- (b) if $h \in \text{Hom}(Z, T)$, then $h \circ (g \circ f) = (h \circ g) \circ f$.

^aSome authors allow collections (and not merely sets) of morphisms in the definition of a category and call *locally small category* what we named a category.

The category is said to be *small* if all its objects (or equivalently all its morphisms) form a set and not merely a collection and *finite* if there exists only a finite number of morphisms (and consequently of objects). It is called *discrete* if the only morphisms are the identities. It is *empty* if there is no objects at all (and no morphisms).

¹They are usually called *small sets* but we do not want to keep this epithet everywhere.

We will usually write $f : X \rightarrow Y$ instead of $f \in \text{Hom}(X, Y)$ and call X (resp. Y) the *domain* (resp. le *codomain*) of f . Note that Id_X is uniquely determined by the conditions (4a). A morphism from X to itself is called an *endomorphism*. In practice, we shall simply say that the collection \mathcal{C} is a category² but we must not forget that it involves some extra structure: morphisms and composition. One may write $\text{Ob}(\mathcal{C})$ (resp. $\text{Mor}(\mathcal{C})$) for the collection of all objects (resp. all morphisms) of \mathcal{C} and $\text{Hom}_{\mathcal{C}}(X, Y)$ (with a subscript \mathcal{C}) for the set of morphisms from X to Y .

- Examples**
1. If G is a monoid, one can then consider the category \mathbf{G} with a unique object \bullet , $\text{End}(\bullet) := G$ and composition given by multiplication on G^{op} ($g \circ f = fg$). This way, we get essentially all categories with a unique object.
 2. If \leq is a preorder on a set X , we will then denote by \mathbf{X} the category whose objects are the elements $x \in X$ and morphisms are couples (x, y) for $x \leq y$ (composition is then uniquely defined). This way, we get essentially all the small categories whose Hom have at most one element.
 3. As a particular case, one will always endow an unordered set with the minimal (also called trivial) preorder “=” and consider any set as a category (small category with no morphisms besides identities). This way, we get all small discrete categories.
 4. We shall consider the category \mathbf{n} associated to the (trivial) set $n := \{0, 1, \dots, n-1\}$ (with only identities as morphisms). But we shall also consider the category $[\mathbf{n}]$ associated to the set $[n] := \{0, \dots, n\}$ with respect to usual order $0 \leq 1 \leq \dots \leq n$. For example, $[\mathbf{0}] = \mathbf{1}$ is the category that has exactly one object and one morphism, $[\mathbf{1}]$ is a category with two distinct objects and a unique morphism between them (plus the identities).
 5. If X is a topological space, then $\text{Open}(X)$ is ordered by inclusion and we shall denote by $\mathbf{Open}(X)$ the corresponding category.
 6. There exists a (small) category Δ with positive integers $[n] := n+1 = \{0, \dots, n\}$ as objects and order preserving maps as morphisms.
 7. We shall denote by \mathbf{Set} the (large) category whose objects are sets and morphisms are maps between them. We shall write $\mathcal{F}(X, Y)$ for the set of all maps between two sets.
 8. In the same way, we shall consider the (large) category \mathbf{Top} whose objects are topological spaces and morphisms are continuous maps. We shall write $\mathcal{C}(X, Y)$ for the set of continuous maps between two topological spaces.
 9. Finally, we will denote by \mathbf{Ab} the (large) category whose objects are abelian groups and morphisms are homomorphisms.

Exercise 1.1.2 Define the categories \mathbf{Ord} , \mathbf{Mon} , \mathbf{Grp} , \mathbf{Ring} , \mathbf{CRing} , $G\text{-Set}$ (resp. $\mathbf{Set}\text{-}G$), $A\text{-Mod}$ (resp. $\mathbf{Mod}\text{-}A$) and $k\text{-Alg}$ of preordered sets, monoids, groups, rings, commutative rings, left (resp. right) G -sets, left (resp. right) A -modules and k -algebras.

The *opposite category* to a category \mathcal{C} is the category \mathcal{C}^{op} with the same objects as \mathcal{C} but $\text{Hom}_{\mathcal{C}^{\text{op}}}(X, Y) = \text{Hom}_{\mathcal{C}}(Y, X)$ (and composition going in the reverse direction).

²As one usually denotes a group by G without explicitly mentioning the multiplication rule.

We have $(\mathcal{C}^{\text{op}})^{\text{op}} = \mathcal{C}$. The *product* $\mathcal{C} \times \mathcal{C}'$ of two categories \mathcal{C} and \mathcal{C}' is made into a category in the obvious way: everything is done termwise. This extends to arbitrary products $\prod_{i \in I} \mathcal{C}_i$ (families of objects and families of morphisms). If \mathcal{C} is a category, then the category $\mathbf{Mor}(\mathcal{C})$ of *morphisms of \mathcal{C}* is defined as follows: an object is a morphism of \mathcal{C} and a morphism between $f : X \rightarrow Y$ and $g : X' \rightarrow Y'$ is a pair of morphisms in \mathcal{C} , $\varphi : X \rightarrow X'$ and $\psi : Y \rightarrow Y'$, such that $g \circ \varphi = \psi \circ f$:

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ \downarrow \varphi & & \downarrow \psi \\ X' & \xrightarrow{f'} & Y'. \end{array}$$

Also, if $X \in \mathcal{C}$, then the category ${}_X\mathcal{C}$ of *objects of \mathcal{C} under X* is defined as follows: an object of ${}_X\mathcal{C}$ is a morphism $f : X \rightarrow Y$ and a morphism from $f : X \rightarrow Y$ to $g : X \rightarrow Z$ is a morphism $h : Y \rightarrow Z$ such that $h \circ f = g$:

$$\begin{array}{ccc} & & Y \\ & \nearrow f & \downarrow h \\ X & & \\ & \searrow g & \downarrow \\ & & Z. \end{array}$$

Also, $\mathcal{C}_{/X} := ({}_X\mathcal{C}^{\text{op}})^{\text{op}}$ is the category of *objects of \mathcal{C} over X* . This is our first example of *dual* construction.

Exercise 1.1.3 Show that, if G is a monoid, then the set of objects of $\mathbf{Mor}(\mathbf{G})$ is G and that, if X is a preordered set, then the set of objects of $\mathbf{Mor}(\mathbf{X})$ is the graph Γ of the relation.

Exercise 1.1.4 Make $\mathcal{C}_{/X}$ explicit when X is an object of a category \mathcal{C} .

1.1.2 Isomorphism

Definition 1.1.5 In a category \mathcal{C} ,

1. a *section* (resp. a *retraction*) of a morphism $f : X \rightarrow Y$ is a morphism $g : Y \rightarrow X$ such that $f \circ g = \text{Id}_Y$ (resp. $\text{Id}_X = g \circ f$):

$$X \xrightleftharpoons[f]{g} Y \quad (\text{resp. } \hookrightarrow X \xrightleftharpoons[f]{g} Y).$$

2. an *isomorphism* is a morphism that has at the same time a section and a retraction. When there exists an isomorphism $X \xrightarrow{\sim} Y$, one says that X and Y are *isomorphic* and writes $X \simeq Y$. An isomorphism between X and itself is called an *automorphism*.

One denotes by $\text{Isom}(X, Y)$ the set of all isomorphisms $X \xrightarrow{\sim} Y$ and by $\text{Aut}(X)$ the group of all automorphisms of X .

Sections and retractions may also be called *right* and *left inverses* respectively. Clearly, f is a section of g if and only if g is a retraction of f . Also retraction (resp.

section) in \mathcal{C} is the same thing as a section (resp. retraction) in \mathcal{C}^{op} (they are *dual* notions).

Examples

1. A map $f : X \rightarrow Y$ (between sets) has a section (resp. retraction) if and only if it is surjective (resp. injective unless $X = \emptyset$).
2. A *retract* A of a topological space X is a subspace such that the inclusion map $A \hookrightarrow X$ has a retraction.
3. A *direct summand* M' of a module M is a submodule such that the inclusion map $M' \hookrightarrow M$ has a retraction.

Proposition 1.1.6 If f is an isomorphism, then it has a unique section and a unique retraction and they are the same.

Proof. If $\text{Id}_Y = f \circ g$ and $h \circ f = \text{Id}_X$, then

$$h = h \circ \text{Id}_Y = h \circ f \circ g = \text{Id}_X \circ g = g. \quad \blacksquare$$

The unique section/retraction of an isomorphism f is called its *inverse* and denoted by f^{-1} .

Exercise 1.1.7 What is an isomorphism in **Set**, in **Top**, in **Ab**, etc. ? In **G** if G is a monoid ? In **X** if X is a preordered set ?

1.1.3 Subcategory

Definition 1.1.8 A *subcategory* of a category \mathcal{C} is the data of

1. a subcollection $\mathcal{C}' \subset \mathcal{C}$,
2. for all $X, Y \in \mathcal{C}'$, a subset $\text{Hom}_{\mathcal{C}'}(X, Y) \subset \text{Hom}_{\mathcal{C}}(X, Y)$, such that
 - (a) if $X \in \mathcal{C}'$, then $\text{Id}_X \in \text{End}_{\mathcal{C}'}(X) := \text{Hom}_{\mathcal{C}'}(X, X)$,
 - (b) if $X, Y, Z \in \mathcal{C}'$, $f \in \text{Hom}_{\mathcal{C}'}(X, Y)$ and $g \in \text{Hom}_{\mathcal{C}'}(Y, Z)$, then $g \circ f \in \text{Hom}_{\mathcal{C}'}(X, Z)$.

It is a *full* subcategory if actually

$$\forall X, Y \in \mathcal{C}', \quad \text{Hom}_{\mathcal{C}'}(X, Y) = \text{Hom}_{\mathcal{C}}(X, Y).$$

A subcategory becomes a category with the induced composition. A full subcategory is uniquely determined by its objects. The subcategory is said to be *wide* if $\text{Ob}(\mathcal{C}') = \text{Ob}(\mathcal{C})$.

Examples

1. **Ab** is a full subcategory of **Grp** which itself is a full subcategory of **Mon** (which itself is a *non-full* subcategory of the category of semigroups or magmas for example).
2. If X is a topological space, then **Open**(X) is a subcategory of **Set** which is *not* full.
3. **Top** is *not* a subcategory of **Set** and neither is **Ab** (but see the notion of a faithful functor below).
4. If X is any topological space, then an *espace étalé* over X is a local homeomorphism $X' \rightarrow X$. They form a full subcategory $\mathbf{Et}(X) \subset \mathbf{Top}_{/X}$.
5. **Δ** is a full subcategory of **Ord**.

6. We may consider the wide subcategory Δ_{inj} of Δ with the same objects but only order preserving *injective* maps as morphism.
7. Conversely, we may consider Δ (resp. Δ_{inj}) as a full subcategory of Δ^+ (resp. Δ_{inj}^+) which is defined by adding $[-1] = \emptyset$.

If \mathcal{C} is a category, we can consider the relation “ \rightarrow ” on the collection \mathcal{C} given by

$$X \rightarrow Y \Leftrightarrow \text{Hom}(X, Y) \neq \emptyset$$

as well as the quotient³

$$\pi_0(\mathcal{C}) := \mathcal{C} / \rightarrow .$$

Definition 1.1.9 A *connected component* of \mathcal{C} is an element of $\pi_0(\mathcal{C})$. The category \mathcal{C} is said to be *connected* if there exists a unique connected component.

The relation “ \rightarrow ” is not an equivalence relation: actually, X and Y will be connected to each other (in the same connected component) when there exists a zigzag of morphisms

$$X = X_0 \rightarrow X_1 \leftarrow X_2 \rightarrow \cdots \leftarrow X_{n-1} \rightarrow X_n = Y.$$

The category \mathcal{C} is the disjoint union of its connected components that will all be viewed as full subcategories of \mathcal{C} .

- Examples**
1. The categories **Set**, **Ab**, **Top**, etc. are all connected.
 2. If a set X is endowed with a preorder relation R , then $\pi_0(\mathbf{X}) = X/R$ and a connected component is the same thing as an equivalence class.

1.2 Functor

1.2.1 Definition/Examples

- Definition 1.2.1**
1. A (*covariant*) *functor* $F : \mathcal{C} \rightarrow \mathcal{C}'$ between two categories is the data for all $X \in \mathcal{C}$ of $F(X) \in \mathcal{C}'$ and for all $f : X \rightarrow Y$ of $F(f) : F(X) \rightarrow F(Y)$, in such a way that we always have $F(\text{Id}_X) = \text{Id}_{F(X)}$ and $F(g \circ f) = F(g) \circ F(f)$.
 2. If $G : \mathcal{C}' \rightarrow \mathcal{C}''$ is another functor, then their *composite* is the functor $G \circ F$ given by $(G \circ F)(X) = G(F(X))$ and $(G \circ F)(f) = G(F(f))$.

We will denote by $\text{Hom}(\mathcal{C}, \mathcal{C}')$ the collection of all functors $\mathcal{C} \rightarrow \mathcal{C}'$. We will often describe the functors by their action on the objects and let the reader guess what happens for morphisms.

There always exists an identity functor $\text{Id}_{\mathcal{C}} : \mathcal{C} \rightarrow \mathcal{C}$ that doesn't change anything. Also, a category \mathcal{C}' is a subcategory of \mathcal{C} if and only if all objects (resp. morphisms) of \mathcal{C}' are objects (resp. morphisms) of \mathcal{C} and inclusion is functorial. A functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is called an *isomorphism* if there exists a functor G such that $\text{Id}_{\mathcal{C}} = G \circ F$ and $F \circ G = \text{Id}_{\mathcal{C}'}$ (but this is not a very interesting notion). A functor $F : \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}'$ is also called a *contravariant* functor from \mathcal{C} to \mathcal{C}' . Any functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ provides a functor $F^{\text{op}} : \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}'^{\text{op}}$ and this construction is “functorial” : $\text{Id}_{\mathcal{C}^{\text{op}}} = \text{Id}_{\mathcal{C}}^{\text{op}}$ and $(G \circ F)^{\text{op}} = G^{\text{op}} \circ F^{\text{op}}$.

³A very large set.

Examples 1. There exists a functor that forgets topology (and continuity) $\mathbf{Top} \rightarrow \mathbf{Set}$ (underlying set). In the other direction, there exists a functor $X \mapsto X^{\text{disc}}$ (resp. $X \mapsto X^{\text{coarse}}$) that endow a set X with the discrete (resp. coarse) topology.

2. There also exists a (contravariant) functor

$$\mathbf{Top}^{\text{op}} \rightarrow \mathbf{Ord}, \quad X \mapsto \text{Open}(X), \quad f \mapsto f^{-1}.$$

3. There exists a functor that forgets the algebraic structure $\mathbf{Ab} \rightarrow \mathbf{Set}$ (underlying set) and, in the other direction, a functor $X \mapsto \mathbb{Z} \cdot X$ (or $\mathbb{Z}^{(X)}$) which sends a set to the free *abelian group* generated by X .

4. There exists an inclusion functor $\mathbf{Grp} \hookrightarrow \mathbf{Mon}$ and two functors $G \mapsto G^{\times}$ and

$$G \mapsto G^{\text{gr}} := \langle x_g, g \in G \mid x_{gh} = x_g x_h, g, h \in G, x_1 = 1 \rangle$$

in the reverse direction.

5. Small categories and functors between them form a category that we shall denote \mathbf{Cat} and there exists functors

$$\mathbf{Mon} \rightarrow \mathbf{Cat}, \quad G \mapsto \mathbf{G} \quad \text{and} \quad \mathbf{Ord} \rightarrow \mathbf{Cat}, \quad X \mapsto \mathbf{X}.$$

6. A (contravariant) functor $F : \mathcal{C}^{\text{op}} \rightarrow \mathcal{D}$ is also called a *presheaf* on \mathcal{C} with values in \mathcal{D} . If X is a topological space, then a presheaf F on $\mathbf{Open}(X)$ with values in \mathbf{Set} (also called a *presheaf of sets on X*) is the following data:

- (a) a set $F(U)$ of *sections* on any open subset U of X , and
- (b) *restriction maps* $F(U) \rightarrow F(V), s \mapsto s|_V$ for $V \subset U$ satisfying $(s|_V)|_W = s|_W$ when $W \subset V$ and $s|_U = s$.

Exercise 1.2.2 What is a functor $\mathbf{G} \rightarrow \mathbf{H}$ between categories associated to monoids ? What is a functor $\mathbf{X} \rightarrow \mathbf{Y}$ between categories associated to preordered sets ?

Exercise 1.2.3 What are the analogs of the “free abelian group” functor $X \mapsto \mathbb{Z} \cdot X$ for the categories \mathbf{Mon} , \mathbf{Grp} , $A\text{-Mod}$ and $k\text{-Alg}$?

Hint. We only describe the first one. The free monoid generated by a set X is $G := \{g : n \rightarrow X, n \in \mathbb{N}\}$ endowed with

$$gh : n + m \rightarrow X, \quad i \mapsto \begin{cases} g(i) & \text{if } i < n \\ h(i - n) & \text{otherwise.} \end{cases}$$

In other words, X represents the alphabet, G denotes the set of words formed using this alphabet, and the operation is concatenation. ■

Exercise 1.2.4 Show that, besides the inclusion functor $\mathbf{Ab} \hookrightarrow \mathbf{Grp}$, there exists an *abelianization* functor $G \mapsto G^{\text{ab}} = G/[G, G]$ in the other direction. Show that the center is *not* functorial in the sense that a group homomorphism $\varphi : G \rightarrow H$ does not necessarily induce a morphism of abelian groups $Z(G) \rightarrow Z(H)$.

Exercise 1.2.5 Show that the categories $\mathbb{Z}\text{-Mod}$ and \mathbf{Ab} are isomorphic. Same thing with the categories $\mathbb{Z}\text{-Alg}$ and \mathbf{Ring} , and, more generally, $k\text{-Alg}$ and a full subcategory of $k\text{-Ring}$ (the image of k must be in the center).

Exercise 1.2.6 Show that the image of a section (resp. a retraction, resp. an inverse) by a functor is a section (resp. a retraction, resp. an inverse).

If we are given two categories \mathcal{C} and \mathcal{C}' , then the projections $\mathcal{C} \times \mathcal{C}'$ on \mathcal{C} and \mathcal{C}' are functorial. The same holds for the obvious partial functors $\mathcal{C}' \rightarrow \mathcal{C} \times \mathcal{C}'$ or $\mathcal{C} \rightarrow \mathcal{C} \times \mathcal{C}'$ associated to a fixed object $X \in \mathcal{C}$ or $X' \in \mathcal{C}'$. If \mathcal{C} is any category, then there exists domain and codomain functors $\mathbf{Mor}(\mathcal{C}) \rightrightarrows \mathcal{C}$ as well as forgetful functors $X \backslash \mathcal{C} \rightarrow \mathcal{C}$ and $\mathcal{C}/X \rightarrow \mathcal{C}$.

Exercise 1.2.7 Let us denote by $\mathbf{Op}(\mathcal{C}) \subset \mathbf{Mor}(\mathcal{C})$ the subcategory whose objects are morphisms with codomain identical to the domain and morphisms have same component on domain and codomain (objects with operator). Show that, if k is a commutative ring, then $\mathbf{Op}(k\text{-Mod})$ is isomorphic to $k[t]\text{-Mod}$.

If \mathcal{C} is any category, then there exists a very important functor

$$\mathrm{Hom} : \mathcal{C}^{\mathrm{op}} \times \mathcal{C} \rightarrow \mathbf{Set}, \quad (X, Y) \mapsto \mathrm{Hom}_{\mathcal{C}}(X, Y)$$

that sends a couple (f, g) of morphisms to the map $\varphi \mapsto g \circ \varphi \circ f$. If we compose with partial functors, we get the (fundamental) functors

$$h_{\mathcal{C}}^X : \mathcal{C} \rightarrow \mathbf{Set}, \quad Y \mapsto \mathrm{Hom}_{\mathcal{C}}(X, Y)$$

and

$$h_Y^{\mathcal{C}} := h_{\mathcal{C}^{\mathrm{op}}}^Y : \mathcal{C}^{\mathrm{op}} \rightarrow \mathbf{Set}, \quad X \mapsto \mathrm{Hom}_{\mathcal{C}}(X, Y).$$

When there is no ambiguity, we shall drop the reference to \mathcal{C} and simply write h^X and h_Y . We may also write $g_*(\varphi) := h^g(\varphi) = g \circ \varphi$ and $f^*(\varphi) = h_f(\varphi) = \varphi \circ f$.

1.2.2 Natural transformation

Definition 1.2.8 1. If $F, G : \mathcal{C} \rightarrow \mathcal{C}'$ are two functors, then a *natural transformation* $\alpha : F \Rightarrow G$ is a collection of morphisms $\alpha_X : F(X) \rightarrow G(X)$ for all $X \in \mathcal{C}$ such that, for all $f : X \rightarrow Y$, we have $\alpha_Y \circ F(f) = G(f) \circ \alpha_X$:

$$\begin{array}{ccc} F(X) & \xrightarrow{F(f)} & F(Y) \\ \downarrow \alpha_X & & \downarrow \alpha_Y \\ G(X) & \xrightarrow{G(f)} & G(Y). \end{array}$$

We shall say *natural isomorphism* and write $F \simeq G$ if all α_X are isomorphisms.

2. If $\beta : G \Rightarrow H$ is another natural transformation, then their *composition* $\beta \circ \alpha : F \Rightarrow H$ is the natural transformation defined by $(\beta \circ \alpha)_X = \beta_X \circ \alpha_X$ for $X \in \mathcal{C}$.

We shall denote by $\mathbf{Hom}(F, G)$ the collection of all natural transformations from F to G . If \mathcal{C} is a small category, then the functors $F : \mathcal{C} \rightarrow \mathcal{C}'$ make a (large) category $\mathbf{Hom}(\mathcal{C}, \mathcal{C}')$ with natural transformations as morphisms. One can check that the isomorphisms are then exactly the natural isomorphisms defined above.

Examples 1. With

$$\begin{array}{ccc} F : \mathbf{CRing} & \longrightarrow & \mathbf{Grp} \\ k & \longmapsto & \mathrm{GL}_n(k) \end{array} \quad \text{and} \quad \begin{array}{ccc} G : \mathbf{CRing} & \longrightarrow & \mathbf{Grp} \\ k & \longmapsto & k^\times, \end{array}$$

there exists a natural transformation $\alpha : F \rightarrow G$ by considering

$$\alpha_k = \det : \mathrm{GL}_n(k) \rightarrow k^\times.$$

This is a natural isomorphism for $n = 1$.

2. If \mathcal{C} is any category, then there exists isomorphisms of categories

$$\mathbf{Hom}(\mathbf{0}, \mathcal{C}) \simeq \mathbf{1}, \quad \mathbf{Hom}(\mathbf{1}, \mathcal{C}) \simeq \mathcal{C},$$

$$\mathbf{Hom}(\mathbf{2}, \mathcal{C}) \simeq \mathcal{C} \times \mathcal{C} \text{ and } \mathbf{Hom}([1], \mathcal{C}) \simeq \mathbf{Mor}(\mathcal{C}).$$

3. If G is a monoid, then there exists an isomorphism $\mathbf{Hom}(\mathbf{G}, \mathbf{Set}) \simeq G\text{-Set}$.
 4. Presheaves on a small category \mathcal{C} with values in a category \mathcal{D} form a category $\hat{\mathcal{C}}(\mathcal{D}) := \mathbf{Hom}(\mathcal{C}^{\mathrm{op}}, \mathcal{D})$ simply denoted $\hat{\mathcal{C}}$ in the case $\mathcal{D} = \mathbf{Set}$. In the case of presheaves of sets on a topological space X , a morphism $\alpha : F \rightarrow G$ is given by a family of maps $\alpha_U : F(U) \rightarrow G(U)$ for U open in X satisfying $\alpha_V(s|_V) = (\alpha_U(s))|_V$ whenever $V \subset U$ and $s \in \mathcal{F}(U)$.

Exercise 1.2.9 Show that

1. if \mathcal{C} is a small category, then there exists an isomorphism of categories $\mathbf{Hom}(\mathcal{C}, \mathcal{D})^{\mathrm{op}} \simeq \mathbf{Hom}(\mathcal{C}^{\mathrm{op}}, \mathcal{D}^{\mathrm{op}})$,
2. if \mathcal{C} and \mathcal{C}' are two small categories, then there exists an isomorphism of categories

$$\mathbf{Hom}(\mathcal{C} \times \mathcal{C}', \mathcal{D}) \simeq \mathbf{Hom}(\mathcal{C}, \mathbf{Hom}(\mathcal{C}', \mathcal{D})).$$

By composition, if \mathcal{C} is a small category, then any functor $G : \mathcal{D} \rightarrow \mathcal{D}'$ induces a functor

$$\mathbf{Hom}(\mathcal{C}, \mathcal{D}) \rightarrow \mathbf{Hom}(\mathcal{C}, \mathcal{D}'), \quad F \mapsto G \circ F.$$

Also, if \mathcal{D} is any category, then any functor $G : \mathcal{C} \rightarrow \mathcal{C}'$ between small categories induces a functor

$$\mathbf{Hom}(\mathcal{C}', \mathcal{D}) \rightarrow \mathbf{Hom}(\mathcal{C}, \mathcal{D}), \quad F \mapsto F \circ G.$$

Exercise 1.2.10 Show that, there exists a functor

$$\mathbf{Cat}^{\text{op}} \times \mathbf{Cat} \rightarrow \mathbf{Cat}, \quad (\mathcal{C}, \mathcal{D}) \mapsto \mathbf{Hom}(\mathcal{C}, \mathcal{D}).$$

Definition 1.2.11 Assume given two functors $S : \mathcal{S} \rightarrow \mathcal{C}$ and $T : \mathcal{T} \rightarrow \mathcal{C}$ with same target \mathcal{C} . An object in the *comma category* $(S \downarrow T)$ is a triple $(X \in \mathcal{S}, Y \in \mathcal{T}, f : S(X) \rightarrow T(Y))$. A morphism $(X, Y, f) \rightarrow (X', Y', f')$ is a pair of morphisms $u : X \rightarrow X', v : Y \rightarrow Y'$ such that $F(v) \circ f = f' \circ F(u)$.

- Examples**
1. In case $\mathcal{C} = \mathbf{1}$, we have $(S \downarrow T) \simeq \mathcal{S} \times \mathcal{T}$.
 2. We have $(S \downarrow T)^{\text{op}} \simeq (T^{\text{op}} \downarrow S^{\text{op}})$.
 3. $(\text{Id}_{\mathcal{C}} \downarrow \text{Id}_{\mathcal{C}}) \simeq \mathbf{Mor}(\mathcal{C}) \simeq \mathbf{Hom}([1], \mathcal{C})$.
 4. $(\text{Id}_{\mathcal{C}} \downarrow 1 \xrightarrow{X} \mathcal{C}) \simeq \mathcal{C}_{/X}$ and $(1 \xrightarrow{X} \mathcal{C} \downarrow \text{Id}_{\mathcal{C}}) \simeq {}_{X \backslash} \mathcal{C}$.
 5. One may also write $\mathcal{C}_{/T} := (\text{Id}_{\mathcal{C}} \downarrow T)$ and ${}_{S \backslash} \mathcal{C} := (S \downarrow \text{Id}_{\mathcal{C}})$.
 6. Or $S_{/X} := (S \downarrow 1 \xrightarrow{X} \mathcal{C})$ and ${}_{X \backslash} T := (1 \xrightarrow{X} \mathcal{C} \downarrow T)$

Recall that the notion of a simplicial set was introduced in section 0.6.

Exercise 1.2.12 Show that the category \mathbf{SSet} of simplicial sets is isomorphic to the category $\widehat{\Delta}$ of presheaves of sets on Δ .

Hint. This is based on the following observations:

1. Any morphism in Δ splits as a surjective morphism followed by an injective morphism,
2. any injective (resp. surjective) morphism $[n] \hookrightarrow [m+1]$ with $n \leq m$ factors through an injective (resp. surjective) morphism $[m] \rightarrow [m+1]$ ($[n] \rightarrow [n-1]$),
3. an injective (resp. surjective) morphism $[m] \rightarrow [m+1]$ (resp. $[n] \rightarrow [n-1]$) is the same thing as a face (resp. degeneracy) map. ■

It is straightforward to generalize the notion of a simplicial set:

Definition 1.2.13 A *simplicial object* of a category \mathcal{D} is a presheaf

$$X_{\bullet} : \Delta^{\text{op}} \rightarrow \mathcal{D}, \quad [n] \mapsto X_n.$$

Dually, a *cosimplicial object* is a functor $X^{\bullet} : \Delta \rightarrow \mathcal{D}, [n] \mapsto X^n$.

Exercise 1.2.14 Show that there exists a cosimplicial object $|\Delta^{\bullet}|$ of \mathbf{Top} sending $[n]$ to $|\Delta^n|$ and $u : [n] \rightarrow [m]$ to the unique linear map sending e_i to $e_{u(i)}$ if (e_0, \dots, e_n) denotes the usual basis of \mathbb{R}^{n+1} .

Exercise 1.2.15 Show that the simplicial set $S_{\bullet}(X)$ associated to a topological space X corresponds to $h_X \circ |\Delta^{\bullet}|$ under $\mathbf{SSet} \simeq \widehat{\Delta}$ and that this provides a functor $\mathbf{Top} \rightarrow \widehat{\Delta}$ sending a topological space X to its associated simplicial set $S_{\bullet}(X)$.

If \mathcal{C} is a small category, then the *nerve* of \mathcal{C} is the simplicial set \mathcal{C}_{\bullet} where \mathcal{C}_n is the set of couples

$$((X_i)_{i=0}^n, (f_i \in \text{Hom}(X_{i-1}, X_i)_{i=1}^n))$$

depicted as $X_0 \xrightarrow{f_1} X_1 \xrightarrow{f_2} \dots \xrightarrow{f_n} X_n$ and such a couple is sent by $u = [m] \rightarrow [n]$ to

$$((X_{u(i)})_{i=0}^m, (f_{u(i)} \circ f_{u(i)-1} \circ \dots \circ f_{u(i-1)} \in \text{Hom}(X_{u(i)-1}, X_{u(i)})_{i=1}^m))$$

(with the convention that the empty composition is the identity).

Exercise 1.2.16 Show that \mathcal{C}_\bullet is indeed a simplicial set and that we obtain a functor $\mathbf{Cat} \rightarrow \widehat{\Delta}$. Give an explicit description in low degree.

Proof. We only do the explicit description: $\mathcal{C}_0 = \mathcal{C}$ is the set of objects X , \mathcal{C}_1 is the set of morphisms $X \rightarrow Y$, \mathcal{C}_2 is the set of pairs of morphisms $X \rightarrow Y \rightarrow Z$. One has

$$s^0(X) = \text{Id}_X, \quad \begin{cases} d^0(X \xrightarrow{f} Y) = Y \\ d^1(X \xrightarrow{f} Y) = X, \end{cases} \quad \begin{cases} s^0(X \xrightarrow{f} Y) = X \xrightarrow{\text{Id}_X} X \xrightarrow{f} Y \\ s^1(X \xrightarrow{f} Y) = X \xrightarrow{f} Y \xrightarrow{\text{Id}_Y} Y, \end{cases}$$

$$\begin{cases} d^0(X \xrightarrow{f} Y \xrightarrow{g} Z) = Y \xrightarrow{g} Z \\ d^1(X \xrightarrow{f} Y \xrightarrow{g} Z) = X \xrightarrow{g \circ f} Z \\ d^2(X \xrightarrow{f} Y \xrightarrow{g} Z) = X \xrightarrow{f} Y, \end{cases} \quad \begin{cases} s^0(X \xrightarrow{f} Y \xrightarrow{g} Z) = X \xrightarrow{\text{Id}_X} X \xrightarrow{f} Y \xrightarrow{g} Z \\ s^1(X \xrightarrow{f} Y \xrightarrow{g} Z) = X \xrightarrow{f} Y \xrightarrow{\text{Id}_Y} Y \xrightarrow{g} Z \\ s^2(X \xrightarrow{f} Y \xrightarrow{g} Z) = X \xrightarrow{f} Y \xrightarrow{g} Z \xrightarrow{\text{Id}_Z} Z. \end{cases}$$

■

1.2.3 Equivalence

Definition 1.2.17 A functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is said to be

1. *faithful* (resp. *full*, resp. *fully faithful*) if for all $X, Y \in \mathcal{C}$, the map

$$\text{Hom}(X, Y) \rightarrow \text{Hom}(F(X), F(Y)), \quad f \mapsto F(f)$$

is injective (resp. surjective, resp. bijective).

2. *essentially surjective* if for all $X' \in \mathcal{C}'$, there exists $X \in \mathcal{C}$ such that $X' \simeq F(X)$.
3. an *equivalence of categories* if there exists $G : \mathcal{C}' \rightarrow \mathcal{C}$ such that $\text{Id}_{\mathcal{C}} \simeq G \circ F$ and $F \circ G \simeq \text{Id}_{\mathcal{C}'}$ (and G is then called a *quasi-inverse*).

One can also define the *essential image* of a functor F as the collection of all $X' \in \mathcal{C}'$ such that there exists $X \in \mathcal{C}$ with $X' \simeq F(X)$. The functor F is then essentially surjective when the essential image is equal to \mathcal{C}' .

The inclusion of a (full) subcategory is a (fully) faithful functor that we shall call an *embedding*. An isomorphism of categories is an equivalence (but not conversely). We shall use the notation $\mathcal{C} \simeq \mathcal{C}'$ for the wider notion of *equivalence* (and not merely isomorphism) of categories. One sometimes says that two categories \mathcal{C} and \mathcal{C}' are *anti-equivalent* if $\mathcal{C}^{\text{op}} \simeq \mathcal{C}'$.

Exercise 1.2.18 Show that the forgetful functors $\mathbf{Top} \rightarrow \mathbf{Set}$ and $\mathbf{Ab} \rightarrow \mathbf{Set}$ are faithful but not fully faithful.

Exercise 1.2.19 Show that there exists a fully faithful functor

$$\mathbf{Mon} \hookrightarrow \mathbf{Cat}, \quad G \mapsto \mathbf{G}, \quad (\text{resp. } \mathbf{Ord} \hookrightarrow \mathbf{Cat}, \quad X \mapsto \mathbf{X}).$$

What is the essential image ?

Exercise 1.2.20 Show that there exists a fully faithful functor $\mathbf{Cat} \hookrightarrow \widehat{\Delta}$.

Exercise 1.2.21 Show that if $F \simeq F'$, then F is faithful (resp. full, resp. fully faithful, essentially surjective, an equivalence) if and only if F' is.

Theorem 1.2.22 A functor is an equivalence of categories if and only if it is fully faithful and essentially surjective.

Proof. To show that the condition is necessary, we first remark that our functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ will be essentially surjective since we will always have $X' \simeq F(X)$ with $X := G(X')$ if G is a quasi-inverse for F . Then, we consider the following sequence of maps

$$\begin{aligned} \text{Hom}(X, Y) &\xrightarrow{F} \text{Hom}(F(X), F(Y)) \\ &\xrightarrow{G} \text{Hom}(G(F(X)), G(F(Y))) \xrightarrow{F} \text{Hom}(F(G(F(X))), F(G(F(Y)))) \end{aligned}$$

Since $G \circ F$ (resp. $F \circ G$) is fully faithful, the composition of the first (resp. last) two arrows is bijective. It follows that all arrows are actually bijections and, in particular, F is fully faithful.

To show that the condition is sufficient, we *choose* for all $X' \in \mathcal{C}'$ an object $X \in \mathcal{C}$ and an isomorphism $\beta_{X'} : F(X) \simeq X'$. We set $G(X') := X$ so that $\beta_{X'} : F(G(X')) \simeq X'$. Since F is fully faithful, there exists for each $f' : X' \rightarrow Y'$ in \mathcal{C}' a unique $f : G(X') \rightarrow G(Y')$ such that $F(f) = \beta_{Y'}^{-1} \circ f' \circ \beta_{X'}$. We then set $G(f') = f$ so that $\beta_{Y'} \circ F(G(f')) = f' \circ \beta_{X'}$. One easily checks that G is a functor and we obtain by construction a natural isomorphism $\beta : F \circ G \simeq \text{Id}_{\mathcal{C}'}$. In particular, if $X \in \mathcal{C}$, then there exists a natural isomorphism $\beta_{F(X)} : F(G(F(X))) \simeq F(X)$ and, since F is fully faithful, there exists a unique morphism $\alpha_X : X \simeq (G \circ F)(X)$ such that $F(\alpha_X) = \beta_{F(X)}^{-1}$. One easily checks that α is indeed a natural isomorphism. ■

Exercise 1.2.23 Show that if X is a preordered set and Y denotes its ordered quotient, then the categories \mathbf{X} and \mathbf{Y} are equivalent.

Exercise 1.2.24 Show that, if A is a ring, then $\mathbf{Mat}(A) := \mathbf{N}$, endowed with $\text{Hom}(m, n) = M_{n \times m}(A)$ and multiplication of matrices, is a small category. Show that if A is a field k , then $\mathbf{Mat}(k)$ is equivalent, but not isomorphic, to the category of finite dimensional k -vector spaces (which is large).

1.2.4 Universal property

Definition 1.2.25 Given a functor $F : \mathcal{C} \rightarrow \mathbf{Set}$, an object $X \in \mathcal{C}$ and $s \in F(X)$, we say that (X, s) is *universal* for F or that it *represents* F if

$$\forall Y \in \mathcal{C}, \forall t \in F(Y), \exists! f : X \rightarrow Y, \quad F(f)(s) = t. \quad (1.1)$$

This may be pictured as follows:

$$\begin{array}{ccc} X & \xrightarrow{\quad} & F(X) \\ \exists! f \downarrow & & \downarrow F(f) \\ Y & \xrightarrow{\quad} & F(Y) \end{array} \quad \begin{array}{c} s \\ \downarrow \\ t. \end{array}$$

We may also say that “the couple $(X \in \mathcal{C}, s \in F(X))$ is universal among all couples $(Y \in \mathcal{C}, t \in F(Y))$ ”, or that “ $s \in F(X)$ is universal for $t \in F(Y)$ ” or that “ X represents F ” (not mentioning s explicitly).

- Examples**
1. $(\mathbb{Z}, 1)$ represents the forgetful functor $\mathbf{Grp} \rightarrow \mathbf{Set}$.
 2. If k is a commutative ring then $(k[t], t)$ represents the forgetful functor $k\text{-}\mathbf{Alg} \rightarrow \mathbf{Set}$ (universal property of polynomials).
 3. $(\mathbb{Z}, 0)$ represents the constant functor $\mathbf{Ring} \rightarrow \mathbf{1} \rightarrow \mathbf{Set}$ (\mathbb{Z} is initial for rings).
 4. The real line \mathbb{R} (together with $\text{Id}_{\mathbb{R}}$) represents the (contravariant) functor

$$\mathcal{O} : X \mapsto \mathcal{C}^k(X, \mathbb{R})$$

on \mathcal{C}^k -manifolds.

5. The inclusion $Y \hookrightarrow X$ of a subspace into a topological space is universal for continuous maps $f : Z \rightarrow X$ such that $f(Z) \subset Y$ (contravariant F).

Exercise 1.2.26 Show that, if F is represented by both X and X' , then $X \simeq X'$. More precisely, show that if both (X, s) and (X', s') are universal for F , then there exists a unique isomorphism $f : X \simeq X'$ such that $F(f)(s) = s'$.

Exercise 1.2.27 Show that usual forgetful functors are representable.

Exercise 1.2.28 Show that $M \otimes_k N$ is universal for (the functor that sends P to the set of) bilinear maps $M \times N \rightarrow P$.

Exercise 1.2.29 Let k be a commutative ring and $f_1, \dots, f_r \in k[t_1, \dots, t_n]$. Show that the functor that sends a commutative k -algebra A (make the morphisms explicit) to the set

$$\mathcal{S}(A) := \{(a_1, \dots, a_n) \in A^n \mid f_1(a_1, \dots, a_n) = \dots = f_r(a_1, \dots, a_n) = 0\}$$

of all solutions with values in A , is representable.

Lemma 1.2.30 — Yoneda. If $F : \mathcal{C} \rightarrow \mathbf{Set}$ is any functor and $X \in \mathcal{C}$, then there exists a natural bijection

$$\text{Hom}(h^X, F) \simeq F(X), \quad \alpha \mapsto \alpha_X(\text{Id}_X).$$

With a contravariant functor $F : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$, it simply reads $\text{Hom}(h_X, F) \simeq F(X)$.

Proof. Given $s \in F(X)$, if $Y \in \mathcal{C}$ and $f : X \rightarrow Y$, then we set $\alpha_Y(f) := F(f)(s)$. This defines a map $\alpha_Y : \text{Hom}(X, Y) \rightarrow F(Y)$ and we shall show that this is natural, meaning that

$$\alpha_Z \circ h^X(g) = F(g) \circ \alpha_Y$$

if $g : Y \rightarrow Z$. Indeed, we do have

$$\begin{aligned} (\alpha_Z \circ h^X(g))(f) &= \alpha_Z(h^X(g)(f)) = \alpha_Z(g \circ f) = F(g \circ f)(s) \\ &= (F(g) \circ F(f))(s) = (F(g)(F(f)(s))) = F(g)(\alpha_Y(f)) = (F(g) \circ \alpha_Y)(f). \end{aligned}$$

It only remains to check that we did define a natural inverse as well as the naturality (exercise). ■

Proposition 1.2.31 A functor $F : \mathcal{C} \rightarrow \mathbf{Set}$ is represented by $X \in \mathcal{C}$ if and only if $h^X \simeq F$.

Proof. We may simply apply Yoneda lemma: the condition means that there exists a natural transformation $\alpha : h^X \rightarrow F$ which is an isomorphism. This α corresponds to some $s \in F(X)$ and property 1.1 exactly means that α_Y is always bijective since, necessarily, $F(f)(s) = \alpha_Y(f)$. ■

In other words, F is represented by X if and only if there exists a natural bijection $\text{Hom}(X, Y) \simeq F(Y)$. For a contravariant functor $F : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$, the condition reads $\text{Hom}(Y, X) \simeq F(Y)$ or, equivalently, $h_X \simeq F$.

Exercise 1.2.32 Show that if \mathcal{C} is a small category, then there exists a fully faithful Yoneda functor

$$\mathfrak{Y} : \mathcal{C} \rightarrow \widehat{\mathcal{C}}, \quad X \mapsto h_X.$$

Solution. If $X, Y \in \mathcal{C}$, then Yoneda's lemma on \mathcal{C}^{op} shows that the map

$$\text{Hom}(h_X, h_Y) \simeq h_Y(X) = \text{Hom}(X, Y), \quad \alpha \mapsto \alpha_X(\text{Id}_X)$$

is bijective. It is therefore sufficient to notice that, for $f : X \rightarrow Y$, we have $h_f^X(\text{Id}_X) = f \circ \text{Id}_X = f$. ■

1.3 Limit

1.3.1 Diagram, cone and limit

Definition 1.3.1 Let I be a small category and \mathcal{C} any category. A *commutative diagram* of shape I in \mathcal{C} is a functor $D : I \rightarrow \mathcal{C}$.

A commutative diagram on I in \mathcal{C} is therefore the data of an object X_i for all $i \in I$ and a morphism $f_\alpha : X_i \rightarrow X_j$ for all $\alpha : i \rightarrow j$ satisfying $f_{\text{Id}_i} = \text{Id}_{X_i}$ and $f_{\beta \circ \alpha} = f_\beta \circ f_\alpha$:

$$\begin{array}{ccc} X_i & \xrightarrow{f_\alpha} & X_j \\ & \searrow f_{\beta \circ \alpha} & \swarrow f_\beta \\ & X_k & \end{array}$$

We shall denote such a diagram as $(f_\alpha : X_i \rightarrow X_j)$ or (X_i, f_α) , and the category of all commutative diagrams on I in \mathcal{C} by $\mathcal{C}^I := \mathbf{Hom}(I, \mathcal{C})$. A morphism $(X_i, f_\alpha) \rightarrow (Y_i, g_\alpha)$ is simply a family of morphisms $u_i : X_i \rightarrow Y_i$ satisfying $g_\alpha \circ u_i = u_j \circ f_\alpha$ for all $\alpha : i \rightarrow j$.

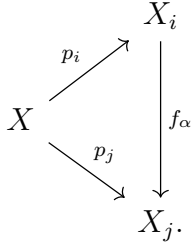
A commutative diagram is said to be *empty*, *discrete* or *finite* when I has this property. In particular, a discrete diagram is the same thing as a family of objects.

By composition, any functor $\lambda : I \rightarrow J$ between small categories will provide a functor $\lambda^* : \mathcal{C}^J \rightarrow \mathcal{C}^I$ between diagrams (and this is functorial). As a particular case, the unique functor $I \rightarrow \mathbf{1}$ induces the *constant diagram* functor

$$\mathcal{C} \simeq \mathcal{C}^1 \rightarrow \mathcal{C}^I, \quad X \mapsto \underline{X}.$$

Definition 1.3.2 A *cone* for a diagram D in \mathcal{C} is an object $X \in \mathcal{C}$ together with a morphism $\underline{X} \rightarrow D$.

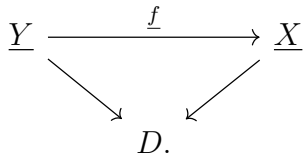
In more down to earth terms, a cone for (X_i, f_α) is an object X together with a family of *structural morphisms* $p_i : X \rightarrow X_i$ satisfying $p_j = f_\alpha \circ p_i$ whenever $\alpha : i \rightarrow j$:



The cone is said to be *empty*, *discrete* or *finite* when the diagram has this property. The dual notion is that of a *cocone*.

Definition 1.3.3 A *limit* X of a commutative diagram D of shape I in \mathcal{C} is a universal cone for D .

In other words, X is a limit for D if and only if X is a cone for D and, given any cone Y for D , there exists a unique morphism $f : Y \rightarrow X$ making commutative the diagram



Formally, it means that the composite (contravariant) functor $h_D \circ _$ is represented by X . Equivalently, there exists a natural isomorphism

$$\mathrm{Hom}_{\mathcal{C}^I}(\underline{Y}, D) \simeq \mathrm{Hom}_{\mathcal{C}}(Y, X)$$

in $Y \in \mathcal{C}$.

In down to earth terms, a commutative diagram (X_i, f_α) has X as a limit if and only if we are given for each $i \in I$ a morphism $p_i : X \rightarrow X_i$ such that for all $\alpha : i \rightarrow j$, we have $p_j = f_\alpha \circ p_i$ with the following universal property: if we are given some $Y \in \mathcal{C}$ endowed for all $i \in I$ with a morphism $g_i : Y \rightarrow X_i$ such that for each $u : i \rightarrow j$, we have $g_j = f_\alpha \circ g_i$, then there exists a unique morphism $g : Y \rightarrow X$ such

that for all $i \in I$, we have $g_i = p_i \circ g$:

$$\begin{array}{ccc}
 & g_i & \rightarrow X_i \\
 & \nearrow p_i & \downarrow f_\alpha \\
 Y & \overset{g}{\dashrightarrow} X & \\
 & \searrow p_j & \downarrow \\
 & g_j & \rightarrow X_j
 \end{array} \tag{1.2}$$

A limit is unique up to a unique isomorphism and we may sometimes say *the* limit and write $X = \varprojlim D$ or $X = \varprojlim_I D$. This should however been understood in the sense that X (together with all structural maps) is *a* limit of D . A limit X of a diagram D in \mathcal{C}^{op} is also called a *colimit* in \mathcal{C} and we shall write $X = \varinjlim D$. Some authors call a limit (resp. colimit) an *inverse* (resp. a *direct*) limit, a *projective* (resp. an *inductive*) limit or a *left* (resp. a *right*) limit. We have the following important formulas:

$$\text{Hom}(\underline{Y}, D) \simeq \text{Hom}(Y, \varprojlim D) \quad \text{et} \quad \text{Hom}(\varinjlim D, Y) \simeq \text{Hom}(D, \underline{Y}).$$

When we write a limit or a colimit, we implicitly assume that it exists.

Exercise 1.3.4 Write down a diagram for the colimit as in (1.2).

Exercise 1.3.5 Show that in an preordered set, a limit (resp. colimit) is a least upper bound or *inf* or *join* (resp. greatest lower bound or *sup* or *meet*). What about cone and cocone ?

Exercise 1.3.6 Show that, if $D \rightarrow D'$ is a morphism of diagrams of shape I in \mathcal{C} with respective limits X and X' , then there exists a unique morphism $X \rightarrow X'$ making commutative the diagram

$$\begin{array}{ccc}
 D & \longrightarrow & D' \\
 \uparrow & & \uparrow \\
 X & \longrightarrow & X'
 \end{array}$$

Definition 1.3.7 A functor $\lambda : J \rightarrow I$ between small categories is said to be *final*^a if the comma category $i \backslash \lambda$ is non empty connected whenever $i \in I$.

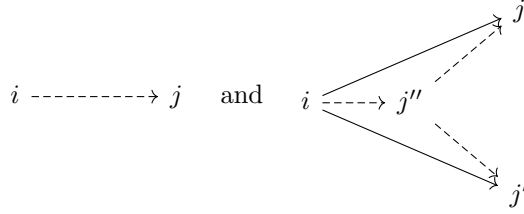
^aThis used to be called cofinal.

For the sake of completeness, recall that an object of $i \backslash \lambda$ is a couple (j, α) with $j \in J$ and $\alpha : i \rightarrow \lambda(j)$. A morphism $(j, \alpha) \rightarrow (j', \alpha')$ is a morphism $\beta : j \rightarrow j'$ in J such that $\lambda(\beta) \circ \alpha = \alpha'$.

We may also say that J is final in I via λ . The dual notion is that of an *initial* functor (λ^{op} is final). An order preserving map $\lambda : I \rightarrow J$ is said to be final/initial if the corresponding functor is

- Examples**
1. $I \rightarrow 1$ is final if and only if I is non empty connected.
 2. A functor $1 \rightarrow \mathcal{C}, 0 \mapsto X$ is final if and only if X is a final object (see below) of \mathcal{C} .
 3. A subset of \mathbb{N} is final (for the usual order) if and only if it is infinite.

4. Assume $J \subset I$ and we can complete (on the right) any of the following diagrams:



(it means that

$$\begin{aligned} &\forall i \in I, \exists j \in J, \exists \alpha : i \rightarrow j \text{ and} \\ &\forall i \in I, \forall j, j' \in J, \forall \alpha : i \rightarrow j, \alpha' : i \rightarrow j', \\ &\exists j'' \in J, \exists \gamma : i \rightarrow j'', \beta' : j'' \rightarrow j, \beta' : j'' \rightarrow j', \beta \circ \gamma = \alpha, \beta' \circ \gamma = \alpha'). \end{aligned}$$

Then J is final in I .

Exercise 1.3.8 Show that, if X is a metric space and \mathcal{V} denotes the set of neighborhoods of $x \in X$, then the map

$$\mathbb{R}_{>0} \rightarrow \mathcal{V}, \quad \epsilon \mapsto \mathbb{B}(x, \epsilon)$$

is initial^a.

^aThis maybe seen as the origin of the surge of ϵ in analysis.

We may always replace the shape of a diagram along a final (or initial) functor:

Proposition 1.3.9 The following are equivalent (and dual):

1. the functor $\lambda : J \rightarrow I$ is final,
2. if D is a diagram of shape I , then $\varinjlim \lambda^*(D) = \varinjlim D$.

Proof. To do. ■

1.3.2 Specific limits

Definition 1.3.10 A empty limit in \mathcal{C} is called a *final* object and denoted by $1_{\mathcal{C}}$. Dually, we get the notion of an *initial* object $0_{\mathcal{C}}$.

If $X \in \mathcal{C}$, there exists a unique morphism $X \rightarrow 1_{\mathcal{C}}$ (resp. $0_{\mathcal{C}} \rightarrow X$).

Examples

1. In **Set**, the initial object is \emptyset and 1 is a final object (defined up to a unique bijection).
2. Same thing in **Top**.
3. In **Ab**, $\{0\}$ is both a final and an initial object.

In general, cones for a diagram D in \mathcal{C} form a category $\mathcal{C}_{/D}$ and a limit of D is the same thing as a final object of $\mathcal{C}_{/D}$.

Definition 1.3.11 A discrete limit is called a *product* and denoted by $\prod_{i \in I} X_i$. The dual notion is that of a *coproduct* denoted by $\coprod_{i \in I} X_i$.

If $X = \prod_{i \in I} X_i$, then there exists structural maps $p_i : X \rightarrow X_i$. It is equivalent to give a morphism $f : Y \rightarrow X$ or its *components* $f_i := p_i \circ f$ and we shall then write $f = (f_i)_{i \in I}$ (and dual).

When there are only a finite number of objects X_1, \dots, X_n , we shall write $X_1 \times \dots \times X_n$ (resp. $X_1 \sqcup \dots \sqcup X_n$). When all X_i are equal to the same X , we shall write X^I (resp. $X^{(I)}$). Note that a final object is nothing else but the empty product (and dual).

Examples

1. In **Set**, the cartesian product is a product and the disjoint union is a coproduct.
2. In **Top**, this is the same thing with the coarser (resp. finer) topology making the projections (resp. injections) continuous.
3. In **Ab**, the cartesian product with termwise addition is a product and the direct sum is a coproduct⁴ (product equals coproduct when I is finite).

The morphism $\delta = (\text{Id}_X, \text{Id}_X) : X \rightarrow X \times X$ is called the *diagonal morphism*. The morphism $\tau = (p_2, p_1) : X \times Y \rightarrow Y \times X$ is called the *flip morphism*. The morphism $(p_1, p_2, p_1, p_3) : X \times Y \times Z \rightarrow X \times Y \times X \times Z$ (resp. (p_1, p_3, p_2, p_3)) is called the *left* (resp. *right*) *distribution morphism*.

Exercise 1.3.12 Show that there exists “canonical” isomorphisms (and dual)

1. $X \times 1_{\mathcal{C}} \simeq X$ for $X \in \mathcal{C}$,
2. $X \times Y \simeq Y \times X$ for $X, Y \in \mathcal{C}$ and
3. $(X \times Y) \times Z \simeq X \times (Y \times Z)$ for $X, Y, Z \in \mathcal{C}$.

Solution. We only do the last one. There exists obvious morphisms $p_1 : (X \times Y) \times Z \rightarrow X \times Y \rightarrow X$, $p_2 : (X \times Y) \times Z \rightarrow X \times Y \rightarrow Y$ and $p_3 : (X \times Y) \times Z \rightarrow Z$ defining a cone for the discrete diagram (X, Y, Z) . Conversely, given such a cone $(f_1 : T \rightarrow X, f_2 : T \rightarrow Y, f_3 : T \rightarrow Z)$, there exists a unique morphism $f' : T \rightarrow X \times Y$ whose components are f_1 and f_2 . Then, there exists a unique morphism $f : T \rightarrow (X \times Y) \times Z$ whose components are f' and f_3 . In other words, there exists a unique morphism satisfying $p_i \circ f = f_i$ for $i = 1, 2, 3$. This shows that $(X \times Y) \times Z$ is a limit for (X, Y, Z) . Following the same lines, one sees that $X \times (Y \times Z)$ also is a limit for this diagram. They are therefore isomorphic. ■

Definition 1.3.13 A limit X of a diagram $(X_1 \xrightarrow{f_1} X_0 \xleftarrow{f_2} X_2)$ is called a *fibered product* of X_1 and X_2 over X_0 and denoted by $X_1 \times_{X_0} X_2$. We shall then also say that the diagram

$$\begin{array}{ccc} X & \xrightarrow{p_1} & X_1 \\ \downarrow p_2 & \lrcorner & \downarrow f_1 \\ X_2 & \xrightarrow{f_2} & X_0 \end{array} \quad (1.3)$$

is *cartesian* or that p_2 is the *pullback* of f_1 along f_2 (and p_1 is the *pullback* of f_2 along f_1). Dually, there exists the notions of a *fibered coproduct* denoted by $X_1 \sqcup_{X_0} X_2$ associated to a diagram $(X_1 \xleftarrow{f_1} X_0 \xrightarrow{f_2} X_2)$, a *cocartesian* square and a *pushout*.

⁴Be careful that what is called *free product* is a coproduct in the category of (non abelian) groups.

Thus, by definition, diagram (1.3) is cartesian if, given $g_1 : Y \rightarrow X_1$ and $g_2 : Y \rightarrow X_2$ such that $f_1 \circ g_1 = f_2 \circ g_2$, there exists a unique $g : Y \rightarrow X$ such that $p_1 \circ g = g_1$ and $p_2 \circ g = g_2$. Note that a product of two objects is nothing but a fibered product over a final object (and dual). Conversely, a fibered product as in the definition is the same thing as the product in $\mathcal{C}_{/X_0}$ (and dual).

Examples 1. In **Set**, we have

$$X_1 \times_{X_0} X_2 = \{(x_1, x_2) \mid f_1(x_1) = f_2(x_2)\} \subset X_1 \times X_2$$

and

$$X_1 \sqcup_{X_0} X_2 = (X_1 \sqcup X_2) / \sim$$

where \sim is the relation $f_1(x_0) \sim f_2(x_0)$ when $x_0 \in X_0$.

2. In **Top**, this is the same thing with the induced (resp. quotient) topology.

3. In **Ab** we have with $(M_1 \xrightarrow{f_1} M_0 \xleftarrow{f_2} M_2)$

$$M_1 \times_{M_0} M_2 = \ker(M_1 \oplus M_2 \rightarrow M_0, (x_1, x_2) \mapsto f(x_2) - f(x_1))$$

and with $(M_1 \xleftarrow{f_1} M_0 \xrightarrow{f_2} M_2)$

$$M_1 \sqcup_{M_0} M_2 = \operatorname{coker}(M_0 \rightarrow M_1 \oplus M_2 \rightarrow M_0, x_0 \mapsto (f_1(x_0), f_2(x_0))).$$

Exercise 1.3.14 Show that the fibered coproduct in the category of commutative rings is tensor product.

Exercise 1.3.15 Show that^a, in a diagram

$$\begin{array}{ccccc} Y_2 & \longrightarrow & Y_1 & \xrightarrow{\quad} & Y_0 \\ \downarrow & & \downarrow & \lrcorner & \downarrow \\ X_2 & \longrightarrow & X_1 & \longrightarrow & X_0, \end{array}$$

if the right hand square is cartesian, then the left hand square is cartesian if and only if the full rectangle is cartesian.

^aThis is an illustration of the more general proposition 1.3.39 below.

Definition 1.3.16 A limit X of a pair $Y \rightrightarrows Z$ is called a *kernel* (or *equalizer*) and denoted by $X = \ker(f, g)$. We shall also say that the sequence

$$X \xrightarrow{i} Y \xrightleftharpoons[f]{g} Z \tag{1.4}$$

is *left exact*. Dually, there exists the notion of a *cokernel* (or *coequalizer*) $\operatorname{coker}(f, g)$ and *right exact* sequence.

Thus, by definition, diagram (1.4) is left exact if, given $h : T \rightarrow Y$ such that $f \circ h = g \circ h$, there exists a unique $k : T \rightarrow X$ such that $h = i \circ k$.

Examples 1. In **Set**, we have

$$\ker(f, g) = \{y \in Y \mid f(y) = g(y)\} \quad \text{and} \quad \text{coker}(f, g) = Z / \sim$$

where \sim is the relation $f(y) \sim g(y)$ for $y \in Y$.

2. Same thing in **Top** with induced and quotient topology respectively.
3. In **Ab**, $\ker(f, g) = \ker(g - f)$ (and dual) .

Exercise 1.3.17 Show that if \mathcal{C} is a small category, then $\mathbf{Op}(\mathcal{C})$ is the kernel of the domain and codomain functors $\mathbf{Mor}(\mathcal{C}) \rightrightarrows \mathcal{C}$ in **Cat**.

Exercise 1.3.18 Make explicit specific limits (final object, products, fibered product and kernel) and colimits (initial object, coproduct, fibered coproduct and cokernel) in **Mon**, **Grp**, $G\text{-Set}$, $A\text{-Mod}$ or **Cat**.

1.3.3 Monomorphism/epimorphism

These notions are usually introduced before that of limits/colimits but are better understood as particular instances of limits/colimits.

Definition 1.3.19 A morphism $i : X \rightarrow Y$ is called a

1. a *monomorphism* if the diagram

$$\begin{array}{ccc} X & \xlongequal{\quad} & X \\ \parallel & & \downarrow i \\ X & \xrightarrow{i} & Y, \end{array} \tag{1.5}$$

is cartesian and we shall then write $i : X \rightarrowtail Y$.

2. a *regular monomorphism* if there exists an exact sequence $X \xrightarrow{i} Y \rightrightarrows Z$.
3. a *split monomorphism* if it has a retraction.

The dual notion is that of a (*regular, split*) *epimorphism* and we shall then write $Y \twoheadrightarrow X$.

- Examples**
1. A morphism in **Set** (resp. **Ab**, resp. **Top**) is a monomorphism/epimorphism if and only if it is injective/surjective.
 2. A regular monomorphism/epimorphism in **Top** is a homeomorphism with a subspace/from a quotient map.
 3. The inclusion map $\mathbb{Q} \rightarrowtail \mathbb{R}$ is at the same time a monomorphism and an epimorphism in the category of *Hausdorff* topological spaces.
 4. Same for the inclusion map $\mathbb{Z} \rightarrowtail \mathbb{Q}$ in the category of rings.

Proposition 1.3.20 The following are equivalent (and dual):

1. The morphism $i : X \rightarrow Y$ is a monomorphism,
2. if $f, g : Z \rightarrow X$ satisfy $i \circ f = i \circ g$, then $f = g$,
3. for all $Z \in \mathcal{C}$, the map $i_* : \text{Hom}(Z, Y) \rightarrow \text{Hom}(Z, X)$ is injective,

4. any commutative diagram

$$\begin{array}{ccc} Y' & \xlongequal{\quad} & Y' \\ \downarrow j & & \downarrow \\ X & \xrightarrow{i} & Y \end{array} \quad (1.6)$$

is cartesian.

Proof. By definition, the morphism i is a monomorphism if and only if, given $f, g : Z \rightarrow X$ satisfying $i \circ f = i \circ g$, there exists a unique h such that $f = g = h$. This is clearly the same thing as the second assertion which in turn is an explicit form of the third one. Assume now that we are given a commutative diagram (1.6) as well as two morphisms $f, g : Z \rightarrow X'$ satisfying $i \circ j \circ f = i \circ g$. If i is a monomorphism, then $j \circ f = g$ and the diagram is therefore cartesian. Finally, in the last assertion, if we put i on the right hand side, then we fall back onto the definition of a monomorphism. ■

Exercise 1.3.21 Show that a split monomorphism is a regular monomorphism and that a regular monomorphism is a monomorphism (and dual).

Exercise 1.3.22 Show that a morphism $f : X \rightarrow Y$ is a monomorphism (resp. an epimorphism) if and only if the induced functor $\mathcal{C}_X \rightarrow \mathcal{C}_Y$ (resp. $Y \backslash \mathcal{C} \rightarrow X \backslash \mathcal{C}$) is fully faithful.

Exercise 1.3.23 Show that^a if both $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ are monomorphisms, then so is $g \circ f$ (and dual). Show that, conversely, if $g \circ f$ is a monomorphism, then so is f (and dual).

^aBe careful that this is not the case for regular monomorphisms in general.

Exercise 1.3.24 Show that, if we are given $f, g : X \rightarrow Y$ and $i : Y \rightarrow Z$ is a monomorphism, then^a $\ker(i \circ f, i \circ g) = \ker(f, g)$ (and dual). Analog for fibered products (and dual) ?

^aIf one of them exists, then so does the other and...

Exercise 1.3.25 Show that, if we are given a commutative diagram of monomorphisms

$$\begin{array}{ccc} X_1 & \xrightleftharpoons{\quad} & X_2 \\ & \searrow & \swarrow \\ & Y & \end{array}$$

then, the upper arrows are inverse isomorphisms to each other (and dual).

Exercise 1.3.26 Show that, in a cartesian diagram

$$\begin{array}{ccc} X' & \xrightarrow{f'} & Y' \\ \downarrow & \lrcorner & \downarrow \\ X & \xrightarrow{f} & Y \end{array}$$

if f is a monomorphism, then f' also is monomorphism (and dual).

Exercise 1.3.27 Show that if $i : X \rightarrowtail Y$ is an injective map and the diagram of sets

$$\begin{array}{ccc} X & \xrightarrow{i} & Y \\ \downarrow i & & \downarrow \\ Y & \longrightarrow & Z \end{array}$$

is cocartesian, then it is also cartesian.

Recall that we always implicitly assume that all limits and colimits that we write do exist. For safety, the reader can assume up to the end of the section that all finite limits and all finite colimits exist.

Definition 1.3.28 The (*regular*) *image* of a morphism $f : X \rightarrow Y$ is

$$\mathrm{im}(f) := \ker(Y \rightrightarrows Y \sqcup_X Y).$$

The dual notion is that of a (*regular*) *coimage*

$$\mathrm{coim}(f) := \mathrm{coker}(X \times_Y X \rightrightarrows X).$$

- Examples**
1. We already know $\mathrm{im}(f)$ in **Set** and $\mathrm{coim}(f) = X / \sim$ with $x \sim x' \Leftrightarrow f(x) = f(x')$.
 2. In **Top**, this is the same thing with induced and quotient topology respectively.
 3. In **Ab** also, we know $\mathrm{im}(f)$ for $f : M \rightarrow N$ and we have $\mathrm{coim}(f) = M / \ker(f)$.

Exercise 1.3.29 Show that, if $f : X \rightarrow Y$ is an epimorphism, then $\mathrm{im}(f) = Y$ (and dual).

Proposition 1.3.30 The factorization $f : X \rightarrow \mathrm{im}(f) \rightarrowtail Y$ is universal for factorizations of a morphism $f : X \rightarrow Y$ through a *regular* monomorphisms (and dual).

Proof. We have to show that, when $Y' \rightarrowtail Y$ is a regular monomorphism, then any commutative diagram may be uniquely completed as follows:

$$\begin{array}{ccccc} X & \longrightarrow & Y' & \rightarrowtail & Y \\ & \searrow & \uparrow & & \nearrow \\ & & \mathrm{im} f & & \end{array}$$

By definition, there exists a left exact sequence $Y' \rightarrow Y \rightrightarrows T$ and we may then contemplate the commutative diagram

$$\begin{array}{ccccc} Y' & \xrightarrow{\quad} & Y & \rightrightarrows & T \\ \uparrow & \nearrow \text{dashed} & \parallel & & \uparrow \\ X & \longrightarrow & \text{im } f & \longrightarrow & Y \rightrightarrows Y \sqcup_X Y. \end{array}$$

If we remove the regular condition in proposition 1.3.30, we then obtain the definition of a “non-regular” image. In general, there always exists a factorization of a morphism $f : X \rightarrow Y$ as follows:

$$\begin{array}{ccccccc} X \times_Y X & \rightrightarrows & X & \xrightarrow{f} & Y & \rightrightarrows & Y \sqcup_X Y \\ & & \downarrow & & \uparrow & & \\ & & \text{coim}(f) & \dashrightarrow & \text{im}(f) & & \end{array}$$

Definition 1.3.31 A morphism f is said to be *strict* if $\text{coim}(f) \simeq \text{im}(f)$.

Example Any morphism is strict in **Set** or **Ab**⁵ but not in **Top**.

Exercise 1.3.32 Show that a strict epimorphism is regular (and dual).

Definition 1.3.33 A category is said to be *balanced* if a morphism which is at the same time a monomorphism and an epimorphism is automatically an isomorphism.

Note that the converse is always true.

Example The categories **Set** and **Ab** are balanced but **Top** and **Ring** are not.

Proposition 1.3.34 Assume \mathcal{C} is balanced.

1. A strict morphism factors uniquely up to an isomorphism as an epimorphism followed by a monomorphism,
2. A strict monomorphism is automatically regular (and dual).

Proof. Factorization is clear since a strict morphism $f : X \rightarrow Y$ splits as $f : X \rightarrow \text{coim}(f) = \text{im}(f) \rightarrow Y$. Conversely, any such decomposition fits into a commutative diagram⁶

$$\begin{array}{ccccccc} X \times_Y X & \rightrightarrows & X & \longrightarrow & Z & \longrightarrow & Y \rightrightarrows Y \sqcup_X Y \\ & & \downarrow & \nearrow \text{dashed} & \nearrow \text{dashed} & & \uparrow \\ & & \text{coim}(f) & \xrightarrow{\quad} & \text{im}(f) & & \end{array}$$

This implies that $\text{coim}(f) \simeq Z \simeq \text{im}(f)$ when the category is balanced. The other assertion is a consequence of exercise 1.3.29. ■

⁵This is Noether’s first isomorphism.

⁶Be careful that we cannot use the universal property of image or coimage here because it is not assumed that the monomorphism nor epimorphism is regular in the factorization.

1.3.4 Constructions of limits

Lemma 1.3.35 We give ourselves a commutative diagram (X_i, f_α) and we assume that both $X' := \prod_i X_i$ and $X'' := \prod_{\alpha: i \rightarrow j} X_j$ exist. Denote by p (resp. q) the morphism $X' \rightarrow X''$ induced by the projections p_j onto the codomain (resp. the composites of the projection p_i onto the domain and f_α). Then, $X = \varprojlim (X_i, f_\alpha)$ if and only if there exists a left exact sequence

$$X \longrightarrow X' \rightrightarrows_q^p X''.$$

In particular, $X \rightarrowtail X'$ is then a regular monomorphism.

Hint. By definition, a cone for our diagram is object Y together with a family of morphisms $g_i : Y \rightarrow X_i$ satisfying $g_j = f_\alpha \circ g_i$ whenever $\alpha : i \rightarrow j$. This is the same thing as a morphism $g : Y \rightarrow X'$ satisfying $p \circ g = q \circ g$. Same cones, same limits. ■

Example If $(f_\alpha : X_i \rightarrow X_j)$ is a commutative diagram of *sets*, then

$$\varprojlim X_i = \left\{ (x_i) \in \prod X_i, \forall \alpha : i \rightarrow j, f_\alpha(x_i) = x_j \right\}$$

and

$$\varinjlim X_i = \coprod X_i / \sim$$

where \sim denotes the (smallest equivalence relation) such that by $x_i \sim x_j$ whenever $f_\alpha(x_i) = x_j$.

Exercise 1.3.36 Assume given $f, g : X \rightarrow Y$ such that $Y \times Y$ exists. Show that $\ker(f, g)$ exists if and only if there exists a cartesian diagram

$$\begin{array}{ccc} Z & \xrightarrow{\quad} & X \\ \downarrow & \lrcorner & \downarrow (f, g) \\ Y & \xrightarrow{\delta} & Y \times Y \end{array}$$

in which case $Z = \ker(f, g)$.

Proposition 1.3.37 Let \mathcal{C} be a category.

1. If all (finite) products and all kernels exist, then all (finite) limits exist (and dual).
2. If all fibered products exist and there is a final object, then all finite limits exist (and dual).

Proof. Immediate consequence of exercises 1.3.35 and 1.3.36. ■

When all (finite) limits exist, then \mathcal{C} is said to be (*finitely*) *complete*. The dual terminology is *cocomplete*. When both properties hold, we say *bicomplete*. Be careful that a complete (or cocomplete) *small* category is automatically (associated to) a preorder:

$$|\mathrm{Mor}(\mathcal{C})| \geq |\mathrm{Hom}(X, Y^{\mathrm{Mor}(\mathcal{C})})| = |\mathrm{Hom}(X, Y)|^{|\mathrm{Mor}(\mathcal{C})|} \Rightarrow |\mathrm{Hom}(X, Y)| \leq 1.$$

■ **Exercise 1.3.38** Show that **Set**, **Top**, **Ab**, etc. are bicomplete.

The next statement is a consequence of coming corollary 1.4.15:

Proposition 1.3.39 If all limits of shape I exist in \mathcal{C} , then all limits of shape I exist in \mathcal{C}^J and, if $D \in (\mathcal{C}^J)^I \simeq \mathcal{C}^{I \times J}$, then

1. $\forall j \in J, \left(\varprojlim_I D \right) (j) = \varprojlim_I D(-, j)$ and
2. $\varprojlim_J \left(\varprojlim_I D \right) \simeq \varprojlim_{I \times J} D \simeq \varprojlim_I \left(\varprojlim_J D \right).$

And dually for colimits.

Proof. Let us write $D'_j := \varprojlim_I D(-, j) \in \mathcal{C}$ for $j \in J$. Any morphism $j \rightarrow k$ in J will provide a compatible family of morphism $D(i, j) \rightarrow D(i, k)$ for $i \in I$, and taking limits on I , a morphism $D'_j \rightarrow D'_k$. This is functorial and provides $D' \in \mathcal{C}^J$. By construction, this is a cone on D : there exists morphisms $D'_j \rightarrow D(i, j)$ for all $(i, j) \in I \times J$ providing a morphism $\underline{D}' \rightarrow D$ in $\mathcal{C}^{I \times J}$. If we are given another cone $\underline{E} \rightarrow D$ with $E \in \mathcal{C}^J$, then there exists for each j , a morphism

$$E_j := \varprojlim_I \underline{E}(-, j) \rightarrow \varprojlim_I D(-, j) = D'_j.$$

Again, they are compatible and provide $E \rightarrow D'$. This shows that $D' = \varprojlim_I D$. The first assertion then follows directly from the definition of D' and it remains to show that $\varprojlim_J D' \simeq \varprojlim_{I \times J} D$. We give ourselves a cone $\underline{X} \rightarrow D$ with $X \in \mathcal{C}$. It induces a cone $\underline{X} \rightarrow D(-, j)$ in \mathcal{C}^I for all $j \in J$ and taking limits on I , a morphism $X \rightarrow D'_j$. Again, they are compatible on J and combine (uniquely) into a morphism $\underline{X} \rightarrow D'$ in \mathcal{C}^J . Taking limits, we obtain $X \rightarrow \varprojlim_J D'$ showing that this is a universal cone. ■

This assertion essentially says that limits and colimits (of diagrams) are computed termwise.

Note that there also exists a morphism

$$\varinjlim_J \left(\varprojlim_I D \right) \longrightarrow \varprojlim_I \left(\varinjlim_J D \right)$$

which is not an isomorphism in general.

■ **Definition 1.3.40** Given a cocone $D \rightarrow \underline{Y}$, we shall say that the colimit of D is *stable under pullback* along $Y' \rightarrow Y$ if

$$\varinjlim (D \times_Y Y') \simeq (\varinjlim D) \times_Y Y'.$$

The dual notion is that of limit being *stable under pushout*.

■ **Exercise 1.3.41** Show that colimits of *sets* are stable under pullback.

As a particular example, epimorphisms are stable under pullback if, given any

cartesian diagram

$$\begin{array}{ccc} X' & \xrightarrow{f'} & Y' \\ \downarrow & \lrcorner & \downarrow \\ X & \xrightarrow{f} & Y, \end{array}$$

if f is an epimorphism, then f' also is an epimorphism.

Examples 1. Epimorphisms are stable under pull back in **Set**, **Ab** or **Top** (and dual).

2. This is not the case in **Haus** or **Ring** :

$$\begin{array}{ccc} \emptyset & \xrightarrow{\quad} & \mathbb{R} \setminus \mathbb{Q} \\ \downarrow & \lrcorner & \downarrow \\ \mathbb{Q} & \xrightarrow{\quad} & \mathbb{R} \end{array} \quad \text{and} \quad \begin{array}{ccc} \mathbb{Q} & \xrightarrow{\quad} & \mathbb{Q}[s] \\ \downarrow & \lrcorner & \downarrow \\ \mathbb{Q}[t] & \xrightarrow{\quad} & \mathbb{Q}[t, s]/(ts - 1). \end{array}$$

3. Be careful that, in the category of sets, we have

$$\varinjlim_{I \times J} (X_i \times Y_j) = \left(\varinjlim_I X_i \right) \times \left(\varinjlim_J Y_j \right) \text{ but } \varinjlim_I (X_i \times Y_i) \neq \left(\varinjlim_I X_i \right) \times \left(\varinjlim_I Y_i \right)$$

in general.

1.3.5 Preservation of limit

Recall that any functor $F : \mathcal{C} \rightarrow \mathcal{D}$ provides by composition a functor

$$F^I : \mathcal{C}^I \rightarrow \mathcal{D}^I, \quad D \mapsto F^I(D) := F \circ D.$$

We shall usually simply write F instead of F^I so that $F(X_i, f_\alpha) = (F(X_i), F(f_\alpha))$.

Definition 1.3.42 If D is a commutative diagram in \mathcal{C} , then a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is said to *preserve* (or *commute with*) the limit of D , if

$$F(\varprojlim D) \simeq \varprojlim F(D).$$

Of course, it is assumed here that the limit of D exists and it implies that the limit of $F(D)$ also does. There exists an obvious analog for colimits. Be careful that a (contravariant) functor $F : \mathcal{C}^{\text{op}} \rightarrow \mathcal{D}$ preserves a limit when it turns a *colimit* in \mathcal{C} into limit in \mathcal{D} .

Definition 1.3.43 A functor is said to be *left exact* (resp. *right exact*) if it preserves all finite limits (resp. colimits). It is said to be *exact* if it is both left and right exact^a.

^aWe should say exact, coexact and biexact respectively but we will follow the mainstream terminology.

A functor that preserves all limits (resp. colimits) is sometimes said to be *continuous* (resp. *cocontinuous*) but this may conflict with the vocabulary of topos theory.

- Examples**
1. The forgetful functor $\mathbf{Top} \rightarrow \mathbf{Set}$ preserves all limits and colimits.
 2. The functor $\mathbf{Set} \rightarrow \mathbf{Top}$ that endows a set with the discrete topology is exact and preserves all colimits. It does not preserve infinite limits however: with $\bar{n} = \{0, 1, \dots, n-1, \infty\}$ (discrete), we have $\varprojlim \bar{n} \simeq \bar{\mathbb{N}} = \mathbb{N} \cup \{\infty\}$ infinite and compact (not discrete).
 3. The functor $\mathbf{Set} \rightarrow \mathbf{Top}$ that endows a set with the coarse topology preserves all limits but is not exact: we have $1 \sqcup 1 = 2$.
 4. The forgetful functor $\mathbf{Ab} \rightarrow \mathbf{Set}$ preserves all limits and the free abelian group functor $\mathbf{Set} \rightarrow \mathbf{Ab}$ preserves all colimits.
 5. The functors h^X and h_X preserve all limits. Be careful that the second one is contravariant: it turns colimits into limits.
 6. The forgetful functor $A\text{-}\mathbf{Mod} \rightarrow \mathbf{Ab}$ preserves all limits and colimits.
 7. For fixed $M \in \mathbf{Mod}\text{-}A$, the functor

$$A\text{-}\mathbf{Mod} \rightarrow \mathbf{Ab}, \quad N \mapsto M \otimes_A N$$

preserves all colimits.

Exercise 1.3.44 Show that, if F is left exact, then F preserves (regular, strict) monomorphisms (and dual).

Exercise 1.3.45 Show that, if \mathcal{C} is a small category and all limits of shape I exist in \mathcal{D} , then all limits of shape I also exist in the category $\widehat{\mathcal{C}}(\mathcal{D})$ of presheaves and they are preserved by the functor

$$\widehat{\mathcal{C}}(\mathcal{D}) \rightarrow \mathcal{D}, \quad F \mapsto F(X)$$

for fixed $X \in \mathcal{C}$ (and dual).

Hint. This is a variant of proposition 1.3.39. ■

Exercise 1.3.46 Show that a representable functor $F : \mathcal{C} \rightarrow \mathbf{Set}$ preserves all limits^a.

^aThere exists no dual statement and the notion of a limit plays a special role.

Solution. We may assume that $F = h^X$ with $X \in \mathcal{C}$. It is then sufficient to check that if D is a commutative diagram in \mathcal{C} , then we have a sequence of bijections

$$\begin{aligned} h^X(\varprojlim D) &\simeq \mathrm{Hom}(X, \varprojlim D) \simeq \mathrm{Hom}(\underline{X}, D) \\ &\simeq \mathrm{Hom}(\{0\}, h^X(D)) \simeq \mathrm{Hom}(\{0\}, \varprojlim h^X(D)) \simeq \varprojlim h^X(D). \end{aligned}$$

Only the middle one needs to be checked by hand: if we write $D =: (X_i, f_\alpha)$, then, giving a morphism $\{0\} \rightarrow h^X(D)$ is equivalent to give a compatible family of maps $\{0\} \rightarrow \mathrm{Hom}(X, X_i)$, or in other words, to give for each $i \in I$, a morphism $g_i : X \rightarrow X_i$ such that $f_\alpha \circ g_i = g_j$, or finally a morphism $\underline{X} \rightarrow D$. ■

Proposition 1.3.47 1. The following are equivalent:

- (a) The functor F preserves all limits,
- (b) F preserves all products and all kernels.

2. The following are equivalent:

- (a) The functor F is left exact,
- (b) F preserves all finite products and all kernels,
- (c) F preserves all fibered products and the final object.

Analogous for colimits.

Proof. Immediate consequence of exercises 1.3.35 and 1.3.36. ■

Definition 1.3.48 1. A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is said to *reflect* the limit of a diagram D in \mathcal{C} if any cone $\underline{X} \rightarrow D$ satisfying $F(X) \simeq \varprojlim F(D)$ also satisfies $X \simeq \varprojlim D$.

2. The functor F is said to be *conservative* if a morphism f of \mathcal{C} is automatically an isomorphism when $F(f)$ is an isomorphism.

One defines dually the property of reflecting colimits. Also, one says that a functor *reflects* monomorphisms (resp. epimorphisms, resp. isomorphisms) if a morphism f of \mathcal{C} is automatically a monomorphism (resp. an epimorphism, resp. an isomorphism) when $F(f)$ is so. Thus, the functor F is conservative if and only if it reflects isomorphisms if and only if it reflects limits of diagrams of shape **1**.

Examples 1. A functor that reflects fibered products reflects monomorphisms (and dual)

2. The forgetful functor $\mathbf{Ab} \rightarrow \mathbf{Set}$ reflects limits but the forgetful functor $\mathbf{Top} \rightarrow \mathbf{Set}$ does not.

Exercise 1.3.49 Show that

- 1. a fully faithful functor reflects all limits and colimits,
- 2. a conservative functor reflects all limits and colimits that it preserves,
- 3. a faithful functor reflects all monomorphisms and epimorphisms.

Exercise 1.3.50 Show that if \mathcal{C} is balanced, then any faithful functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is conservative.

1.3.6 Filtered colimit

Definition 1.3.51 A small category I is said to be *filtered* if any finite diagram in I has a cocone.

A *filtered diagram* is a diagram $I \rightarrow \mathcal{C}$ with I filtered. A *filtered colimit* is a colimit of a filtered diagram (and dual).

Example The category associated to a preordered set is filtered if and only if the preordered set is directed (any finite subset has an upper bound).

Exercise 1.3.52 Show that a small category I is filtered if and only if

- 1. $I \neq \emptyset$,
- 2. $\forall i, j \in I, \exists k \in I, i \rightarrow k, j \rightarrow k$,
- 3. $\forall u, v : i \rightarrow j, \exists k \in I, c : j \rightarrow k / c \circ u = c \circ v$.

In the computation of colimits, one may always replace a filtered category with a directed set:

Proposition 1.3.53 If I is a filtered category, then there exists a directed set J and a functor $u : J \rightarrow I$ such that, for all diagram $D : I \rightarrow \mathcal{C}$, if $\varinjlim (D \circ u)$ exists, then $\varinjlim D$ exists and $\varinjlim (D \circ u) \simeq \varinjlim D$.

Proof. To do (hard). ■

Exercise 1.3.54 Show that

1. a set (resp. a small category) is the filtered colimit of its finite subsets (resp. finite subcategories),
2. a category with finite colimits (resp. finite coproducts) and filtered colimits has all colimits (resp. all coproducts).
3. a functor that preserves finite colimits (resp. finite coproducts) and filtered colimits preserves all colimits (resp. coproducts).

Filtered colimits of *sets* are exact:

Proposition 1.3.55 If I is a *filtered* category, J is a *finite* category and D is a diagram of *sets* of shape $I \times J$, then

$$\varinjlim_I \varprojlim_J D \simeq \varprojlim_J \varinjlim_I D$$

Proof. According to proposition 1.3.53, we may assume that I is a directed set. Moreover, thanks to proposition 1.3.47 it is sufficient to treat the cases of a final object, a product of two objects or a kernel. We shall use the fact that colimits of diagrams are computed termwise (proposition 1.3.39). In particular, the case of a final object is trivial (colimit of a constant diagram). We consider now two families of morphisms $(f_{ij} : X_i \rightarrow X_j)$ and $(g_{ij} : Y_i \rightarrow Y_j)$ defined for $i < j$ with $f_{jk} \circ f_{ij} = f_{ik}$ and $g_{jk} \circ g_{ij} = g_{ik}$ when $i < j < k$. Since I is directed, we have

$$\varinjlim X_i = \coprod X_i / \sim \quad \text{with} \quad x_i \sim x_j \Leftrightarrow \exists k \geq i, j, f_{ik}(x_i) = f_{jk}(x_j)$$

for $x_i \in X_i$ and $x_j \in X_j$. We shall denote by \bar{x}_i the class of $x_i \in X_i$ and use analogous notations for all diagrams of shape I . By considering colimits on the projections, we have the obvious morphisms

$$\varinjlim (X_i \times Y_i) \rightarrow \varinjlim X_i \quad \text{and} \quad \varinjlim (X_i \times Y_i) \rightarrow \varinjlim Y_i.$$

The universal property of products provides us with a map

$$\varinjlim (X_i \times Y_i) \rightarrow \varinjlim X_i \times \varinjlim Y_i, \quad \overline{(x_i, y_i)} \mapsto (\bar{x}_i, \bar{y}_i)$$

and we have to show that it is bijective. Assume that $(\bar{x}_i, \bar{y}_i) = (\bar{x}'_j, \bar{y}'_j)$. As explained above, there exists $k \geq i, j$ such that $f_{ik}(x_i) = f_{jk}(x'_j)$ and $\ell \geq i, j$ such that $g_{i\ell}(y_i) = g_{j\ell}(y'_j)$. Since I is directed, there exists $m \geq k, \ell$ and we will have $(f_{im}(x_i), g_{im}(y_i)) = (f_{jm}(x'_j), g_{jm}(y'_j))$ so that $\overline{(x_i, y_i)} = \overline{(x'_j, y'_j)}$. This shows injectivity. Surjectivity is proven in the same way. This takes care of the product of two sets and we will now treat the case of kernels. We give ourselves two families of

maps $(\varphi_i, \psi_i : X_i \rightarrow Y_i)$ compatible with all f_{ij} 's and g_{ij} 's. The universal property of kernels provides us with a map

$$\varinjlim \ker(\varphi_i, \psi_i) \rightarrow \ker(\varinjlim \varphi_i, \varinjlim \psi_i).$$

Injectivity and surjectivity are then shown as before, using the explicit description of filtered colimits of sets. ■

Exercise 1.3.56 Show that filtered colimits are exact in **Ab**, **Top**, etc. and that they are preserved by the forgetful functors to **Set** (one may show the second assertion first).

Exercise 1.3.57 Show that **Ab** satisfies AB6 extra condition: filtered colimits commute with arbitrary products:

$$\prod_{j \in J} \varinjlim_{i_j \in I_j} M_{i_j} \simeq \varinjlim_{\prod_{j \in J} I_j} \prod_{j \in J} M_{i_j}$$

when each I_j is filtered.

Definition 1.3.58 An *ind-object* “ $\varinjlim X_i$ ” of a category \mathcal{C} is a filtered diagram $(X_i)_{i \in I}$. They form a category $\text{Ind}(\mathcal{C})$ with

$$\text{Hom}(\varinjlim X_i, \varinjlim Y_j) = \varprojlim_{i \in I} \varinjlim_{j \in J} \text{Hom}(X_i, Y_j).$$

The dual notion is that of a *pro-object* “ $\varprojlim X_i$ ” and they form a category $\text{Pro}(\mathcal{C}) := \text{Ind}(\mathcal{C}^{\text{op}})^{\text{op}}$.

One gets an equivalent category by considering only directed sets instead of filtered categories.

Example The category of profinite sets (pro-objects of the category of finite sets) is equivalent to the category of totally disconnected compact Hausdorff spaces.

Exercise 1.3.59 Show that the obvious functor $\mathcal{C} \rightarrow \text{Ind}(\mathcal{C})$ is fully faithful, exact and preserves all limits.

1.4 Adjointness

1.4.1 Definition

Definition 1.4.1 A functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is said to be *adjoint* to a functor $G : \mathcal{C}' \rightarrow \mathcal{C}$ if there exists a natural isomorphism

$$\forall X \in \mathcal{C}, X' \in \mathcal{C}', \quad \Phi_{X, X'} : \text{Hom}(F(X), X') \simeq \text{Hom}(X, G(X')).$$

One sometimes writes $F \dashv G$ or $F : \mathcal{C} \rightleftarrows \mathcal{C}' : G$ for short. The dual notion is that of a *coadjoint* so that G is coadjoint to F if and only if F is adjoint to G . It may be useful to write down explicitly what it means for $\Phi_{X, X'}$ and its inverse to be natural: given $g : Y \rightarrow X, g' : X' \rightarrow Y'$, then

$$\forall f : F(X) \rightarrow X', \quad \Phi_{Y, Y'}(g' \circ f \circ F(g)) = G(g') \circ \Phi_{X, X'}(f) \circ g \quad \text{and}$$

$$\forall f' : X' \rightarrow G(X), \quad \Phi_{Y, Y'}^{-1}(G(g') \circ f \circ g) = g' \circ \Phi_{X, X'}^{-1}(f) \circ F(g).$$

Examples 1. The forgetful functor $\mathbf{Top} \rightarrow \mathbf{Set}$ has both an adjoint (discrete topology) and a coadjoint (coarse topology):

$$\mathcal{C}(X^{\text{disc}}, Y) \simeq \mathcal{F}(X, Y) \quad \text{et} \quad \mathcal{F}(X, Y) \simeq \mathcal{C}(X, Y^{\text{coarse}}).$$

2. The forgetful functor $\mathbf{Ab} \rightarrow \mathbf{Set}$ has an adjoint (free abelian group):

$$\text{Hom}(\mathbb{Z} \cdot X, M) \simeq \mathcal{F}(X, M).$$

3. The inclusion $\mathbf{Grp} \hookrightarrow \mathbf{Mon}$ has both an adjoint and a coadjoint: for a group G and a monoid H , we have

$$\text{Hom}(H^{\text{gr}}, G) \simeq \text{Hom}(H, G) \quad \text{et} \quad \text{Hom}(G, H) \simeq \text{Hom}(G, H^{\times}).$$

Exercise 1.4.2 Show that most forgetful and inclusion functors we have already met have an adjoint (and sometimes a coadjoint) and make them explicit.

Exercise 1.4.3 Show that the adjoint to the forgetful functor $k\text{-}\mathbf{Alg} \rightarrow k\text{-}\mathbf{Mod}$ is the tensor algebra functor $M \mapsto T(M)$. Same thing with $S(M)$ when we restrict to commutative algebras.

Exercise 1.4.4 Show that, if $f : A \rightarrow B$ is a morphism of rings, then the forgetful functor $B\text{-}\mathbf{Mod} \rightarrow A\text{-}\mathbf{Mod}$ has both an adjoint $M \mapsto B \otimes_A M$ and a coadjoint $M \mapsto \text{Hom}_A(B, M)$.

Exercise 1.4.5 Show that (for fixed Y) the functor $X \mapsto X \times Y$ from \mathbf{Set} to itself is adjoint to the functor $Z \mapsto \mathcal{F}(Y, Z)$:

$$\mathcal{F}(X \times Y, Z) \simeq \mathcal{F}(X, \mathcal{F}(Y, Z)).$$

This is called *Currying*. Write down the analogous statements for \mathbf{Cat} and \mathbf{Ab} .

Exercise 1.4.6 Show that if both F_1 and F_2 are adjoint to G , then $F_1 \simeq F_2$ (and dual).

Solution. Both $F_1(X)$ and $F_2(X)$ represent the same functor $X' \mapsto \text{Hom}(X, G(X'))$ and there exists therefore an isomorphism $F_1(X) \simeq F_2(X)$ which is easily seen to be natural. ■

Exercise 1.4.7 Show that if $F_1 : \mathcal{C} \rightleftarrows \mathcal{C}' : G_1$ and $F_2 : \mathcal{C}' \rightleftarrows \mathcal{C}'' : G_2$ then, $F_2 \circ F_1 : \mathcal{C} \rightleftarrows \mathcal{C}'' : G_1 \circ G_2$.

1.4.2 Unit and counit

Proposition 1.4.8 A functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is adjoint to $G : \mathcal{C}' \rightarrow \mathcal{C}$ if and only if there exists $\alpha : \text{Id}_{\mathcal{C}} \Rightarrow G \circ F$ and $\beta : F \circ G \Rightarrow \text{Id}_{\mathcal{C}'}$ such that $\beta_F \circ F(\alpha) = \text{Id}_F$ and $G(\beta) \circ \alpha_G = \text{Id}_G$.

The condition means that the diagrams

$$\begin{array}{c} \circlearrowleft \\ F(X) \xrightleftharpoons[\beta_{F(X)}]{F(\alpha_X)} F(G(F(X))) \end{array} \text{ and } \begin{array}{c} G(F(G(X'))) \xrightleftharpoons[\alpha_{G(X')}]^{G(\beta_{X'})} G(X') \circlearrowright \end{array}$$

are always commutative.

Proof. Assume first that there exists a natural isomorphism

$$\Phi_{X,X'} : \text{Hom}(F(X), X') \simeq \text{Hom}(X, G(X')).$$

We may then set

$$\alpha_X := \Phi_{X,F(X)}(\text{Id}_{F(X)}) \text{ and } \beta_{X'} := \Phi_{G(X'),X'}^{-1}(\text{Id}_{G(X')}).$$

Then, we have

$$\begin{aligned} \beta_{F(X)} \circ F(\alpha_X) &= \Phi_{G(F(X)),F(X)}^{-1}(\text{Id}_{G(F(X))}) \circ F(\Phi_{X,F(X)}(\text{Id}_{F(X)})) \\ &= \Phi_{X,F(X)}^{-1}(\Phi_{X,F(X)}(\text{Id}_{F(X)})) \\ &= \text{Id}_{F(X)} \end{aligned}$$

and dually.

Conversely, given $f : F(X) \rightarrow X'$, we set $\Phi_{X,X'}(f) = G(f) \circ \alpha_X$ and define dually for $f' : X \rightarrow G(X')$, $\Psi_{X,X'}(f') := \beta_{X'} \circ F(f')$. We shall then have

$$(\Psi_{X,X'} \circ \Phi_{X,X'})(f) = \beta_{X'} \circ F(G(f)) \circ F(\alpha_X) = f \circ \beta_{F(X)} \circ F(\alpha_X) = f$$

and symmetrically by duality. ■

Definition 1.4.9 The natural transformations α and β are then called *unit* and *counit* for the adjunction.

- Examples**
1. In the case of topological spaces and sets, for the discrete/forget (resp. forget/coarse) adjunction, we get the identity for both unit and counit, the counit (resp. unit) being the canonical continuous map $X^{\text{disc}} \rightarrow X$ (resp. $X \rightarrow X^{\text{coarse}}$).
 2. In the case of abelian groups and sets, for the free/forget adjunction, we have unit $E \mapsto \mathbb{Z} \cdot E, x \mapsto 1 \cdot x$ and counit $\mathbb{Z} \cdot M \twoheadrightarrow M, \sum a_i \cdot x_i \mapsto \sum a_i x_i$.
 3. In the case of groups and monoids, for the gr/forget (resp. forget/ \times) adjunction, we get inclusion $G \mapsto G^{\text{gr}}$ and identity (resp. identity and inclusion $G \mapsto G^\times$) for unit and counit respectively.

Proposition 1.4.10 If F is adjoint to G with unit α and counit β , then F is faithful (resp. fully faithful) if and only if α_X is always a monomorphism (resp. an isomorphism). And dually.

Proof. For $X, Y \in \mathcal{C}$, there exists a commutative diagram

$$\begin{array}{ccc} & \text{Hom}(Y, X) & \\ F \swarrow & & \searrow \alpha_{X*} \\ \text{Hom}(F(Y), F(X)) & \xrightarrow{\sim} & \text{Hom}(Y, G(F(X))) \end{array}$$

This follows from the fact that α is a natural transformation so that

$$\Phi_{Y, F(X)}(F(f)) = G(F(f)) \circ \alpha_Y = \alpha_X \circ f.$$

Thus we see that the map F is injective (resp. bijective) for all X, Y in \mathcal{C} if and only if this is the case for α_{X*} , which means that α_X is a monomorphism (resp. an isomorphism) for all $X \in \mathcal{C}$.

Now, we have G^{op} is adjoint to F^{op} and the unit for this adjunction is β^{op} . Moreover, $\beta_{X'}^{\text{op}}$ is a monomorphism (resp. an isomorphism) if and only if $\beta_{X'}$ is an epimorphism (resp. an isomorphism). Therefore, G is faithful (resp. fully faithful) if and only if G^{op} is faithful (resp. fully faithful) if and only if $\beta_{X'}$ is an epimorphism (resp. an isomorphism) for all X' . ■

Exercise 1.4.11 Describe unit and counit in all the examples studied so far. Deduce in each case faithfulness or full faithfulness of the functors.

Exercise 1.4.12 Show that, if a small category \mathcal{C} has (self) coproducts, then all representable functors F on \mathcal{C} have an adjoint.

Hint. We may assume that $F = h^X$, consider the functor

$$\text{Set} \rightarrow \mathcal{C}, \quad I \mapsto X^{(I)} = \coprod_I X$$

define unit $I \rightarrow \text{Hom}(X, X^{(I)})$ and counit $X^{(\text{Hom}(X, Y))} \rightarrow Y$ and check the properties. ■

1.4.3 Adjoint and limit

Proposition 1.4.13 All limits of shape I exist in a category \mathcal{C} if and only if the functor $X \mapsto \underline{X}$ has a coadjoint which is then given by^a $D \mapsto \varprojlim D$.

^aOnce a specific choice of the limit is made for each diagram D .

Proof. We have indeed a natural isomorphism in X given by

$$\text{Hom}(\underline{X}, D) \simeq \text{Hom}(X, \varprojlim D)$$

and it follows from exercise 1.3.6 that it is also natural in D . ■

Exercise 1.4.14 Show that any adjunction between two functors F and G extends to an adjunction on diagrams of a given shape I :

$$\text{Hom}(F(D), E) \simeq \text{Hom}(D, G(E)).$$

Corollary 1.4.15 A functor that has an adjoint preserves all limits (and dual).

Proof. We have to show that, if G has an adjoint F and D is a diagram, then $G(\varprojlim D) \simeq \varprojlim G(D)$. This follows by adjunction from $F(\underline{X}) = \underline{F(X)}$. ■

As a consequence, we recover the result from proposition 1.3.39: limits commute with limits (and dual).

Example As we already noticed, the forgetful functor $\mathbf{Top} \rightarrow \mathbf{Set}$ preserves all limits and all colimits, the forgetful functor $\mathbf{Ab} \rightarrow \mathbf{Set}$ preserves all limits and the inclusion functor $\mathbf{Mon} \hookrightarrow \mathbf{Grp}$ preserves all limits and colimits.

Theorem 1.4.15 has a partial converse:

Theorem 1.4.16 — Freyd adjunction theorem. If \mathcal{C}' be a complete category and $G : \mathcal{C}' \rightarrow \mathcal{C}$ a functor that preserves all limits and satisfies the *solution-set condition* below, then it has an adjoint (and dual).

Solution-set condition: given any X in \mathcal{C} , there exists a *set* of morphisms $X \rightarrow G(Y_i)$ such that any morphism $X \rightarrow G(Y)$ factors through some $G(Y_i)$.

Proof. (Sketch) It is sufficient to set

$$F(X) := \varprojlim_{X \rightarrow G(Y_i)} Y_i. \quad \blacksquare$$

Theorem 1.4.17 Assume $F : \mathcal{C} \rightarrow \mathcal{C}'$ is adjoint to a fully faithful functor G . If D' is a diagram in \mathcal{C}' and $X = \varprojlim G(D')$, then $X' := F(X) = \varprojlim D'$ and $X \simeq G(X')$.

Proof. To do (corollary 5.6.6 of [Riehl16]). ■

1.4.4 Reflective subcategory

Definition 1.4.18 A full subcategory $\mathcal{C}' \subset \mathcal{C}$ is said to be *reflective* if the inclusion functor has an adjoint, called *reflection*. The dual terminology is *coreflective*.

In other words, a functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is a reflection if there exists a natural bijection

$$\mathrm{Hom}(F(X), X') \simeq \mathrm{Hom}(X, X')$$

when $X' \in \mathcal{C}'$. It means that there exists a natural morphism $X \rightarrow F(X)$ with the following universal property:

$$\begin{array}{ccc} X & \xrightarrow{\forall f} & X' \\ \downarrow & \searrow \exists! \bar{f} & \\ F(X) & & \end{array}$$

- Examples**
1. The category of preordered sets is a reflective subcategory of the category of partially ordered sets.
 2. The category of abelian groups is a reflective subcategory of the category of all groups.

3. The category of groups is both a reflective and a coreflective subcategory of the category of all monoids.
4. The category of compact Hausdorff spaces is a reflective subcategory of the category of all topological spaces.
5. The category of sets is isomorphic to the reflective (resp. coreflective) subcategory of discrete (resp. chaotic) topological spaces.

Exercise 1.4.19 Show that if \mathcal{C}' is a reflective subcategory of \mathcal{C} , then any diagram D' in \mathcal{C}' that has a limit (resp. a colimit) in \mathcal{C} has a limit (resp. a colimit) in \mathcal{C}' . Show also that \mathcal{C}' is stable under limits that exist in \mathcal{C} . Finally, show that, if a colimit in \mathcal{C} of objects of \mathcal{C}' is an object of \mathcal{C}' , then this is a colimit in \mathcal{C}' .

Exercise 1.4.20 Show that, if \mathcal{C}' is a full subcategory of \mathcal{C} , then a functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is a reflection if and only if there exists a natural morphism $\alpha_X : X \rightarrow F(X)$ for $X \in \mathcal{C}$ such that α_X is an isomorphism when $X \in \mathcal{C}'$.

Solution. The subcategory \mathcal{C}' is reflective with reflection F if and only if there exists a natural morphism $\alpha_X : X \rightarrow F(X)$ for $X \in \mathcal{C}$ and a natural isomorphism $\beta_X : F(X) \simeq X$ when $X \in \mathcal{C}'$ such that $\beta_{F(X)} \circ F(\alpha_X) = \text{Id}_{F(X)}$ when $X \in \mathcal{C}$ and $\beta_X \circ \alpha_X = \text{id}_X$ when $X \in \mathcal{C}'$. This is clearly equivalent to our condition with $\beta_X := \alpha_X^{-1}$. ■

Exercise 1.4.21 Show that, if filtered colimits exist in \mathcal{C} , then \mathcal{C} is a reflective subcategory of $\text{Ind}(\mathcal{C})$ with adjoint “ $\varinjlim X_i \mapsto \varinjlim X_i$ ”.

1.4.5 Kan extension

Definition 1.4.22 Let $p : \mathcal{C} \rightarrow \mathcal{C}'$ be a functor between small categories. The *(left) Kan extension* of a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ along p is a functor $p!F$ which is universal for all functors $G : \mathcal{C}' \rightarrow \mathcal{D}$ and natural transformations $F \Rightarrow p^{-1}G := G \circ p$.

In other words, $p!F$ represents the functor $G \mapsto \text{Hom}(F, G \circ p)$ on the category $\mathbf{Hom}(\mathcal{C}', \mathcal{D})$. It means that $p!F : \mathcal{C}' \rightarrow \mathcal{D}$ is endowed with a natural transformation $\alpha : F \Rightarrow p^{-1}p!F$ such that, given any natural transformation $\gamma : F \Rightarrow p^{-1}G$, there exists a unique natural transformation $\tilde{\gamma} : p!F \Rightarrow G$ such that $\gamma = p^{-1}(\tilde{\gamma}) \circ \alpha$:

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{F} & \mathcal{D} \\ \downarrow p & \nearrow p!F \Rightarrow G & \\ \mathcal{C}' & & \end{array} \quad \text{and} \quad \begin{array}{ccc} F & \xRightarrow{\alpha} & p^{-1}G \\ \downarrow & \nearrow & \\ p^{-1}p!F & & \end{array}$$

(not commutative)

There exists the dual notion of a *right Kan extension* p_*F with $\gamma : p^{-1}G \Rightarrow F$ this time.

Examples 1. A diagram $D : I \rightarrow \mathcal{C}$ has colimit X if and only if the constant functor $\mathbf{1} \rightarrow \mathcal{C}, 0 \mapsto X$ is the Kan extension of D along the projection $I \rightarrow \mathbf{1}$:

$$\begin{array}{ccc} I & \xrightarrow{D} & \mathcal{C} \\ \downarrow & \nearrow X \rightarrow Y & \\ \mathbf{1} & & \end{array} \quad \text{and} \quad \begin{array}{ccc} D & \longrightarrow & \underline{Y} \\ \downarrow & \nearrow & \\ \underline{X} & & \end{array}$$

2. A functor $F : \mathcal{C} \rightarrow \mathcal{D}$ between small categories has a coadjoint G if and only if the Kan extension of $\text{Id}_{\mathcal{C}}$ along F exists and $F \circ F! \text{Id}_{\mathcal{C}} = F!F$, in which case $G = F! \text{Id}_{\mathcal{C}}$:

$$\begin{array}{ccc} \mathcal{C} & \xRightarrow{\text{Id}_{\mathcal{C}}} & \mathcal{C} \xrightarrow{F} \mathcal{D} \\ F \downarrow & \nearrow G & \nearrow F \circ G \\ \mathcal{D} & & \end{array}, \quad \begin{array}{ccc} \text{Id}_{\mathcal{C}} & \xRightarrow{\alpha} & G' \circ F \\ \downarrow & \nearrow & \\ G \circ F & & \end{array} \quad \text{and} \quad \begin{array}{ccc} F & \xRightarrow{\alpha} & H \circ F \\ \downarrow & \nearrow & \\ F \circ G \circ F & & \end{array}$$

Proposition 1.4.23 Let $p : \mathcal{C} \rightarrow \mathcal{C}'$ be functor between small categories. Then the functor

$$p^{-1} : \mathbf{Hom}(\mathcal{C}', \mathcal{D}) \rightarrow \mathbf{Hom}(\mathcal{C}, \mathcal{D}), \quad G \mapsto G \circ p.$$

has an adjoint $p_!$ if and only if all Kan extensions along p with values in \mathcal{D} exist (and dual).

Proof. Follows immediately from the definition. ■

Proposition 1.4.24 If all colimits exist in \mathcal{D} , then the Kan extension of $F : \mathcal{C} \rightarrow \mathcal{D}$ along $p : \mathcal{C} \rightarrow \mathcal{C}'$ always exists (and dual).

Proof. (Sketch) We set

$$(p_! F)(X') := \varinjlim_{p(X) \rightarrow X'} F(X)$$

and check. ■

Exercise 1.4.25 Show that if \mathcal{C} is small, all colimits exist in \mathcal{D} and $F : \mathcal{C} \rightarrow \mathcal{D}$ is fully faithful, then (the above diagram commutes:) $F \simeq p_! F \circ p$ (and dual).

Exercise 1.4.26 Show that, if $g : \mathcal{C} \rightarrow \mathcal{C}'$ is a functor between small categories and all colimits exist in \mathcal{D} , then the functor

$$\widehat{g}^{-1} : \widehat{\mathcal{C}'}(\mathcal{D}) \rightarrow \widehat{\mathcal{C}}(\mathcal{D}), \quad T' \mapsto T' \circ g.$$

has an adjoint $\widehat{g}_!$ (resp. a coadjoint $\widehat{g}_* : \widehat{\mathcal{C}} \rightarrow \widehat{\mathcal{C}'}$).

Solution. Follows from propositions 1.4.23 and 1.4.24. ■

1.5 Miscellaneous

1.5.1 Algebraic structure

A category is said to be *cartesian*⁷ if all finite products exist.

Definition 1.5.1 Let \mathcal{C} be a *cartesian* category. A *monoid* of \mathcal{C} is an object G endowed with a *multiplication morphism* $\mu : G \times G \rightarrow G$ and a *unit morphism*^a $\epsilon : 1 \rightarrow G$ making commutative the following diagrams:

$$\begin{array}{ccc} G \times G \times G & \xrightarrow{\mu \times \text{Id}_G} & G \times G \\ \downarrow \text{Id}_G \times \mu & & \downarrow \mu \\ G \times G & \xrightarrow{\mu} & G \end{array} \quad \text{and} \quad \begin{array}{ccc} G & \xrightarrow{\epsilon \times \text{Id}_G} & G \times G \\ \downarrow \text{Id}_G \times \epsilon & \searrow & \downarrow \mu \\ G \times G & \xrightarrow{\mu} & G \end{array}$$

(under the identification $1 \times G \simeq G \simeq G \times 1$). A *morphism of monoids* $G \rightarrow G'$ of \mathcal{C} is a morphism $f : G \rightarrow G'$ in \mathcal{C} making the following commutative:

$$\begin{array}{ccc} G \times G & \xrightarrow{f \times f'} & G' \times G' \\ \downarrow \mu & & \downarrow \mu' \\ G & \xrightarrow{f} & G' \end{array} \quad \text{and} \quad \begin{array}{ccc} G & \xrightarrow{f} & G' \\ \epsilon \swarrow & & \searrow \epsilon' \\ & 1 & \end{array}$$

⁷This is a particular case of the notion of a *monoidal category*.

It is a *group* if there exists an *inversion morphism* $\iota : G \rightarrow G$ making commutative

$$\begin{array}{ccc}
 G & \xrightarrow{(\iota, \text{Id}_G)} & G \times G \\
 \downarrow (\text{Id}_G, \iota) & \searrow & \downarrow \mu \\
 & 1 & \\
 G \times G & \xrightarrow{\mu} & G
 \end{array}$$

ϵ

It is *abelian* if $G^{\text{op}} = G$ where G^{op} denotes the monoid obtained by exchanging factors in $G \times G$ (composing μ with the flip map).

^aAutomatically unique.

Monoids (resp. groups, resp. abelian groups) of \mathcal{C} make a category $\mathbf{Mon}(\mathcal{C})$ (resp. $\mathbf{Grp}(\mathcal{C})$, resp. $\mathbf{Ab}(\mathcal{C})$). We shall mostly concentrate on abelian groups.

- Examples**
1. An group of the category \mathbf{Set} is nothing but a usual group.
 2. An group of \mathbf{Top} is a topological group (with continuous multiplication and continuous inversion).
 3. A group in the category of smooth manifolds is a *Lie group*.
 4. An abelian group in the category of abelian groups is an abelian group (!).

Exercise 1.5.2 Show that $k[t]$ (resp. $k[t]_t$) endowed with $t \mapsto t \otimes 1 + 1 \otimes t$ (resp. $t \mapsto t \otimes t$) is an abelian group of the category opposite to the category of k -algebras (it is called a *bialgebra*).

Exercise 1.5.3 Show that, if all limits exist in \mathcal{C} , then the same holds in $\mathbf{Ab}(\mathcal{C})$ and they are preserved by the forgetful functor $\mathbf{Ab}(\mathcal{C}) \rightarrow \mathcal{C}$.

Exercise 1.5.4 Let $F : \mathcal{C} \rightarrow \mathcal{C}'$ be a functor between cartesian categories. Show that

1. if F preserves finite products, then F induces a functor still written $F : \mathbf{Ab}(\mathcal{C}) \rightarrow \mathbf{Ab}(\mathcal{C}')$,
2. if moreover, F preserves all limits of \mathcal{C} , then it preserves all limits of $\mathbf{Ab}(\mathcal{C})$.

The category of abelian groups in a category of presheaves of sets can be identified with the category of presheaves of abelian groups:

Exercise 1.5.5 Show that if \mathcal{C} is a small category, then there exists an isomorphism of categories $\widehat{\mathcal{C}}(\mathbf{Ab}) \simeq \mathbf{Ab}(\widehat{\mathcal{C}})$.

Solution. For fixed $X \in \mathcal{C}$, the functor $\widehat{\mathcal{C}} \rightarrow \mathbf{Set}, F \mapsto F(X)$ preserves finite products (actually all limits and colimits: see for example exercise 1.3.45 below). Therefore, if $\mathcal{M} \in \mathbf{Ab}(\widehat{\mathcal{C}})$ and $X \in \mathcal{C}$, then $\mathcal{M}(X)$ is a usual abelian group and this is clearly functorial. Conversely, if $\mathcal{M} \in \widehat{\mathcal{C}}(\mathbf{Ab})$, then the obvious family of maps

$$\mu_X : \mathcal{M}(X) \times \mathcal{M}(X) \rightarrow \mathcal{M}(X), \quad \epsilon_X : \{0\} \rightarrow \mathcal{M}(X) \quad \text{and} \quad \iota_X : \mathcal{M}(X) \rightarrow \mathcal{M}(X)$$

define a structure of abelian group on the underlying presheaf of sets of \mathcal{M} . These two constructions are clearly inverse to each other. ■

Exercise 1.5.6 Show that if \mathcal{C} is a small cartesian category, then there exists a fully faithful functor $\mathbf{Ab}(\mathcal{C}) \hookrightarrow \widehat{\mathcal{C}}(\mathbf{Ab})$.

Definition 1.5.7 Let \mathcal{C} be a cartesian category.

1. Let (G, μ) be a monoid of \mathcal{C} with unit ϵ . An *action* of G on an object E of \mathcal{C} is a *morphism* $\rho : G \times E \rightarrow E$ making commutative the following diagrams:

$$\begin{array}{ccc} G \times G \times E & \xrightarrow{\mu \times \text{Id}_E} & G \times E \\ \downarrow \text{Id}_G \times \rho & & \downarrow \mu \\ G \times E & \xrightarrow{\mu} & G \end{array} \quad \text{and} \quad \begin{array}{ccc} E & \xrightarrow{\epsilon \times \text{Id}_G} & G \times E \\ & \searrow & \downarrow \mu \\ & & E. \end{array}$$

2. A *ring* of \mathcal{C} is an object A endowed with a structure of abelian group (given by μ) and a structure of monoid (given by ν) making commutative

$$\begin{array}{ccc} A \times A \times A & \xrightarrow{\text{Id}_A \times \mu} & A \times A \\ \downarrow & & \searrow \nu \\ A \times A \times A \times A & \xrightarrow{\nu \times \nu} & A \times A \end{array} \quad \begin{array}{c} \nearrow \mu \\ \nearrow \nu \\ \nearrow \mu \end{array} \rightarrow A$$

with left distribution on the left as well as the analogous diagram for right distribution:

$$\begin{array}{ccc} A \times A \times A & \xrightarrow{\mu \times \text{Id}_M} & A \times A \\ \downarrow & & \searrow \nu \\ A \times A \times A \times A & \xrightarrow{\nu \times \nu} & A \times A \end{array} \quad \begin{array}{c} \nearrow \nu \\ \nearrow \mu \\ \nearrow \nu \end{array} \rightarrow A.$$

3. Let (A, μ, ν) be a ring of \mathcal{C} . An *A-module* is an abelian group (M, μ) of \mathcal{C} endowed with an action ρ of (A, ν) making commutative both

$$\begin{array}{ccc} A \times M \times M & \xrightarrow{\text{Id}_A \times \mu} & A \times M \\ \downarrow & & \searrow \rho \\ A \times M \times A \times M & \xrightarrow{\rho \times \rho} & M \times M \end{array} \quad \begin{array}{c} \nearrow \rho \\ \nearrow \mu \\ \nearrow \rho \end{array} \rightarrow M$$

and

$$\begin{array}{ccc} A \times A \times M & \xrightarrow{\mu \times \text{Id}_M} & A \times M \\ \downarrow & & \searrow \rho \\ A \times M \times A \times M & \xrightarrow{\rho \times \rho} & M \times M \end{array} \quad \begin{array}{c} \nearrow \rho \\ \nearrow \mu \\ \nearrow \rho \end{array} \rightarrow M.$$

They form a category $G\text{-Set}(\mathcal{C})$ (resp. $\mathbf{Ring}(\mathcal{C})$, resp. $A\text{-Mod}(\mathcal{C})$).

Example We consider the category $\mathcal{C}^k\text{-Man}$ of \mathcal{C}^k -differentiable manifolds and \mathcal{C}^k -differentiable maps (works with $k = 0, 1, \dots, \infty$, and even hol). A group in this category is a \mathcal{C}^k -differentiable manifold G endowed with a group structure such that both multiplication $\mu : G \times G \rightarrow G$ and inversion $i : G \rightarrow G$ are \mathcal{C}^k -differentiable. A G^{op} -object is a \mathcal{C}^k -differentiable manifold X endowed with a right

action $X \times G \rightarrow X$ which is \mathcal{C}^k -differentiable. For example, the *trivial action* is the (second) projection. Assume that X is endowed with the *trivial action* and consider the category $G^{\text{op}}\text{-}\mathcal{C}^k\text{-}\mathbf{Man}/_X$ of G -objects $\pi : Y \rightarrow X$ over X . Such an object restricts any open subset U of X by setting $Y|_U := \pi^{-1}(U)$ and considering $Y|_U \rightarrow U$ together with $Y|_U \times G \rightarrow Y|_U$. Such an object is said to be *trivial* if it is isomorphic to the first projection $X \times G \rightarrow X$ endowed with $\text{Id}_X \times \mu : X \times G \times G \mapsto X \times G$. It is called a *G-torsor* or a *principal G-bundle* if it is *locally trivial*: there exists a covering of X by open subsets U such that $Y|_U$ is trivial. We can consider the full subcategory of all G -torsors over X . It is an example of a *groupoid*: any morphism is an isomorphism.

Exercise 1.5.8 Define the category $k\text{-}\mathbf{Alg}(\mathcal{C})$ of k -algebras of a cartesian category \mathcal{C} when k is a commutative ring of \mathcal{C} .

Exercise 1.5.9 Show that $k[t]$ endowed with $t \mapsto t \otimes 1 + 1 \otimes t$ and $t \mapsto t \otimes t$ is a commutative ring of the category opposite to the category of k -algebras.

1.5.2 Projective/Injective

Definition 1.5.10 An object X of a category \mathcal{C} is said to be *projective* if h^X preserves epimorphisms.

The dual notion is that of an *injective* object: h_X sends monomorphisms to epimorphisms. Thus X is projective (resp. injective) if whenever $Z \twoheadrightarrow Y$ is an epimorphism (resp. $Y \hookrightarrow Z$ is a monomorphism), then the corresponding map

$$\text{Hom}(X, Z) \twoheadrightarrow \text{Hom}(X, Y) \quad (\text{resp. } \text{Hom}(Z, X) \twoheadrightarrow \text{Hom}(Y, X))$$

is surjective. In other words, X is projective (resp. injective) when any diagram

$$\begin{array}{ccc} & Z & \\ \nearrow & \downarrow & \\ X & \longrightarrow & Y \end{array} \quad (\text{resp.} \quad \begin{array}{ccc} & Z & \\ \uparrow & \searrow & \\ Y & \longrightarrow & X \end{array})$$

can be completed with the dotted arrow.

- Examples**
1. In **Set**, all objects (resp. non empty objects) are projective (resp. injective).
 2. In **Ab**, projective objects are free abelian groups and injective objects are divisible groups (proposition 1.5.14 below).
 3. In $R\text{-}\mathbf{Mod}$, projective objects are direct factors of free R -modules ($\mathbb{Z}/2$ is projective – but not free – over $\mathbb{Z}/6$).
 4. If E is a fiber bundle on a compact manifold X , then $E(X)$ is a projective $\mathcal{O}(X)$ -module (requires an argument).
 5. The projective objects of the category of compact Hausdorff spaces are the *Stonean* (meaning extremally disconnected⁸) compact Hausdorff spaces.

⁸Closure of open is open.

Exercise 1.5.11 Show that a coproduct of projectives is projective (and dual).

Exercise 1.5.12 Show that if X is projective, then

1. any epimorphism $Y \twoheadrightarrow X$ is split,
2. if a $Y \rightarrowtail X$ (resp. $X \twoheadrightarrow Y$) is a split monomorphism (resp epimorphism), then Y also is projective (and dual).

Proposition 1.5.13 When epimorphisms are stable under pullback, an object X is projective if and only if any epimorphism $Y \twoheadrightarrow X$ is split (and dual).

Proof. Pulling back an epimorphism $Z \twoheadrightarrow Y$ along a morphism $X \rightarrow Y$ provides an epimorphism $X \times_Y Z \twoheadrightarrow X$ that has a section $X \rightarrow X \times_Y Z$ that we can compose with the projection $X \times_Y Z \rightarrow Z$ to get a lifting $X \rightarrow Z$ of the original map $X \rightarrow Y$ along the epimorphism $Z \twoheadrightarrow Y$:

$$\begin{array}{ccc} X \times_Y Z & \longrightarrow & Z \\ \downarrow & \nearrow & \downarrow \\ X & \longrightarrow & Y \end{array}$$

■

Proposition 1.5.14 An abelian group Q is injective if and only if it is divisible.

Proof. Assume Q injective. If $a \in Q$ and $n > 0$, then the map $\mathbb{Z} \rightarrow Q, 1 \mapsto a$ extends along the injective map $\mathbb{Z} \xrightarrow{n} \mathbb{Z}$ to a map $1 \mapsto b$ with $b = na$. Contradiction.

Assume conversely that Q is divisible and we are given a morphism $f : M \rightarrow Q$ and an injective map $M \rightarrowtail N$. By Zorn's lemma, there exists a maximal subgroup M' of N such that f extends to M' . We may assume $M' = M$. If $a \notin M$, then $M \cap \langle a \rangle = \langle na \rangle$. In Q , we can write $f(na) = nx$ and extend f to $M + \langle a \rangle$ by sending a to x (works also for $n = 0$). ■

Definition 1.5.15 A category \mathcal{C} is said to have *enough projectives* if, given any $X \in \mathcal{C}$, there exists a projective Y and an epimorphism $Y \twoheadrightarrow X$ (and dual).

Thus \mathcal{C} has *enough injectives* if there always exists a monomorphism $X \rightarrowtail Y$ with Y injective.

- Examples**
1. The category **Set** has enough projectives and injectives.
 2. The category $R\text{-Mod}$ has enough projectives and injectives (we shall prove this last statement later).
 3. The category of compact Hausdorff spaces has enough projectives.

Exercise 1.5.16 Show that

1. if each X_i is projective, then $(X_i)_{i \in I}$ is projective in $\mathcal{C} := \prod_{i \in I} \mathcal{C}_i$,
2. if each \mathcal{C}_i has enough projectives, so does \mathcal{C} .

We shall extend later the next statement to A -modules.

Proposition 1.5.17 The category of abelian groups has enough injectives.

Proof. We write $M = F/N$ with F free and consider the embedding $M \hookrightarrow F_{\mathbb{Q}}/N$ with $F_{\mathbb{Q}} := \mathbb{Q} \otimes_{\mathbb{Z}} F$ (\mathbb{Q} -vector space is divisible and quotient of divisible is divisible). ■

Definition 1.5.18 An object X of a category \mathcal{C} is said to be *compactly presented* (or *compact*) if h^X preserves filtered colimits.

In other words,

$$\varinjlim \operatorname{Hom}(X, Y_i) \simeq \operatorname{Hom}(X, \varinjlim Y_i)$$

when (Y_i) is filtered. It means that any morphism $X \rightarrow \varinjlim Y_i$ factors through some Y_i .

Examples

1. A set is compactly presented if and only if it is finite.
2. A topological space is compactly presented if and only if it is finite discrete.
3. An abelian group is compactly presented (resp. compactly presented projective) if and only if it is finitely generated (resp. free of finite rank).
4. An open subset U of a topological space X is compact (in the usual sense) if and only if U is a compactly presented object of $\operatorname{Open}(X)$.

Definition 1.5.19 A set $S \subset \mathcal{C}$ is a *set of generators* (or *separators*) for a category \mathcal{C} if the functor $\prod_{G \in S} h^G$ is faithful^a. When $S = \{G\}$, we say that G is a *generator*.

^aOne should say separators (resp. generators) when the functor is faithful (resp. conservative)

In down to earth terms, it means that if $f_1 \neq f_2 : X \rightarrow Y$, then there exists $g : G \rightarrow X$ with $G \in S$ such that $f_1 \circ g \neq f_2 \circ g$.

Examples

1. $1 := \{0\}$ is a generator for \mathbf{Set} .
2. A is a generator for $A\text{-Mod}$.
3. If \mathcal{C} is a category, then $\{h_X, X \in \mathcal{C}\}$ is a set of generators for $\widehat{\mathcal{C}}$.

Exercise 1.5.20 Show that, if \mathcal{C} has all coproducts, then S is a set of generators if and only if there exists for all $X \in \mathcal{C}$ an epimorphism $\coprod_{i \in I} G_i \twoheadrightarrow X$ with $G_i \in S$.

Proof. For the converse, choose

$$\coprod_{G \in S, f: G \rightarrow X} G \twoheadrightarrow X. \quad \blacksquare$$

1.5.3 Localization

In this sketchy section, we sweep set theoretical questions under the rug.

Definition 1.5.21 The *localization* of a (small) category \mathcal{C} with respect to a collection of morphisms W of \mathcal{C} is a category $W^{-1}\mathcal{C}$ which is universal for functors $\mathcal{C} \rightarrow \mathcal{D}$ sending W to isomorphisms in \mathcal{D} .

It means that there exists a functor $\gamma : \mathcal{C} \rightarrow W^{-1}\mathcal{C}$ sending W to isomorphisms such that, given any category \mathcal{D} , the functor

$$\gamma^* : \operatorname{Hom}(W^{-1}\mathcal{C}, \mathcal{D}) \rightarrow \operatorname{Hom}(\mathcal{C}, \mathcal{D})$$

induces an equivalence with the full subcategory of functors sending W to isomorphisms.

Exercise 1.5.22 An *isogeny* is a homomorphism of abelian groups $f : M \rightarrow N$ such that

1. $\forall y \in N, \forall n \in \mathbb{Z} \setminus 0, \exists x \in M, f(nx) = y,$
2. $\forall x \in M, f(x) = 0 \Rightarrow \exists n \in \mathbb{Z} \setminus 0, nx = 0.$

Show that, if W is the collection of all isogenies, then $W^{-1}\mathbf{Ab} \simeq \mathbf{Q}\text{-Vec} (:= \mathbf{Q}\text{-Mod}).$

Exercise 1.5.23 The category of topological spaces up to homotopy has topological spaces as objects but sets of morphisms $[X, Y] := \mathcal{C}(X, Y) / \sim$ and composition induced by usual composition. Show that it is equivalent to $W^{-1}\mathbf{Top}$ if W denotes the collection of homotopy equivalences.

Proposition 1.5.24 The localization $W^{-1}\mathcal{C}$ of a (small) category \mathcal{C} with respect to a set of morphisms W always exists.

Proof. (Sketch) We may assume that W contains all isomorphisms in \mathcal{C} . Then, the objects of $W^{-1}\mathcal{C}$ are the objects of \mathcal{C} and morphisms⁹ are finite chains

$$X = X_0 \xleftarrow{W} X_1 \rightarrow X_2 \xleftarrow{W} \cdots \rightarrow X_{n-1} \xleftarrow{W} X_n \rightarrow X_{n+1} = Y$$

up to some equivalence. ■

Definition 1.5.25 A category \mathcal{C} admits *right calculus of fractions* with respect to a set of morphisms W if

1. W contains all identities and is stable under composition,
2. given any $f : X \rightarrow Y$ in \mathcal{C} and $\varphi : Y' \rightarrow Y$ in W , there always exists $f' : X' \rightarrow Y'$ and $\varphi' : X' \rightarrow X$ in W with $\varphi \circ f' = f \circ \varphi'$,
3. given any $f, g : X \rightarrow Y'$ in \mathcal{C} and $\psi : Y' \rightarrow Y$ in W such that $\psi \circ f = \psi \circ g$, there exists $\varphi : X' \rightarrow X$ in W such that $f \circ \varphi = g \circ \varphi$.

Proposition 1.5.26 If a (small) category \mathcal{C} admits right calculus of fraction with respect to W , then $W^{-1}\mathcal{C}$ is the category having the same objects as \mathcal{C} and

$$\mathrm{Hom}_{W^{-1}\mathcal{C}}(X, Y) = \varinjlim_{X' \rightarrow X \in W} \mathrm{Hom}_{\mathcal{C}}(X', Y).$$

Proof. (Sketch) By definition, morphisms and composition are described, up to equivalence, by the following diagram

$$\begin{array}{ccccc} X'' & & & & \\ \downarrow W & \searrow & & & \\ X' & & Y' & & \\ \downarrow W & \searrow & \downarrow W & \searrow & \\ X & \cdots \cdots \cdots & Y & \cdots \cdots \cdots & Z. \end{array}$$

It is then a matter of checking the various properties. ■

⁹They form a collection and not merely a set.

Exercise 1.5.27 Show that if \mathcal{C} admits right calculus of fraction with respect to W , then the localization functor $Q : \mathcal{C} \rightarrow W^{-1}\mathcal{C}$ is exact.

2. Linear algebra

In the same way as general categories are modeled on the category of sets, there exist the notion of an abelian category that is modeled on the category of abelian groups.

2.1 Additive structure

2.1.1 Preadditive category

Definition 2.1.1 A *pre-additive category* (also called **Ab-category**^a or *ringoid*) is a category \mathcal{P} endowed with a factorization of the Hom functor:

$$\begin{array}{ccc} \text{Hom} : \mathcal{P}^{\text{op}} \times \mathcal{P} & \longrightarrow & \mathbf{Set} \\ & \searrow & \uparrow \\ & & \mathbf{Ab}. \end{array}$$

^aThis is a particular instance of the notion of an *enriched* category.

In other words, \mathcal{P} is preadditive if and only if, for all $M, N \in \mathcal{P}$, $\text{Hom}(M, N)$ is (endowed with the structure of) an abelian group and for all $M, N, P \in \mathcal{P}$, the map

$$\begin{array}{ccc} \text{Hom}(M, N) \times \text{Hom}(N, P) & \longrightarrow & \text{Hom}(M, P) \\ (f, g) & \longmapsto & g \circ f \end{array}$$

is bilinear (check). In other words, distributivity holds. We shall denote by $0_{MN} : M \rightarrow N$ (or $0_M : M \rightarrow M$) for the zero morphism but we may as well write $0 := 0_{MN}$. If $M \in \mathcal{P}$, then $\text{End}(M)$ is now a ring and not merely a monoid and we may as well write $1 := 1_M := \text{Id}_M$ (which should not be confused with the final morphism).

A preadditive category is not a usual category satisfying some specific property but a usual category *endowed* with some extra structure. Unlike the case of an abelian group versus a set for instance, there is no practical reason to distinguish between the two since, as we shall see, this additional structure is unique when the category is cartesian.

- Examples**
1. The category **Ab** of abelian groups is preadditive.
 2. More generally, if A is a ring, then the category $A\text{-Mod}$ of left A -modules is preadditive.
 3. If \mathcal{C} is a small category, then the category $\widehat{\mathcal{C}}(\mathbf{Ab})$ of presheaves of abelian groups on \mathcal{C} is preadditive.
 4. If \mathcal{C} is any cartesian category, then the category $\mathbf{Ab}(\mathcal{C})$ of abelian groups of \mathcal{C} is preadditive. Idem for the category $A\text{-Mod}(\mathcal{C})$ of A -modules of \mathcal{C} (over some ring A of \mathcal{C}).
 5. As a particular case, the category **AbTop**, **AbHaus** or **AbCHaus** of topological (resp. Hausdorff, resp. compact Hausdorff) abelian groups are preadditive.
 6. The category **A** with only one object, whose morphisms are the elements of a ring A and composition is given by multiplication ($g \circ f = fg$) is preadditive. Any preadditive category with exactly one object has this form.
 7. The category \mathbf{Mat}_A whose morphisms are matrices with coefficients in A (and natural numbers as objects) is preadditive.
 8. The category $\mathbf{Vect}(X)$ of vector bundles on a manifold X is preadditive.
 9. The categories **Set**, **Top**, **Grp** or **Ring** are *not* preadditive.

Exercise 2.1.2 Show that, if \mathcal{P} is a preadditive category, then \mathcal{P}^{op} is preadditive too.

Proposition 2.1.3 If \mathcal{C} is a small category and \mathcal{P} is a preadditive category, then $\mathbf{Hom}(\mathcal{C}, \mathcal{P})$ is preadditive.

Proof. Given two natural transformations $\alpha, \beta : F \rightarrow G$, on simply sets $(\alpha + \beta)_X = \alpha_X + \beta_X$ for $X \in \mathcal{C}$. Details are left to the reader. For example, if $f : X \rightarrow Y$ is a morphism in \mathcal{C} , then distributivity in \mathcal{P} implies

$$\begin{aligned} (\alpha + \beta)_Y \circ F(f) &= \alpha_Y \circ F(f) + \beta_Y \circ F(f) \\ &= G(f) \circ \alpha_X + G(f) \circ \beta_X = G(f) \circ (\alpha + \beta)_X. \end{aligned} \quad \blacksquare$$

It follows that the category $\widehat{\mathcal{C}}(\mathcal{P}) := \mathbf{Hom}(\mathcal{C}^{\text{op}}, \mathcal{P})$ of presheaves on a small category \mathcal{C} with values in a preadditive category \mathcal{P} is also preadditive. Or else, if I is a small category, then the category $\mathcal{P}^I := \mathbf{Hom}(I, \mathcal{P})$ of diagrams of shape I in \mathcal{P} is preadditive.

Definition 2.1.4 A *preadditive subcategory* of a preadditive category \mathcal{P} is a subcategory \mathcal{P}' such that for all $M, N \in \mathcal{P}$, $\mathbf{Hom}_{\mathcal{P}'}(M, N)$ is a subgroup of $\mathbf{Hom}_{\mathcal{P}}(M, N)$.

A preadditive subcategory is preadditive. A full subcategory of a preadditive category is automatically a preadditive category.

If $M \in \mathcal{P}$, we will now write

$$h^M : \mathcal{P} \rightarrow \mathbf{Ab} \quad \text{and} \quad h_M : \mathcal{P}^{\text{op}} \rightarrow \mathbf{Ab}$$

(replacing henceforth \mathbf{Set} with \mathbf{Ab}).

Exercise 2.1.5 Show that $h^M : \mathcal{P} \rightarrow \mathbf{Ab}$ still preserves all limits and is in particular left exact (and dual).

Definition 2.1.6 A functor $F : \mathcal{P} \rightarrow \mathcal{Q}$ between two preadditive categories is *additive* if for all $M, N \in \mathcal{P}$, the map

$$\text{Hom}(M, N) \rightarrow \text{Hom}(F(M), F(N)), \quad f \mapsto F(f)$$

is a group homomorphism.

- Examples**
1. A subcategory \mathcal{P}' of a preadditive category \mathcal{P} is a preadditive subcategory if and only if there exists the structure of a preadditive category on \mathcal{P}' turning the inclusion into an additive functor.
 2. If \mathcal{P} is any preadditive category, then the functors h^M and h_M are additive.
 3. If A is a ring and M is a right A -module, then the functor $N \mapsto M \otimes_A N$ is additive.
 4. If $f : A \rightarrow B$ is a ring homomorphism, then the forgetful functor $B\text{-}\mathbf{Mod} \rightarrow A\text{-}\mathbf{Mod}$ as well as its adjoint $M \mapsto B \otimes_A M$ and coadjoint $M \mapsto \text{Hom}_A(B, N)$ are all additive.

We shall denote by $\text{Hom}_+(\mathcal{P}, \mathcal{Q})$ the collection of all additive functors.

Exercise 2.1.7 Show that if A, B are two rings, then $\text{Hom}_{\mathbf{Ring}}(A, B) \simeq \text{Hom}_+(\mathbf{A}, \mathbf{B})$.

Proposition 2.1.8 Let $F : \mathcal{P} \rightarrow \mathcal{Q}$ be an additive functor. Then,

1. if $G : \mathcal{D} \rightarrow \mathcal{E}$ is additive, so is $G \circ F$,
2. $F^{\text{op}} : \mathcal{P}^{\text{op}} \rightarrow \mathcal{Q}^{\text{op}}$ also is additive,
3. if I a small category, $F^I : \mathcal{P}^I \rightarrow \mathcal{Q}^I$ also is additive.

Proof. Clear. ■

If \mathcal{P} is a small preadditive category and \mathcal{Q} is a preadditive category, then additive functors form a (full) preadditive subcategory $\mathbf{Hom}_+(\mathcal{P}, \mathcal{Q})$ of $\mathbf{Hom}(\mathcal{P}, \mathcal{Q})$. We shall write

$$\mathcal{P}\text{-}\mathbf{Mod} := \mathbf{Hom}_+(\mathcal{P}, \mathbf{Ab}) \quad (\text{resp.} \quad \mathbf{Mod}\text{-}\mathcal{P} = \mathcal{P}^{\text{op}}\text{-}\mathbf{Mod})$$

and call them left (resp. right) \mathcal{P} -modules (this is the additive analog to $\widehat{\mathcal{C}}$).

Exercise 2.1.9 Show that $\mathbf{A}\text{-}\mathbf{Mod} \simeq A\text{-}\mathbf{Mod}$.

Exercise 2.1.10 Prove the additive Yoneda lemma: if \mathcal{P} is a small^a preadditive category, $M \in \mathcal{P}$ and $F \in \mathcal{P}\text{-}\mathbf{Mod}$, then there exists a natural isomorphism of

abelian groups $\text{Hom}(h^M, F) \simeq F(M)$ (and dual).

^aThe statement generalizes to large categories.

2.1.2 Additive category

Definition 2.1.11 An initial object in a preadditive category \mathcal{P} is called a *zero* object and denoted by 0. A coproduct of a family $(M_i)_{i \in I}$ is called a *direct sum* and denoted by $\bigoplus_{i \in I} M_i$.

Proposition 2.1.12 Let \mathcal{P} be a preadditive category. Then, for M, M_1, \dots, M_n in \mathcal{P} , the following are equivalent

1. $M \simeq M_1 \times \dots \times M_n$,
2. $M \simeq M_1 \oplus \dots \oplus M_n$,
3. there exists $p_k : M \rightarrowtail M_k$ and $i_k : M_k \rightarrowtail M$ for $k = 1, \dots, n$, such that $p_k \circ i_k = \text{Id}_{M_k}$, $p_\ell \circ i_k = 0_{M_k M_\ell}$ when $k \neq \ell = 1, \dots, n$ and $\sum_{k=1}^n i_k \circ p_k = \text{Id}_M$. Moreover, the maps $p_k : M \rightarrow M_k$ (resp. $i_k : M_k \rightarrow M$) are the structural morphisms for the product (resp. direct sum).

Proof. Since (3) is autodual and (2) is dual to (1), it is sufficient to prove that (1) \Rightarrow (3) \Rightarrow (2).

1. (1) \Rightarrow (3) : We already got the structural morphisms p_k and the existence of the morphisms i_k satisfying the commutation rules then come from the universal property of the product. It only remains to check the last equality which formally follows from $p_\ell \circ (\sum_{k=1}^n i_k \circ p_k) = p_\ell$ for $\ell = 1, \dots, n$.
2. (3) \Rightarrow (2) : If we are given $f_k : M_k \rightarrow N$ for $k = 1, \dots, n$, we then set $f = \sum_{\ell=1}^n f_\ell \circ p_\ell$. We will have for $k = 1, \dots, n$, $f \circ i_k = \sum_{\ell=1}^n f_\ell \circ p_\ell \circ i_k = f_k$. Assume conversely that $f : M \rightarrow N$ satisfies $f \circ i_k = f_k$ for $k = 1, \dots, n$. We will then have

$$\sum_{\ell=1}^n f_\ell \circ p_\ell = \sum_{\ell=1}^n f \circ i_\ell \circ p_\ell = f \circ \sum_{\ell=1}^n i_\ell \circ p_\ell = f. \quad \blacksquare$$

- Examples**
1. ($n = 0$) A final object M is the same thing as a zero object (meaning an initial object) and characterized by $\text{Id}_M = 0_M$ or equivalently $\text{End}(M) = 1 = \{0\}$.
 2. ($n = 2$) A product M of two objects M_1 and M_2 is the same thing as a direct sum (meaning a coproduct) and characterized by the existence of $p_1 : M \rightarrowtail M_1$, $p_2 : M \rightarrowtail M_2$, $i_1 : M_1 \rightarrowtail M$ and $i_2 : M_2 \rightarrowtail M$ such that

$$\begin{cases} p_1 \circ i_1 = \text{Id}_{M_1}, \\ p_2 \circ i_2 = \text{Id}_{M_2} \end{cases}, \quad \begin{cases} p_2 \circ i_1 = 0_{M_1 M_2}, \\ p_1 \circ i_2 = 0_{M_2 M_1} \end{cases} \quad \text{and} \quad i_1 \circ p_1 + i_2 \circ p_2 = \text{Id}_M. \quad (2.1)$$

Definition 2.1.13 An *additive* category is a cartesian preadditive category.

It means that the preadditive category admits finite products and this is equivalent to admitting finite direct sums in which case they are the same. Equivalently, by induction, it admits a final (resp. zero) object and products (resp. direct sums) of two objects.

In an additive category the natural morphism $\bigoplus_{i \in I} M_i \rightarrow \prod_{i \in I} M_i$ is an isomorphism when I is finite (but not otherwise) and we shall mostly use the first notation.

- Examples**
1. The categories **Ab**, **A-Mod**, **Mat_A** and **Vect(X)** are additive.
 2. If \mathcal{C} is a small (resp. a cartesian) category, then $\widehat{\mathcal{C}}(\mathbf{Ab})$ (resp. **Ab**(\mathcal{C}) or **A-Mod**(\mathcal{C})) is additive.
 3. If \mathcal{P} is a small preadditive category, then $\mathcal{P}\text{-Mod}$ is additive.
 4. If A is a non-zero ring, then the category **A** is *not* an additive category.

Exercise 2.1.14 Show that, if \mathcal{P} is an additive category, then \mathcal{P}^{op} is additive too. Show also that, if I is a small category, then \mathcal{P}^I is additive.

Lemma 2.1.15 In an additive category, if $M = \bigoplus_{k=1}^m M_k$ and $N = \bigoplus_{j=1}^n N_j$, then

$$\text{Hom}(M, N) \simeq \bigoplus_{1 \leq j \leq n, 1 \leq k \leq m} \text{Hom}(M_k, N_j).$$

Proof. Follows from universal properties of direct sum which is at the same time a product and a coproduct. ■

We shall use matrix notation and write $f = [f_{jk}]$ with $f_{jk} = p_j \circ f \circ i_k : M_k \rightarrow N_j$. Composition is then reduced to matrix multiplication:

Exercise 2.1.16 Show that, if we are given two morphisms $f : M \rightarrow N$ and $g : N \rightarrow P$ in an additive category, and decompositions $M = \bigoplus_{k=1}^m M_k$ and $N = \bigoplus_{j=1}^n N_j$ and $P = \bigoplus_{i=1}^p N_p$, then $g \circ f = [h_{ik}]$ with $h_{ik} = \sum_j g_{ij} \circ f_{jk}$.

Example

1. In the case $M = M_1 \oplus M_2$, we have $\text{Id}_M = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$, $i_1 = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$, $i_2 = \begin{bmatrix} 0 \\ 1 \end{bmatrix}$, $p_1 = \begin{bmatrix} 1 & 0 \end{bmatrix}$ and $p_2 = \begin{bmatrix} 0 & 1 \end{bmatrix}$ and we can check identities 2.1 by matrix operations, the last one reading

$$\begin{bmatrix} 1 \\ 0 \end{bmatrix} \begin{bmatrix} 1 & 0 \end{bmatrix} + \begin{bmatrix} 0 \\ 1 \end{bmatrix} \begin{bmatrix} 0 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}.$$

2. If $p_1, p_2 : M \oplus M \rightarrow M$ and $i_1, i_2 : M \rightarrow M \oplus M$ denote the structural morphisms, and Δ denotes the *diagonal morphism* ($p_1 \circ \Delta = p_2 \circ \Delta = \text{Id}_M$), then

$$\Delta = \begin{bmatrix} 1 \\ 1 \end{bmatrix} = i_1 + i_2 \quad (\text{and dually } \nabla = \begin{bmatrix} 1 & 1 \end{bmatrix} = p_1 + p_2).$$

Proposition 2.1.17 If \mathcal{P} is an additive category, then the preadditive structure is unique.

Proof. If $f_1, f_2 : M \rightarrow N$, then

$$f_1 + f_2 = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} f_1 \\ f_2 \end{bmatrix}. \quad (2.2)$$

Exercise 2.1.18 Show that the forgetful functor $\mathbf{Ab}(\mathcal{P}) \rightarrow \mathcal{P}$ is an isomorphism when \mathcal{P} is additive.

Solution. If an object M of \mathcal{P} is endowed with an abelian group law $\begin{bmatrix} f & g \end{bmatrix} : M \oplus M \rightarrow M$, then the neutral element for this law is necessarily the unique map $0 : 0 \rightarrow M$ and the following diagram is therefore required to commute:

$$\begin{array}{ccc} M & \xrightarrow{\begin{bmatrix} 0 \\ 1 \end{bmatrix}} & M \oplus M \\ \begin{bmatrix} 1 \\ 0 \end{bmatrix} \downarrow & \searrow 1 & \downarrow \begin{bmatrix} f & g \end{bmatrix} \\ M \oplus M & \xrightarrow{\begin{bmatrix} f & g \end{bmatrix}} & M. \end{array}$$

It follows that $f = g = 1$ and the law is therefore the codiagonal morphism ∇ . One easily checks that, conversely, ∇ defines an abelian group structure on M and that this is functorial. This provides an inverse to the forgetful functor. ■

Exercise 2.1.19 Show that a category \mathcal{P} is additive if and only if there exists a (not unique) cartesian category \mathcal{C} such that $\mathbf{Ab}(\mathcal{C}) \simeq \mathcal{P}$.

Definition 2.1.20 An *additive subcategory* of an additive category is a preadditive subcategory which is additive.

Exercise 2.1.21 Show that

1. a subcategory \mathcal{P}' of an additive category \mathcal{P} is an additive subcategory if and only if it is stable under finite direct sums: if $M_1, \dots, M_n \in \mathcal{P}'$, then $\bigoplus_{k=1}^n M_k \in \mathcal{P}'$.
2. a functor $F : \mathcal{P} \rightarrow \mathcal{Q}$ between additive categories is additive if and only if it preserves finite direct sums: if $M_1, \dots, M_n \in \mathcal{P}$, then $F(\bigoplus_{k=1}^n M_k) = \bigoplus_{k=1}^n F(M_k)$.

Exercise 2.1.22 Let \mathcal{Q} be an additive category. Show that, if a fully faithful functor $\mathcal{P} \hookrightarrow \mathcal{Q}$ has an adjoint or a coadjoint, then \mathcal{P} also is additive.

Exercise 2.1.23 Show that, if \mathcal{P} is a small preadditive category and \mathcal{Q} is additive, then $\mathbf{Hom}_+(\mathcal{P}, \mathcal{Q})$ is additive.

Proposition 2.1.24 If a functor F between two additive categories is adjoint to a functor G , then both functors are additive and there exists a natural isomorphism of abelian groups

$$\mathbf{Hom}(FM, N) \simeq \mathbf{Hom}(M, GN). \quad (2.3)$$

Proof. Since an adjoint (resp. coadjoint) preserves all colimits (resp. all limits), it preserves finite direct sums and is therefore additive. Now, given two morphisms $f, g : FM \rightarrow N$, we have

$$f = \begin{bmatrix} 1 & 0 \end{bmatrix} \begin{bmatrix} f \\ g \end{bmatrix}, g = \begin{bmatrix} 0 & 1 \end{bmatrix} \begin{bmatrix} f \\ g \end{bmatrix} \quad \text{and} \quad f + g = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} f \\ g \end{bmatrix}.$$

Thus, if we denote by Φ the natural isomorphism (2.3), we shall have first

$$\Phi(f) = \begin{bmatrix} 1 & 0 \end{bmatrix} \Phi \left(\begin{bmatrix} f \\ g \end{bmatrix} \right) \quad \text{and} \quad \Phi(g) = \begin{bmatrix} 0 & 1 \end{bmatrix} \Phi \left(\begin{bmatrix} f \\ g \end{bmatrix} \right)$$

from which we derive

$$\Phi \left(\begin{bmatrix} f \\ g \end{bmatrix} \right) = \begin{bmatrix} \Phi(f) \\ \Phi(g) \end{bmatrix}.$$

It follows that

$$\Phi(f + g) = \begin{bmatrix} 1 & 1 \end{bmatrix} \Phi \left(\begin{bmatrix} f \\ g \end{bmatrix} \right) = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} \Phi(f) \\ \Phi(g) \end{bmatrix} = \Phi(f) + \Phi(g). \quad \blacksquare$$

2.1.3 Preabelian category

We fix an additive category \mathcal{P} .

Definition 2.1.25 The *kernel* of a morphism $f : M \rightarrow N$ is $\ker f := \ker(f, 0)$ if it exists. The dual notion is that of a *cokernel* $\operatorname{coker} f$.

Be careful that the kernel actually comes with a morphism $i : \ker f \rightarrowtail M$ and is only defined up to isomorphism (and dual). It has the following universal property: $f \circ i = 0$ and given any $g : M' \rightarrow M$ such that $f \circ g = 0$, there exists a unique morphism $\bar{g} : M' \rightarrow \ker f$ such that $i \circ \bar{g} = g$ (and dual):

$$\begin{array}{ccc} \ker f & \xrightarrow{i} & M \\ & \nwarrow \bar{g} & \uparrow g \\ & & M' \end{array} \quad \begin{array}{ccc} M & \xrightarrow{f} & N \\ & \searrow 0 & \downarrow g \\ & & N' \end{array} \quad \begin{array}{ccc} & & \nearrow \bar{g} \\ & & \end{array} \quad \text{(and} \quad \begin{array}{ccc} M & \xrightarrow{f} & N \xrightarrow{p} \operatorname{coker} f \\ & \searrow 0 & \downarrow g \\ & & N' \end{array}$$

Definition 2.1.25 has a kind of converse in the sense that $\ker(f, g) = \ker(g - f)$ when $f, g : M \rightarrow N$ (and dual).

Definition 2.1.26 A *preabelian category* is an additive category with all kernels and cokernels.

Equivalently, it is a preadditive category which is finitely bicomplete. Up to the end of the section, you can assume that \mathcal{P} is preabelian or else follow our convention to implicitly assume that a limit or a colimit exists when we write it down.

Proposition 2.1.27 If $f : M \rightarrow N$ is a morphism, then

1. f is a monomorphism if and only if $\ker f = 0$ (and dual),
2. $f = 0$ if and only if $\ker f = M$ (and dual),
3. f is a monomorphism and $f = 0$ if and only if $M = 0$ (and dual).

Proof. If f is a monomorphism and $g : M' \rightarrow M$ satisfies $f \circ g = 0$ ($= f \circ 0$), then $g = 0$ and therefore, g factorizes through 0 (uniquely). Conversely, if $\ker f = 0$ and $g, h : M' \rightarrow M$ satisfy $f \circ g = f \circ h$, then $f \circ (g - h) = 0$ and therefore $g - h$ factor through $\ker f = 0$ which implies that $g - h = 0$ and finally $g = h$.

If $f = 0$, then *any* morphism $g : M' \rightarrow M$ satisfies $g \circ f = 0$ so that $\ker f = M$. Conversely, if $\ker f = M$, then $f = f \circ \operatorname{Id}_M = 0$.

The last assertion follows from the first two. \blacksquare

Definition 2.1.28 1. A sequence

$$0 \rightarrow M' \rightarrow M \xrightarrow{f} M'' \quad (2.4)$$

is said to be *(left) exact* if the sequence

$$M' \longrightarrow M \xrightarrow[0]{f} M''$$

is left exact. The dual notion is that of a *(right) exact sequence* $M' \rightarrow M \rightarrow M'' \rightarrow 0$.

2. A *short exact sequence*^a is a sequence

$$0 \rightarrow M' \xrightarrow{i} M \xrightarrow{p} M'' \rightarrow 0 \quad (2.5)$$

which is exact both on the left and on the right.

^aBe careful that there exists alternative weaker definitions at this level.

In other words, a sequence (2.4) is (left) exact if and only if $M' = \ker f$ (and dual) and a sequence (2.5) is exact if and only if $M' = \ker p$ and $\operatorname{coker} i = M''$.

Exercise 2.1.29 Show that

1. a sequence $0 \rightarrow M \xrightarrow{f} N \rightarrow 0$ is left exact if and only if it is right exact^a if and only if f is an isomorphism,
2. if $0 \rightarrow M' \xrightarrow{f} M \xrightarrow{g} M'' \rightarrow 0$ is left exact (resp. short exact), then f is an isomorphism if and only if $g = 0$ (resp. $M'' = 0$) – and dual.

^aSo that there is no ambiguity in saying that such a sequence is exact.

Solution. 1. By duality, only the case of a left exact sequence requires our attention. Now, this is the kind of situation where we have to take into account the fact that a kernel is only defined up to a unique isomorphism: the sequence is left exact if and only if f induces an isomorphism between M and the kernel of $N \rightarrow 0$ which is exactly N .

2. We know that $g = 0$ if and only if $\ker g = M$ which means that f is an isomorphism when the sequence is left exact. If the sequence is moreover right exact, then g is an epimorphism and zero at the same time which means that $M'' = 0$. ■

Exercise 2.1.30 Show that any commutative diagram with exact rows

$$\begin{array}{ccccccc} M' & \longrightarrow & M & \longrightarrow & M'' & \longrightarrow & 0 \\ & \downarrow & \downarrow & & \downarrow & & \\ 0 & \longrightarrow & N' & \longrightarrow & N & \longrightarrow & N'' \end{array}$$

may be uniquely completed as described (and dual).

Left (resp. right, resp. short) exact sequences in \mathcal{P} form a full additive subcategory

of the category $\mathcal{P}^{[2]}$ of all diagrams of the form $M' \rightarrow M \rightarrow M''$ so that the zero on the left (resp. the right, resp. both sides) is only a decoration to specify what kind of sequence we are considering.

Proposition 2.1.31 Limits of left exact sequences are computed termwise (and dual).

Proof. Limits always preserve kernels since they preserve all limits. ■

It means that, if we are given a commutative diagram of short exact sequences, then

$$\varprojlim_I (0 \rightarrow M'_i \rightarrow M_i \rightarrow M''_i) = 0 \rightarrow \varprojlim_I M'_i \rightarrow \varprojlim_I M_i \rightarrow \varprojlim_I M''_i,$$

and in particular implies that the left hand side is a left exact sequence. Be careful however that the question of limits or colimits of *short* exact sequences is way more complicated.

Corollary 2.1.32 If all columns and both bottom rows are exact in a commutative diagram

$$\begin{array}{ccccccc} & & 0 & & 0 & & 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & K' & \longrightarrow & K & \longrightarrow & K'' \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & M' & \longrightarrow & M & \longrightarrow & M'' \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & N' & \longrightarrow & N & \longrightarrow & N'', \end{array}$$

then the top row also is exact (and dual). ■

It is important to notice that there is no similar assertion for short exact sequences and that it will be necessary to introduce later long exact sequences to solve this problem.

Exercise 2.1.33 Show that, in $A\text{-Mod}$, any short exact sequence is isomorphic to some

$$0 \rightarrow M' \rightarrow M \rightarrow M/M' \rightarrow 0$$

where M' is a submodule of M .

Definition 2.1.34 A short exact sequence is said to *split* if it is isomorphic to

$$0 \longrightarrow M_1 \xrightarrow{i_1} M_1 \oplus M_2 \xrightarrow{p_2} M_2 \longrightarrow 0.$$

Proposition 2.1.35 A short exact sequence

$$0 \longrightarrow M' \xrightarrow{i} M \xrightarrow{p} M'' \longrightarrow 0$$

splits if and only if p is a split epimorphism (and dual).

Proof. The condition is clearly necessary (use i_2). Assume conversely given some $s : M'' \rightarrow M$ such that $p \circ s = 1$. Since $p \circ (1 - s \circ p) = 0$, there exists a unique $r : M \rightarrow M'$ such that $i \circ r = 1 - s \circ p$. And we have in particular $i \circ r + s \circ p = 1$. Since we already know that $p \circ i = 0$ and $p \circ s = 1$, it only remains to check that $r \circ s = 0$ and $r \circ i = 1$. Since i is a monomorphism, it is sufficient to notice that $i \circ (r \circ s) = (1 - s \circ p) \circ s = 0$ and $i \circ (r \circ i) = (1 - s \circ p) \circ i = i$. ■

It is implicit in the previous argument that we have an isomorphism of short exact sequences with given by $\begin{bmatrix} r \\ p \end{bmatrix} : M \xrightarrow{\sim} M' \oplus M''$ with inverse $\begin{bmatrix} i & s \end{bmatrix}$.

Exercise 2.1.36 Show that a diagram

$$\begin{array}{ccc} M' & \xrightarrow{f'} & N' \\ \downarrow g' & & \downarrow g \\ M & \xrightarrow{f} & N \end{array}$$

is cartesian (resp. cocartesian, resp. both cartesian and cocartesian) if and only if the sequence

$$0 \longrightarrow M' \xrightarrow{\begin{bmatrix} f' \\ g' \end{bmatrix}} M \oplus N' \xrightarrow{\begin{bmatrix} f & g \end{bmatrix}} N \longrightarrow 0$$

is left exact (resp. right exact, resp. short exact).

The following provides a very nice alternative definition for (short) exact sequences:

Exercise 2.1.37 Show that a sequence

$$0 \rightarrow M' \rightarrow M \rightarrow M'' \rightarrow 0$$

is left exact (resp. right exact, resp. short exact) if and only if the diagram

$$\begin{array}{ccc} M' & \longrightarrow & M \\ \downarrow & & \downarrow \\ 0 & \longrightarrow & M'' \end{array}$$

is cartesian (resp. cocartesian, resp. both cartesian and cocartesian).

Proposition 2.1.38 If a diagram

$$\begin{array}{ccc} M' & \xrightarrow{f'} & N' \\ \downarrow g' & \lrcorner & \downarrow g \\ M & \xrightarrow{f} & N \end{array}$$

is cartesian, then $\ker f' = \ker f$ and $\ker g' = \ker g$ (and dual).

Solution. There exists a sequence of cartesian diagrams

$$\begin{array}{ccccc} \ker g' & \longrightarrow & M' & \xrightarrow{f'} & N' \\ \downarrow & \lrcorner & \downarrow g' & \lrcorner & \downarrow g \\ 0 & \longrightarrow & M & \xrightarrow{f} & N \end{array}$$

showing thanks to exercise 1.3.15 that the rectangle is cartesian and therefore $\ker g' \simeq \ker g$. The original statement is obtained by symmetry. ■

Exercise 2.1.39 Show that if A is a commutative ring, M is an A -module and $f, g \in A$ are comaximal (the ideal generated by f and g is the full ring) then the sequence

$$0 \rightarrow M \rightarrow M_f \times M_g \xrightarrow{\bar{}} M_{fg} \rightarrow 0$$

is exact.

Solution. The condition means that there exists $h_1, h_2 \in A$ such that $h_1 f_1 + h_2 f_2 = 1$. Before going any further, note that f^n and g^m are always comaximal when $n, m \geq 0$ and that $M_f = M_{f^n}$ when $n > 0$. We now turn to the proof. First of all, composition is clearly zero. Let us show that the first map is injective: if the image of $s \in M$ is zero in M_{f_i} for $i = 1, 2$, it means that there exists $n_i \in \mathbb{N}$ such that $f_i^{n_i} s = 0$. Up to replacing f_i by $f_i^{n_i}$, we may assume that $f_i s = 0$ and we will have

$$s = 1s = h_1 f_1 s + h_2 f_2 s = 0.$$

We give now ourselves $s_i / f_i^{n_i} \in M_{f_i}$ for $i = 1, 2$ with same image in $M_{f_1 f_2}$ and we have to find $s \in M$ whose image in M_{f_i} is $s_i / f_i^{n_i}$. We may again assume that $n_i = 1$ and our hypothesis means that $(f_1 f_2)^m (f_1 s_2 - f_2 s_1) = 0$ for some m . Actually, up to replacing f_i by f_i^{m+1} and s_i by $f_i^m s_i$, we may assume that $f_1 s_2 - f_2 s_1 = 0$. It is then sufficient to set $s := h_1 s_1 + h_2 s_2$. We will have

$$f_1 s = f_1 (h_1 s_1 + h_2 s_2) = (h_1 f_1 + h_2 f_2) s_1 = s_1$$

and symmetrically for s_2 . It remains to show that the last map is surjective: we give ourselves $s / (f_1 f_2)^n \in M_{f_1 f_2}$ and we show that it comes from $M_{f_1} \times M_{f_2}$. Again, we may assume that $n = 1$ and consider $(h_2 s / f_1, -h_1 s / f_2)$. We will have

$$\frac{h_2 s}{f_1} - \frac{-h_1 s}{f_2} = \frac{h_2 f_2 s + h_1 f_1 s}{f_1 f_2} = \frac{s}{f_1 f_2}. \quad \blacksquare$$

2.1.4 Abelian category

Definition 2.1.40 An *abelian category* is a preabelian category where all monomorphisms and all epimorphisms are regular.

The condition means that any monomorphism is a kernel and any epimorphism is a cokernel (the converse being always true).

- Examples**
1. The category **Ab** of abelian groups is abelian.
 2. The category $A\text{-Mod}$ of left modules over a ring A is abelian.
 3. More generally, if \mathcal{P} is a small preadditive category, then the category $\mathcal{P}\text{-Mod}$ of additive functors from \mathcal{P} to **Ab** is abelian.
 4. Also, if \mathcal{C} is a small category, then the category $\widehat{\mathcal{C}}(\mathbf{Ab})$ of presheaves of abelian groups on \mathcal{C} is abelian.
 5. The category $\mathbf{Op}(k\text{-Mod})$ of k -modules endowed with an operator is abelian.
 6. The category **AbCHaus** of compact Hausdorff abelian groups is abelian.
 7. The category **AbTop** is not abelian however because the identity $\mathbb{Z}^{\text{disc}} \rightarrow \mathbb{Z}^{\text{coarse}}$ is not a regular monomorphism.
 8. The category **AbHaus** is not abelian either because the map $\ell^1(\mathbb{R}) \rightarrow \ell^2(\mathbb{R})$ is not a regular monomorphism.
 9. The category **Mat** $_{\mathbb{Z}}$ is not abelian because $[2]$ is not a regular monomorphism.

Exercise 2.1.41 A *representation* of a monoid G (or a G -module) is a morphism $\rho : G \rightarrow \text{End}_k(M)$ where M is a module on some fixed commutative ring k . A morphism of representations is a k -linear map $M \rightarrow N$ which is compatible with the actions of G . Show that they form an abelian category $\mathbf{Rep}_k(G)$.

Hint. One can check that $\mathbf{Rep}_k(G) \simeq A\text{-Mod}$ with $A = k \cdot G$. ■

Exercise 2.1.42 Let $\mathcal{O} \subset \mathbb{C}[[t]]$ be the ring of power series with positive radius of convergence. A *derivation* on an \mathcal{O} -module M is a group homomorphism $\partial_M : M \rightarrow M$ such that $\forall s \in M, \partial(fs) = f's + f\partial_M(s)$. An \mathcal{O} -linear map $f : M \rightarrow N$ between modules with derivations is said to be *horizontal* if it commutes with the derivations. Show that modules with derivations and horizontal maps form an abelian category.

Hint. One can check that this category is equivalent to the category $\mathcal{D}\text{-Mod}$ where

$$\mathcal{D} := \left\{ \sum_{k=0}^d f_k \partial^k, f_k \in \mathcal{O} \right\}$$

is the (non-commutative) ring of *differential operators* with the commutation rule

$$[\partial, f] := \partial f - f\partial = f'. \quad \blacksquare$$

Proposition 2.1.43 If \mathcal{A} is an abelian category, then \mathcal{A}^{op} is also an abelian category, as well as \mathcal{A}^I if I is a small category.

Proof. For the second assertion, use proposition 1.3.39. ■

Definition 2.1.44 An *abelian subcategory* of an abelian category is a *full* subcategory which is stable under all finite limits and finite colimits.

Exercise 2.1.45 Show that an abelian subcategory is indeed an abelian category.

Exercise 2.1.46 Show that a functor $F : \mathcal{A} \rightarrow \mathcal{B}$ between two preabelian categories is left exact if and only if it is additive and preserves left exact sequences (and dual).

Exercise 2.1.47 Let \mathcal{B} be an abelian category. Show that, if a fully faithful functor $\mathcal{A} \hookrightarrow \mathcal{B}$ has an exact adjoint or coadjoint, then \mathcal{A} also is abelian.

Exercise 2.1.48 Show that the category of free abelian groups of finite rank (or equivalently $\mathbf{Mat}_{\mathbb{Z}}$) is preabelian but that the inclusion in the category of all abelian groups is not an exact functor.

Hint. Everything works fine besides the fact that the cokernel Q in \mathbf{Ab} of a morphism of free abelian groups of finite rank may not be free. However, one can write $Q = F \oplus T$ with F free and T torsion. An then, the cokernel in the original category will be F as one easily checks. ■

Lemma 2.1.49 In a preabelian category, a morphism $f : M \rightarrow N$ is a regular monomorphism if and only if

$$0 \rightarrow M \xrightarrow{f} N \rightarrow \operatorname{coker} f \rightarrow 0$$

is a *short*^a exact sequence (and dual).

^aRecall that we mean that it is simultaneously left and right exact.

Proof. The condition is clearly sufficient. Conversely, if f is the kernel of a morphism $g : N \rightarrow P$, we can then build the following commutative diagram (using the universal property of the kernel):

$$\begin{array}{ccccc}
 M = \ker g & & & & \operatorname{coker} f \\
 \uparrow & \searrow f & & \nearrow p & \downarrow \\
 & & N & & \\
 \downarrow & \nearrow & \searrow f & & \\
 \ker p & & & & P.
 \end{array}$$

As shown in exercise 1.3.25, both left hand morphisms are necessarily inverse to each other, showing that $M = \ker p$. ■

Recall the notion of a balanced category from definition 1.3.33: a morphism which is at the same time a monomorphism and an epimorphism is automatically an isomorphism.

Proposition 2.1.50 An abelian category is balanced.

Proof. If f is an epimorphism, then $\text{coker } f = 0$. If f is also a monomorphism, then it is regular since we assumed our category to be abelian and it then follows from lemma 2.1.49 that the sequence $0 \rightarrow M \xrightarrow{f} N \rightarrow 0$ is (left) exact. It means that f is an isomorphism. ■

Recall that we introduced the notions of regular image and coimage in definition 1.3.31.

Proposition 2.1.51 If $f : M \rightarrow N$ is a morphism in a (pre) abelian category, then the sequence

$$0 \rightarrow \text{im } f \rightarrow N \rightarrow \text{coker } f \rightarrow 0$$

is (left) exact (and dual).

Proof. According to proposition 1.3.30, for left exactness, it is sufficient to show that $\ker(N \xrightarrow{\pi} \text{coker } f)$ is universal for factorization of f through a regular monomorphism $j : N' \rightarrow N$. We can then simply contemplate the following commutative diagram:

$$\begin{array}{ccccccc}
 0 & \longrightarrow & N' & \xrightarrow{j} & N & \longrightarrow & \text{coker } j \longrightarrow 0 \\
 & & \uparrow & \swarrow & \parallel & & \uparrow \\
 & & & \ker \pi & & & \\
 & & \swarrow & & \searrow & & \\
 M & \xrightarrow{f} & N & \xrightarrow{\pi} & \text{coker } f & \longrightarrow & 0.
 \end{array}$$

If the category is abelian, then π is a regular epimorphism with $\ker \pi = \text{im } f$ and the sequence is therefore also right exact. ■

Corollary 2.1.52 A morphism $f : M \rightarrow N$ in a preabelian category is an epimorphism if and only if $\text{im } f = N$ (and dual).

Proof. We have $\text{im}(f) = N$ if and only if $\pi : N \rightarrow \text{coker } f$ is the zero morphism. But since π is an epimorphism, it means that $\text{coker } f = 0$ or equivalently that f is an epimorphism. ■

Recall that the notion of a *strict morphism* was introduced in definition 1.3.31: it means that regular image and regular coimage coincide.

Proposition 2.1.53 A preabelian category is abelian if and only if all morphisms are strict.

Proof. We already know from exercise 1.3.32 that a strict monomorphism/epimorphism is automatically regular.

Conversely, any morphism $f : M \rightarrow N$ splits as $M \xrightarrow{\pi} \text{im}(f) \rightarrow N$. Now π splits in turn as $M \rightarrow \text{im}(\pi) \rightarrow \text{im}(f)$. Since the morphism $\text{im}(\pi) \rightarrow N$ is

still a monomorphism and all monomorphisms are regular in an abelian category, proposition 1.3.30 implies that $\text{im}(\pi) = \text{im}(f)$. Corollary 2.1.52 implies that π is an epimorphism and we are done. More precisely, this implies that $\text{coim } f \rightarrow \text{im } f$ is an epimorphism. The dual argument shows that this is a monomorphism and we know from proposition 2.1.50 that an abelian category is balanced. ■

Corollary 2.1.54 In an abelian category, any morphism $f : M \rightarrow N$ factors uniquely up to an isomorphism as an epimorphism followed by a monomorphism. ■

Most statements about abelian categories can be proven by reducing to a category of modules using the following embedding theorem¹:

Theorem 2.1.55 — Freyd-Mitchell. If \mathcal{A} is a small abelian category, then there exists a ring A and a fully faithful exact functor $\mathcal{A} \hookrightarrow A\text{-Mod}$.

Proof. To do. ■

In other words, a small category is abelian if and only if it is equivalent to an abelian subcategory of a category of modules.

2.1.5 Exactness

Up to the end of this section, we work in a fixed abelian category \mathcal{A} .

Definition 2.1.56 A sequence $M' \xrightarrow{f} M \xrightarrow{g} M''$ is said to be *exact in the middle* or *exact in M* or even simply *exact* if $\text{im } f = \ker g$.

Example A sequence $0 \rightarrow M' \xrightarrow{f} M$ is exact if and only if f is a monomorphism (and dually, a sequence $M \xrightarrow{f} M'' \rightarrow 0$ is exact if and only if f is an epimorphism).

Exercise 2.1.57 Show that the condition for being exact is autodual: if $\text{coker } f = \text{im } g$, then $\text{im } f = \ker g$.

Exercise 2.1.58 Show that a sequence

$$0 \longrightarrow M' \xrightarrow{i} M \xrightarrow{p} M'' \longrightarrow 0$$

is left exact (resp. right exact, resp. short exact) if and only if it is exact in M' , M (resp. M, M'' , resp. M', M, M'').

Exercise 2.1.59 Show that a sequence

$$0 \longrightarrow M' \longrightarrow M \longrightarrow M'' \longrightarrow 0$$

is short exact if and only if it is left exact and p is an epimorphism (and dual).

Exercise 2.1.60 Let \mathcal{B} be another abelian category and $F : \mathcal{A} \rightarrow \mathcal{B}$ an additive functor. Show that the following are equivalent:

1. F is exact,

¹Feel free to use it if you are too lazy to draw big diagrams.

2. if $M' \rightarrow M \rightarrow M''$ is exact (in the middle), then $F(M') \rightarrow F(M) \rightarrow F(M'')$ is exact in the middle,
3. F preserves short exact sequences.

Proposition 2.1.61 If $f : M \rightarrow N$ is an epimorphism, then any cartesian diagram

$$\begin{array}{ccc} M' & \xrightarrow{f'} & N' \\ \downarrow & \lrcorner & \downarrow \\ M & \xrightarrow{f} & N \end{array}$$

is also cocartesian (and dual).

Proof. We know from exercise 2.1.36 that the sequence

$$0 \rightarrow M' \rightarrow M \oplus N' \rightarrow N \rightarrow 0$$

is *left* exact. Moreover, the composite morphism $M \oplus N' \rightarrow M \rightarrow N$ is an epimorphism. This shows that the sequence is also *right* exact and therefore (exercise 2.1.36 again) that the square is cocartesian. ■

Corollary 2.1.62 Epimorphisms are stable under pull back (and dual).

Proof. Use for example (the dual version of) exercise 1.3.26. ■

It is sometimes hard to provide a formal proof that a given sequence is exact and the next criterion will prove itself quite valuable:

Exercise 2.1.63 Show that a sequence $M' \xrightarrow{f} M \xrightarrow{g} M''$ is exact in M if and only if $g \circ f = 0$ and, given any $h : L \rightarrow M$ such that $g \circ h = 0$, there exists an epimorphism $\pi : L' \rightarrow L$ such that $h \circ \pi$ factors through f :

$$\begin{array}{ccccc} L' & \overset{\pi}{\dashrightarrow} & L & & \\ \downarrow & & \downarrow h & \searrow 0 & \\ M' & \xrightarrow{f} & M & \xrightarrow{g} & M'' \\ & \searrow 0 & & & \end{array}$$

Solution. If the sequence $M' \xrightarrow{f} M \xrightarrow{g} M''$ is exact, we can simply compose the cartesian squares

$$\begin{array}{ccccc} L' & \xrightarrow{\quad} & L & \xlongequal{\quad} & L \\ \downarrow & \lrcorner & \downarrow & \lrcorner & \downarrow h \\ M' & \longrightarrow & \operatorname{im} f = \ker g & \hookrightarrow & M \end{array}$$

and use the fact that epimorphisms are stable under pullback (corollary 2.1.62). Conversely, if the condition is satisfied, we can consider the case $L = \ker(g)$. Since f

factors through $\ker g$, there exists a commutative diagram

$$\begin{array}{ccccc} L' & \twoheadrightarrow & \ker(g) & & \\ \downarrow & \nearrow & \downarrow & \searrow 0 & \\ M' & \xrightarrow{f} & M & \xrightarrow{g} & M'' \end{array}$$

with an epimorphism upstairs. This implies that $M' \twoheadrightarrow \ker g$ is an epimorphism and therefore that $\operatorname{im} f = \ker g$ (by unique factorization). ■

Proposition 2.1.64 — Snake lemma. If

$$\begin{array}{ccccccc} 0 & \longrightarrow & M' & \xrightarrow{i} & M & \xrightarrow{p} & M'' \longrightarrow 0 \\ & & \downarrow f' & & \downarrow f & & \downarrow f'' \\ 0 & \longrightarrow & N' & \xrightarrow{j} & N & \xrightarrow{q} & N'' \longrightarrow 0 \end{array}$$

is a commutative diagram with exact rows, then there exists a natural autodual (long) exact sequence

$$0 \rightarrow \ker f' \xrightarrow{i} \ker f \xrightarrow{p} \ker f'' \xrightarrow{\delta} \operatorname{coker} f' \xrightarrow{\bar{j}} \operatorname{coker} f \xrightarrow{\bar{q}} \operatorname{coker} f'' \rightarrow 0.$$

Proof. The construction of δ is shown on the following diagram:

$$\begin{array}{ccccccc} & & & P & \twoheadrightarrow & \ker f'' & \\ & & & \downarrow & \lrcorner & \downarrow & \\ 0 & \longrightarrow & M' & \longrightarrow & M & \longrightarrow & M'' \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \longrightarrow & N' & \longrightarrow & N & \longrightarrow & N'' \longrightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ & & \operatorname{coker} f' & \twoheadrightarrow & Q & & \end{array}$$

Let us give some details. We first form the cocartesian (which is also cartesian by proposition 2.1.61) diagram

$$\begin{array}{ccc} N' & \twoheadrightarrow & N \\ \downarrow & & \downarrow \\ \operatorname{coker} f' & \twoheadrightarrow & Q. \end{array}$$

Note that exercise 2.1.38 provides us with short exact sequences $0 \rightarrow \operatorname{coker} f' \rightarrow Q \rightarrow N'' \rightarrow 0$ (so that we turned a cokernel into a kernel) and $0 \rightarrow \operatorname{im} f' \rightarrow N \rightarrow Q \rightarrow 0$

(that we shall only use in the end). We can consider the following factorization

$$\begin{array}{ccccccc}
 M' & \longrightarrow & M & \longrightarrow & M'' & \longrightarrow & 0 \\
 \downarrow f' & & \downarrow f & & \searrow \delta & & \\
 0 & \curvearrowright & N' & \longrightarrow & N & & \\
 \downarrow & & \downarrow & & \downarrow & & \\
 \text{coker } f' & \longrightarrow & Q & & & &
 \end{array}$$

and then the induced morphism

$$\begin{array}{ccccccc}
 0 & \longrightarrow & \ker f'' & \longrightarrow & M'' & \xrightarrow{f''} & N'' \\
 & & \downarrow \delta & & \downarrow \delta & & \parallel \\
 0 & \longrightarrow & \text{coker } f' & \longrightarrow & Q & \longrightarrow & N''.
 \end{array}$$

It follows from proposition 2.1.32 that $\ker \delta = \ker \delta$ (but this will only be used in the end). Dually, we have a cartesian square and a factorization

$$\begin{array}{ccc}
 P & \twoheadrightarrow & \ker f'' \\
 \downarrow & \lrcorner & \downarrow \\
 M & \twoheadrightarrow & M''
 \end{array}
 \quad \text{and} \quad
 \begin{array}{ccc}
 P & \twoheadrightarrow & M \\
 \downarrow \tilde{\delta} & & \downarrow f \\
 N' & \twoheadrightarrow & N.
 \end{array}$$

The fact that all horizontal morphisms are monomorphisms and all faces besides the left hand side are commutative in

$$\begin{array}{ccccc}
 & \ker f'' & \twoheadrightarrow & M'' & \\
 & \downarrow \delta & & \downarrow \delta & \\
 P & \twoheadrightarrow & M & \twoheadrightarrow & M'' \\
 \downarrow \tilde{\delta} & & \downarrow f & & \downarrow \delta \\
 & \text{coker } f' & \twoheadrightarrow & Q & \\
 N' & \twoheadrightarrow & N & \twoheadrightarrow & Q.
 \end{array}$$

implies that it is indeed a commutative diagram so that our definition is autodual. According to exercise 2.1.32, we are only left with the proof of exactness in $\ker f''$.

As an intermediate step, we first assume that f' is an epimorphism and we have to prove that $\ker f \rightarrow \ker f''$ also is an epimorphism. For this purpose, we consider the following commutative diagram with cartesian upleft square:

$$\begin{array}{ccccc}
 & & u & & \\
 P' & \twoheadrightarrow & P & \twoheadrightarrow & M \\
 \downarrow u' & \lrcorner & \downarrow \tilde{\delta} & & \downarrow f \\
 M' & \xrightarrow{f'} & N' & \xrightarrow{j} & N \\
 \downarrow i & & \downarrow j & & \\
 M & \xrightarrow{f} & N & &
 \end{array}$$

and set $v := u - i \circ u' : P' \rightarrow M$ so that $f \circ v = 0$. It follows that v factors through $\ker f$. Then, since $p \circ v = p \circ u$ (because $p \circ i = 0$), the diagram

$$\begin{array}{ccccc}
 & & u & & \\
 & \nearrow & & \searrow & \\
 P' & \xrightarrow{\quad} & P & \xrightarrow{\quad} & M \\
 \downarrow & & \downarrow & & \downarrow \\
 \ker f & \xrightarrow{\quad} & \ker f'' & & \\
 \downarrow & & \downarrow & & \downarrow \\
 M & \xrightarrow{p} & M'' & &
 \end{array}$$

is commutative and the middle horizontal morphism must be an epimorphism.

We now turn to the general case. We have to show that p induces an epimorphism $\ker f \twoheadrightarrow \ker \delta = \ker \underline{\delta}$. It is sufficient to apply the intermediate step (case f' surjective) to the following commutative diagram with exact rows and columns:

$$\begin{array}{ccccccc}
 & & 0 & & 0 & & \\
 & & \downarrow & & \downarrow & & \\
 & & \ker f & \longrightarrow & \ker \underline{\delta} & & \\
 & & \downarrow & & \downarrow & & \\
 0 & \longrightarrow & M' & \xrightarrow{i} & M & \xrightarrow{p} & M'' \longrightarrow 0 \\
 & & \downarrow & & \downarrow f & & \downarrow \underline{\delta} \\
 0 & \longrightarrow & \operatorname{im} f' & \longrightarrow & N & \longrightarrow & Q \longrightarrow 0.
 \end{array}$$

Exercise 2.1.65 Show that if

$$\begin{array}{ccccccc}
 M' & \longrightarrow & M & \longrightarrow & M'' & \longrightarrow & 0 \\
 \downarrow f' & & \downarrow f & & \downarrow f'' & & \\
 0 & \longrightarrow & N' & \longrightarrow & N & \longrightarrow & N''
 \end{array}$$

is a commutative diagram with exact rows, then there exists a natural autodual (long) exact sequence

$$\ker f' \rightarrow \ker f \rightarrow \ker f'' \rightarrow \operatorname{coker} f' \rightarrow \operatorname{coker} f \rightarrow \operatorname{coker} f''.$$

Exercise 2.1.66 Prove the four-lemma: if

$$\begin{array}{ccccccc}
 M_1 & \longrightarrow & M_2 & \longrightarrow & M_3 & \longrightarrow & M_4 \\
 \downarrow f_1 & & \downarrow f_2 & & \downarrow f_3 & & \downarrow f_4 \\
 N_1 & \longrightarrow & N_2 & \longrightarrow & N_3 & \longrightarrow & N_4
 \end{array}$$

is a commutative diagram with exact rows with f_1 an epimorphism and f_2, f_4 monomorphisms, then f_3 is a monomorphism (and dual).

Solution. Split into

$$\begin{array}{ccccccc} M_1 & \longrightarrow & M_2 & \longrightarrow & M'_2 & \longrightarrow & 0 \\ & & \downarrow & & \downarrow f'_2 & & \\ 0 & \longrightarrow & N'_1 & \longrightarrow & N_2 & \longrightarrow & N'_2 \longrightarrow 0 \end{array}$$

and

$$\begin{array}{ccccccc} 0 & \longrightarrow & M'_2 & \longrightarrow & M_3 & \longrightarrow & M'_4 \longrightarrow 0 \\ & & \downarrow f'_2 & & \downarrow f_3 & & \downarrow \\ 0 & \longrightarrow & N'_2 & \longrightarrow & N_3 & \longrightarrow & N_4 \end{array}$$

and apply snake lemma (version of exercise 2.1.65) to both. ■

Exercise 2.1.67 Prove the five-lemma: if

$$\begin{array}{ccccccccc} M_1 & \longrightarrow & M_2 & \longrightarrow & M_3 & \longrightarrow & M_4 & \longrightarrow & M_5 \\ \downarrow f_1 & & \simeq \downarrow f_2 & & \downarrow f_3 & & \simeq \downarrow f_4 & & \downarrow f_5 \\ N_1 & \longrightarrow & N_2 & \longrightarrow & N_3 & \longrightarrow & N_4 & \longrightarrow & N_5 \end{array}$$

is a commutative diagram with exact rows with f_1 an epimorphism and f_2, f_4 isomorphisms and f_5 a monomorphism, then f_3 is an isomorphism.

Proposition 2.1.68 Assume that, in the commutative diagram

$$\begin{array}{ccccccc} 0 & \longrightarrow & M' & \longrightarrow & M & \longrightarrow & M'' \longrightarrow 0 \\ & & \parallel & & \downarrow f & & \downarrow f'' \\ 0 & \longrightarrow & N' & \longrightarrow & N & \longrightarrow & N'' \longrightarrow 0, \end{array}$$

the bottom row is exact. Then, the top row also is exact if and only if the right hand square is cartesian (and dual).

Solution. It follows from exercise 2.1.38, proposition 2.1.62 and exercise 2.1.59 that the condition is sufficient. For the converse, we can first pull back the bottom row and apply the five lemma upstairs:

$$\begin{array}{ccccccc} 0 & \longrightarrow & M' & \longrightarrow & M & \longrightarrow & M'' \longrightarrow 0 \\ & & \parallel & & \downarrow & & \parallel \\ 0 & \longrightarrow & N'' & \longrightarrow & \widetilde{M} & \longrightarrow & M'' \longrightarrow 0 \\ & & \parallel & & \downarrow \ulcorner & & \downarrow f'' \\ 0 & \longrightarrow & N' & \longrightarrow & N & \longrightarrow & N'' \longrightarrow 0. \end{array}$$
■

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